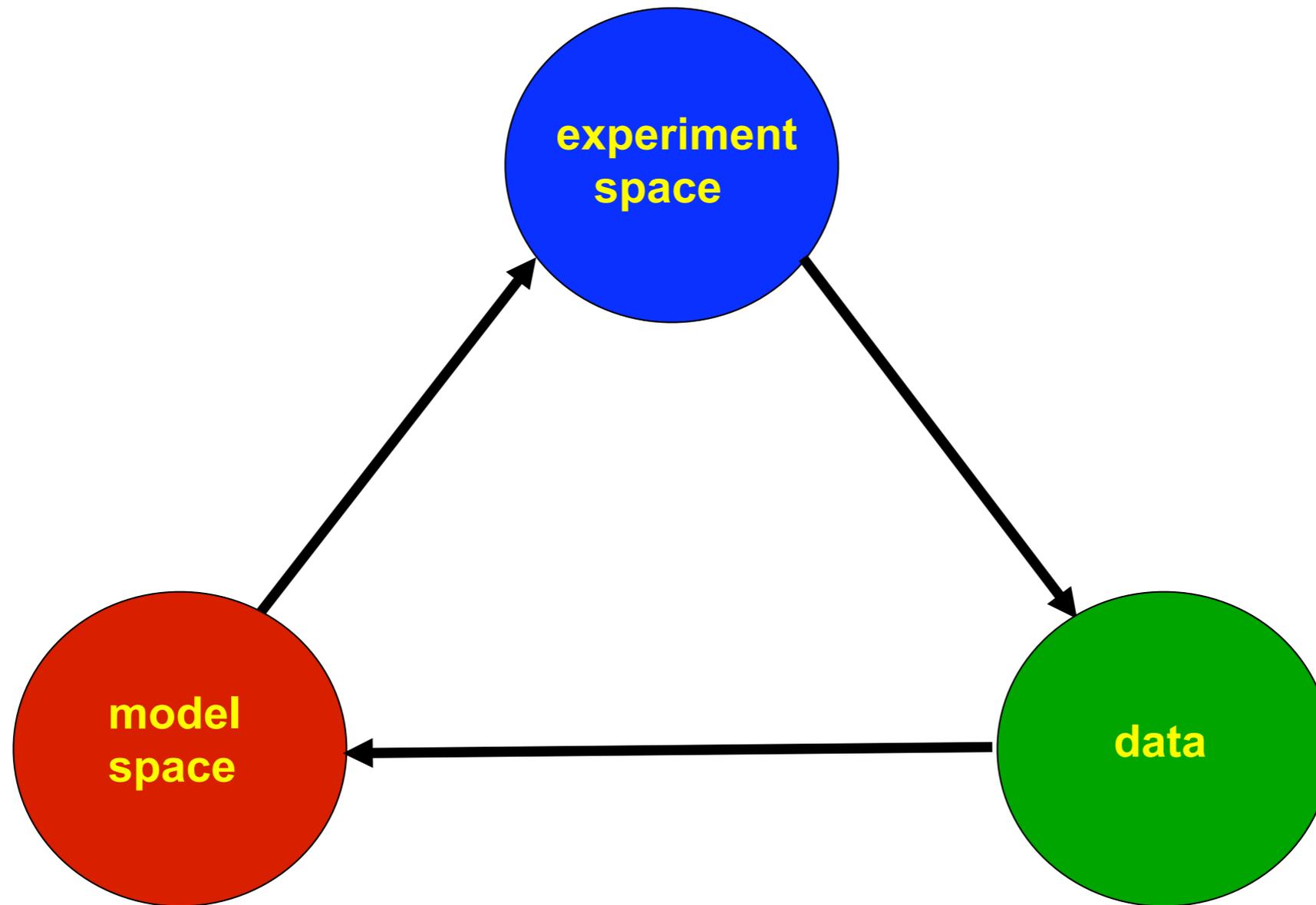


# Adaptive Sensing and Sparse Interactions



BIRS Workshop on Sparse Statistics  
Optimization and Machine Learning  
January 16-21 2011

Rob Nowak

[www.ece.wisc.edu/~nowak](http://www.ece.wisc.edu/~nowak)

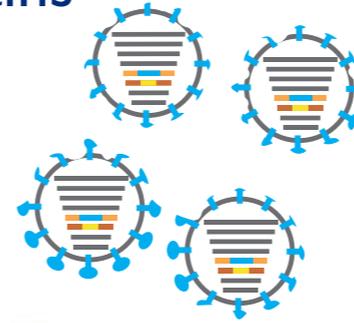
Joint work with R. Castro, J. Haupt,  
M. Malloy and B. Nazer

