

Computing regular subalgebras of simple Lie algebras

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July 8, 2010

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- A $(\mathfrak{g}, \mathfrak{l})$ -module M is of *finite type* if for any fixed irreducible finite-dimensional \mathfrak{l} -module V the Jordan-Hölder multiplicities of V in all finite-dimensional \mathfrak{l} -submodules of M are bounded.

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$$\mathfrak{l} = \mathfrak{h} \oplus \underbrace{\bigoplus_{\alpha \in \Delta(\mathfrak{l})} \mathfrak{g}^{\alpha}}_{\text{subalgebra of interest}} \quad \mathfrak{k} = \mathfrak{h} \oplus \bigoplus_{\substack{\alpha \in \Delta(\mathfrak{l}): \\ -\alpha \in \Delta(\mathfrak{l})}} \mathfrak{g}^{\alpha}; \quad \mathfrak{n} = \bigoplus_{\substack{\alpha \in \Delta(\mathfrak{l}): \\ -\alpha \notin \Delta(\mathfrak{l})}} \mathfrak{g}^{\alpha} .$$

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- Done in "Semisimple Lie subalgebras of simple Lie algebras", Dynkin.

Definition (Penkov)

- (a) *Cone condition.* \mathfrak{l} satisfies the cone condition if $\text{Cone}_{\mathbb{Q}}(\Delta(\mathfrak{n})) \cap \text{Cone}_{\mathbb{Q}}(\text{Sing}_{\mathfrak{b} \cap \mathfrak{k}}(\mathfrak{g}/\mathfrak{l})) = \{0\}$, where $\text{Sing}_{\mathfrak{b} \cap \mathfrak{k}}$ stands for $\mathfrak{b} \cap \mathfrak{k}$ -singular. Motivation: I. Penkov, V. Serganova, G. Zuckerman, 2004.
- (b) *Centralizer condition.* \mathfrak{l} satisfies the centralizer condition if $(C([\mathfrak{k}, \mathfrak{k}] \cap N(\mathfrak{n}))_{ss})$ has simple constituents of type A and C only. Motivation: S. Fernando, 1990.

Theorem

Let $\mathfrak{l} = \mathfrak{k} \oplus \mathfrak{n}$ be a subalgebra containing a Cartan subalgebra of the simple Lie algebra $\mathfrak{g} \simeq \mathfrak{sl}(n), \mathfrak{so}(2n+1), \mathfrak{sp}(2n), \mathfrak{so}(2n), E_6, E_7, F_4$ or G_2 . Then \mathfrak{l} is a Fernando-Kac subalgebra of finite type if and only if the cone condition and the centralizer condition are satisfied.

If the nilradical is zero, the cone condition is trivially satisfied. The theorem follows directly from [Fer90] (non-existence if centralizer condition fails) and the construction in [BL87] (if centralizer condition holds).

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Proposition

Suppose you can find a relation

$$a_1\alpha_1 + \cdots + a_l\alpha_l = b_1\beta_1 + \cdots + b_k\beta_k$$

where $\alpha \in \text{Sing}_{\mathfrak{g} \cap \mathfrak{k}}(\mathfrak{g}/\mathfrak{l})$, $\beta_i \in \Delta(\mathfrak{n})$, $a_i, b_j \in \mathbb{Z}_{>0}$, and in addition

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The failure of the cone condition alone is sufficient for the existence of such a relation in types $A_n, B_n, D_n, E_6, E_7, G_2$.

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- Rant about C++, the “vector partition” program, etc.
http://vector-partition.jacobs-university.de/cgi-bin/vector_partition_linux.cgi?rootSAs

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- Step 3. Check whether Δ' is isomorphic via a root system isomorphism of $\Delta(\mathfrak{g})$ to a simple basis of an element already present in R . If so, terminate the current branch of computation. Otherwise, add Δ to R and go to Step 1.

Note: step 3 - not needed in the sense that one can compare only two pairs of Dynkin diagrams.

Proposition

For the two root subsystems Δ_1 and Δ_2 to be isomorphic via isomorphism of $\Delta(\mathfrak{g})$ it is necessary and sufficient that their Dynkin diagrams and the Dynkin diagrams of Δ_1^\perp and Δ_2^\perp are isomorphic, where \perp stands for strongly orthogonal ($\alpha \perp \beta$ if $\alpha \pm \beta$ is not a root).

