# Geometry of the set of synchronous quantum correlations 

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## Correlations

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Figure: Alice
$n=$ number of experiments, $m=$ number of possible outcomes.


Figure: Bob

## Correlations



Figure: Alice
$n=$ number of experiments, $m=$ number of possible outcomes.

A correlation is a tuple

$$
\{p(i, j \mid x, y) \geq 0\} ; \quad i, j \leq m, \quad x, y \leq n
$$

satisfying

$$
\sum_{i, j} p(i, j \mid x, y)=1
$$

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\left\{E_{x, i}\right\}_{i=1}^{m},\left\{F_{y, j}\right\}_{j=1}^{m} \subseteq \mathfrak{A}
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where $E_{x, i} F_{y, j}=F_{y, j} E_{x, i}$. Then

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p(i, j \mid x, y)=\phi\left(E_{x, i} F_{y, j}\right)
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Set of all qc correlations: $C_{q c}(n, m)$.
Set of all quantum correlations: $C_{q}(n, m)$.
Set of all local correlations: $C_{l o c}(n, m)$.

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Each $C_{*}(n, m)$ is convex and satisfies:

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Theorem
(Junge-Navascues-Palazuelos-Perez-Garcia-Scholz-Werner, Fritz, Ozawa)
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Connes' embedding conjecture is true if and only if $\overline{C_{q}(n, m)}=C_{q c}(n, m)$ for every $n, m$.

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## Theorem (Paulsen-Severini-Stahlke-Todorov-Winter)

A correlation $p \in C_{q c}^{s}(n, m)$ iff there exists a $C^{*}$-algebra $\mathfrak{A}$, $\left\{E_{x, i}\right\}_{i=1}^{m} \subset \mathfrak{A}$, and a tracial state $\tau: \mathfrak{A} \rightarrow \mathbb{C}$ such that

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## Question

What is the geometry of $C_{q}^{s}(n, m)$ and $C_{q c}^{s}(n, m)$ ?

## Theorem (Dykema-Paulsen)

Connes' embedding conjecture is true if and only if $\overline{C_{q}^{s}(n, m)}=C_{q c}^{s}(n, m)$ for every $n, m$.

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## Theorem (R.)

The set $C_{q}^{s}(3,2)$ is closed. Moreover, if $p \in C_{q}^{s}(3,2)$, then there exists a $\mathfrak{A} \subset \mathbb{M}_{16}$, projection valued measures $\left\{E_{x, i}\right\} \subset \mathfrak{A}$ and a trace $\tau$ such that

$$
p(i, j \mid x, y)=\tau\left(E_{x, i} E_{y, j}\right)
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When $m=2$,

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p(i, j \mid x, x)=\left(\begin{array}{cc}
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\end{array}\right), \\
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$$

For each $\left(r_{1}, r_{2}, r_{3}\right) \in[0,1]^{3}$, we will determine the corresponding set of $\left\{\left(w_{1,2}, w_{1,3}, w_{2,3}\right)\right\} \subseteq \mathbb{R}^{3}$, denoted $S_{\vec{r}}\left[C_{q}^{s}(3,2)\right]$.

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\begin{aligned}
S_{d}\left(n_{1}, n_{2}, n_{3}\right) & :=\left\{\frac{1}{d}\left(\operatorname{Tr}\left(E_{1} E_{2}\right), \operatorname{Tr}\left(E_{1} E_{3}\right), \operatorname{Tr}\left(E_{2} E_{3}\right)\right): \operatorname{Tr}\left(E_{x}\right)=n_{x}\right\} \\
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Goal: Describe $S_{d}\left(n_{1}, n_{2}, n_{3}\right) \subseteq S_{\left(n_{1} / d, n_{2} / d, n_{3} / d\right)}\left[C_{q}^{s}(3,2)\right]$.

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Set $S_{d}(n):=S_{d}(n, n, n)$.

## Theorem

For every $n, S_{2 n}(n)=S_{2}(1)=S_{(.5,5, .5)}\left[C_{q}^{s}(3,2)\right]$ is an affine image of the $3 \times 3$ elliptope.

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Need to determine the geometry of $S_{d}\left(n_{1}, n_{2}, n_{3}\right)$, for all $n_{1}, n_{2}, n_{3} \leq d$.

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## Lemma

Assume $n_{1}+n_{2}<d$. Then

$$
S_{d}\left(n_{1}, n_{2}, n_{3}\right) \subseteq \frac{d-1}{d} \operatorname{co}\left\{S_{d-1}\left(n_{1}, n_{2}, n_{3}\right), S_{d-1}\left(n_{1}, n_{2}, n_{3}-1\right)\right\}
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(2) Calculate the closure of $\cup S_{d}\left(n_{1}, n_{2}, n_{3}\right)$. This is equal to $\overline{C_{q}^{s}(3,2)}$.
(3) Observe that every correlation in $\overline{C_{q}^{s}(3,2)}$ can be realized with $\mathfrak{A} \subseteq \mathbb{M}_{16}$.

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## Theorem (R.)

Assume $r_{1} \leq r_{2} \leq r_{3} \leq 1 / 2, \vec{r}=\left(r_{1}, r_{2}, r_{3}\right) \in[0,1]^{3}$. Then

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S_{(r, r, r)}\left[C_{q}^{s}(3,2)\right]=\operatorname{co}\left\{\max (0,6 r-2) S_{2}(1), 2 r S_{2}(1)\right\}
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Is $C_{q}^{s}(3,2)=C_{q c}^{s}(3,2)$ ?

Thanks for your attention!