# Motivation: Inferring Biological Pathways

virus



13,071 single-gene knock-down cell strains



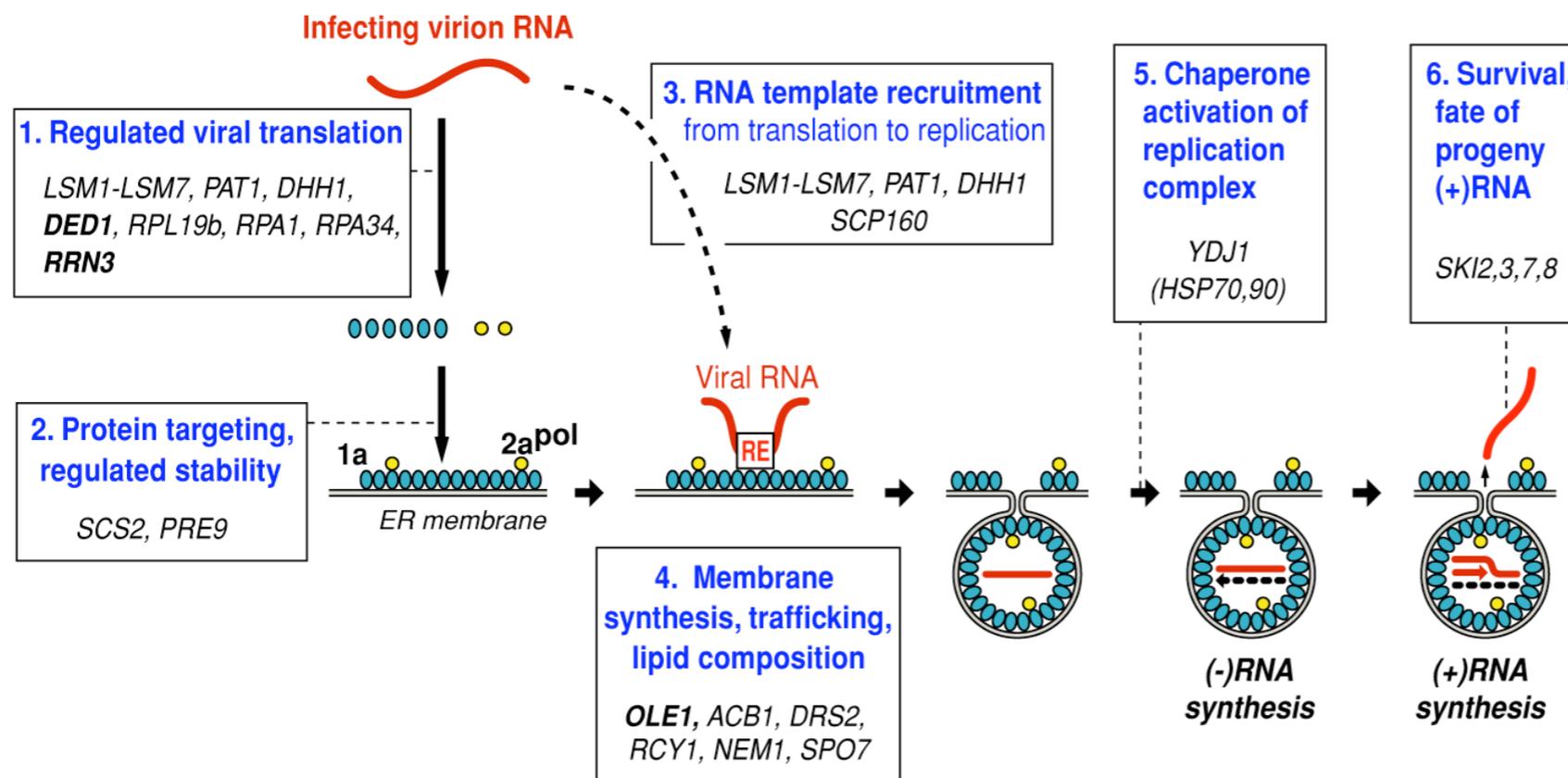
microwell array



infect each strain with fluorescing virus



fruit fly



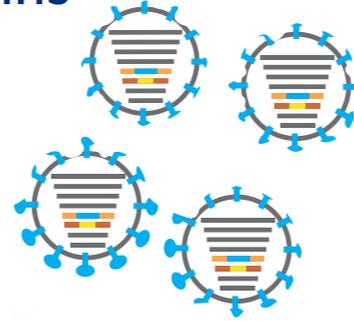
approximately 100 significant genes/proteins discovered by Ahlquist and Kawaoka Labs at UW-Madison