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Note: if Δ parametrizes a regular subalgebra $[\mathfrak{k}, \mathfrak{k}]$, then Δ^\perp parametrizes the root system of the centralizer of $[\mathfrak{k}, \mathfrak{k}]$. The centralizer of a regular subalgebra consists of a regular subalgebra and a piece of the Cartan subalgebra.

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- Take an arbitrary α' in $\Delta(\mathfrak{g}) \setminus \Delta(\mathfrak{k})$. Start adding positive roots of \mathfrak{k} to α' until possible. The root obtained in the end is the $\mathfrak{b} \cap \mathfrak{k}$ -singular vector in the simple \mathfrak{k} -submodule containing $g^{\alpha'}$.

Generate all nilradicals up to isomorphism that can be attached to \mathfrak{k} (one representative per isomorphism class)

Let the \mathfrak{k} -module decomposition of $\mathfrak{g}/\mathfrak{k}$ be $M_1 \oplus \cdots \oplus M_N$.

- Pairing table. We say M_i and M_j pair to M_k if there exist $\alpha \in \text{Weights}(M_i)$ and $\beta \in \text{Weights}(M_j)$ such that $\alpha + \beta \in \text{Weights}(M_k)$.

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- Choose an arbitrary order \prec' on the set of all subsets of $\{M_1, \dots, M_N\}$. Using W' and \prec' induce a partial order \prec on all subsets of $\{M_1, \dots, M_N\}$.

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- To enumerate all nilradicals up to isomorphism (getting one representative in each W' -class) one enumerates all subsets of the M_i 's that respect the pairing table, have no opposite modules, and respect \prec .

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- Pairing table. We say M_i and M_j pair to M_k if there exist $\alpha \in \text{Weights}(M_i)$ and $\beta \in \text{Weights}(M_j)$ such that $\alpha + \beta \in \text{Weights}(M_k)$.
- Opposite modules. M_i is opposite to M_j if $\text{Weights}(M_i) = -\text{Weights}(M_j)$.
- Compute the group W' of all root system automorphisms of $\Delta(\mathfrak{g})$ that preserve $\Delta(\mathfrak{b} \cap \mathfrak{k})$. Example: $\Delta(\mathfrak{k}) = 7A_1 \subset E_7$. Then W' has 168 elements.
- Choose an arbitrary order \prec' on the set of all subsets of $\{M_1, \dots, M_N\}$. Using W' and \prec' induce a partial order \prec on all subsets of $\{M_1, \dots, M_N\}$.
- To enumerate all nilradicals up to isomorphism (getting one representative in each W' -class) one enumerates all subsets of the M_i 's that respect the pairing table, have no opposite modules, and respect \prec . Numerology: E_6 : 64580 possibilities.

Generate all $\mathfrak{sl}(2)$ -subalgebras of \mathfrak{g} : starting facts from Dynkin

Some facts from “Semisimple Lie subalgebras of semisimple Lie algebras”.

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- Let our $\mathfrak{sl}(2)$ be given by h, e, f with $[h, e] = 2e$, $[h, f] = -2f$, $[e, f] = h$. h can be assumed to lie in a Cartan s.a. of \mathfrak{g} .

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- This is a quadratic system of m equations, where m equals the sum of $\text{rk}[\mathfrak{k}, \mathfrak{k}]$ and the number of roots of the form $\alpha_i - \alpha_j$. Solve it!
- Find a simple basis of \mathfrak{g} with respect to h to recover the characteristic of h in \mathfrak{g} .

The “vector partition” program

- Project started December 2008.
- Can compute:
 - Everything described in the talk.
 - Algebraic expressions in closed form for the Kostant partition function. Can go up to A_6 , D_4 , B_4 , C_4 .
 - Hyperplane arrangements (needed to describe the combinatorial chambers for the Kostant partition function).
 - Weyl groups, Kazhdan-Lusztig coefficients, structure constants of simple Lie algebras.
 - Simplex algorithm (basic implementation).
 - Has its own large integer/rational number library, classes implementing quasipolynomials over \mathbb{Q} .
 - Uses its own classes for hashing arrays. No external packages!
- Current size of the mathematical part: $\sim 30\,000$ lines of code. Total project size $> 35\,000$ lines of code.
- 560+ commits in the public source code repository.

Thank you for Your attention!



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