# Motivation: Inferring Biological Pathways

virus



13,071 single-gene  
knock-down cell strains



microwell array



infect each strain  
with fluorescing virus



fruit fly

## Challenges:

1. very low SNR data and huge number of experiments and tests
2. non-linear interactions

# Challenge 1: High-Dimensionality and Low SNR

nature

Vol 454 | 14 August 2008 | doi:10.1038/nature07151

## ***Drosophila* RNAi screen identifies host genes important for influenza virus replication**

Linhui Hao<sup>1,2\*</sup>, Akira Sakurai<sup>3\*†</sup>, Tokiko Watanabe<sup>3</sup>, Ericka Sorensen<sup>1</sup>, Chairul A. Nidom<sup>5,6</sup>, Michael A. Newton<sup>4</sup>, Paul Ahlquist<sup>1,2</sup> & Yoshihiro Kawaoka<sup>3,7,8,9</sup>

How do they confidently determine the ~100 out of 13K genes hijacked for virus replication from extremely noisy data?

### **Sequential Experimental Design:**

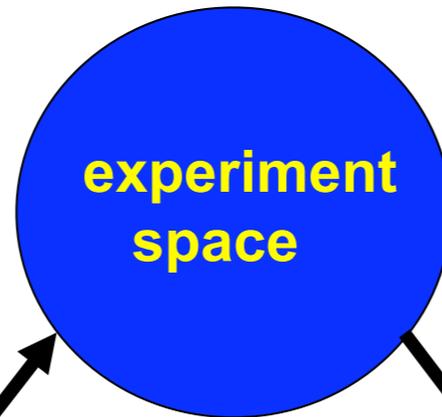
**Stage 1:** assay all 13K strains, **twice**; keep all with significant fluorescence in one or both assays for 2nd stage (13K → 1K)

**Stage 2:** assay remaining 1K strains, **6-12 times**; retain only those with statistically significant fluorescence (1K → 100)

vastly more efficient than replicating all 13K experiments many times

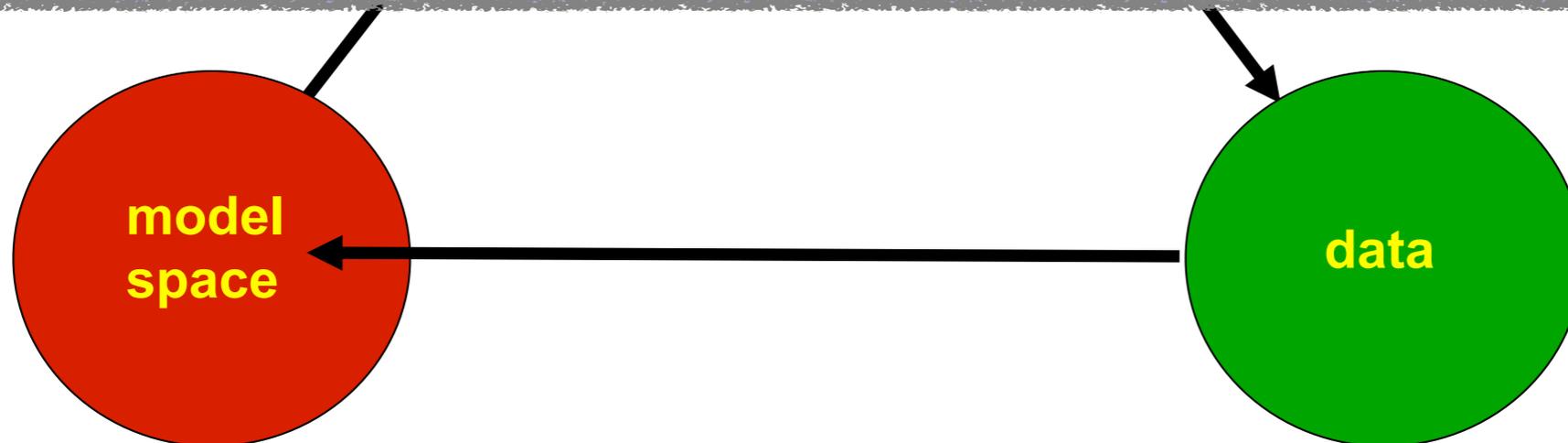
# Feedback from Data Analysis to Data Collection

high-throughput  
experiments



Optimized multi-stage designs controlling the false discovery or the family-wise error rate

S. Zehetmayer, P. Bauer and M. Posch, Statist. Med. 2008; 27:4145–4160



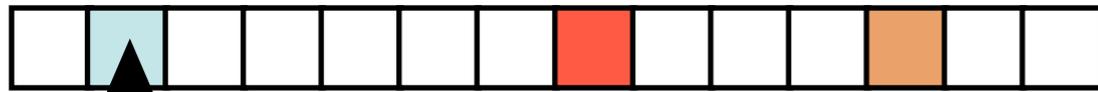
sets of genes critical to a  
certain function/process

microarray or  
assay datasets

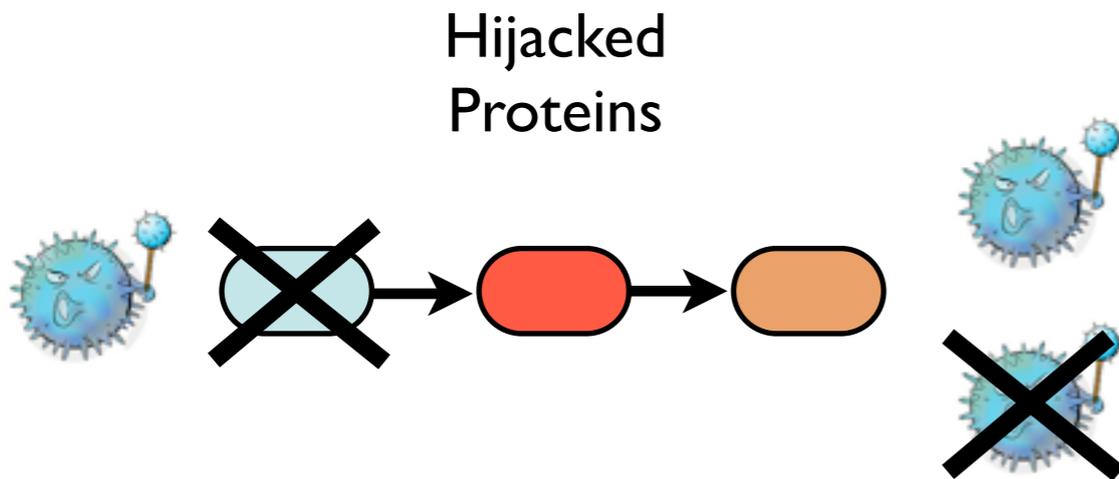
# Challenge 2: Sparsity & Nonlinear Effects

## Linear Pathway

Genome



Knockdown

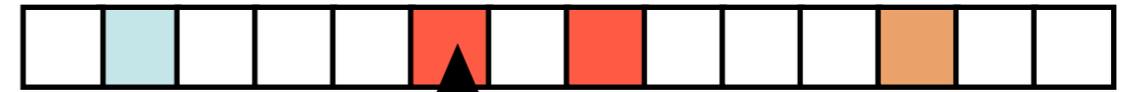


detectable effect

**Linear Sparsity**

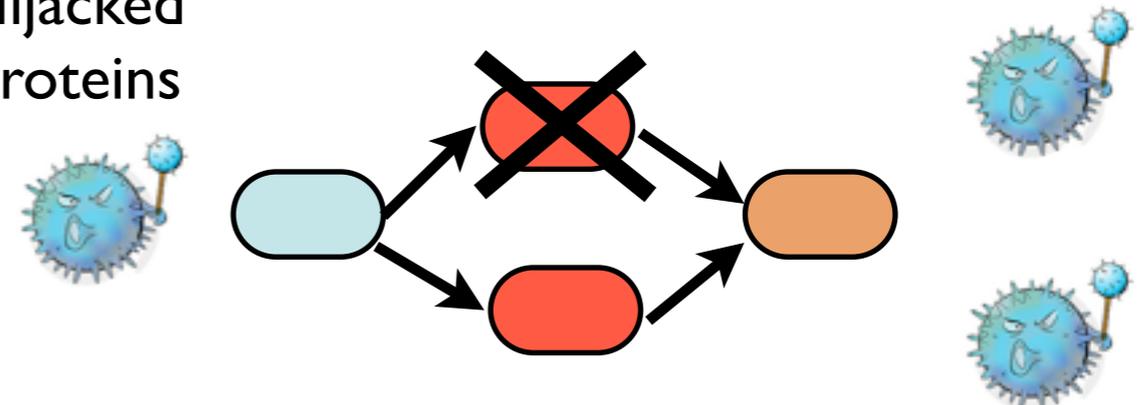
## Redundant Pathways

Genome



Knockdown

Hijacked Proteins



little or no observable effect !

**Multilinear Sparsity**

# Outline of Talk

## I. Sequential Experimental Designs for High-Dimensional Testing

thresholds for recovery in high-dimensional limit:

**non-adaptive designs**      SNR  $\sim \log n$

**sequential designs**      SNR  $\sim$  arbitrarily slowly growing function of  $n$

## 2. Compressed Sensing of Sparse Multilinear Functions

number of compressed sensing measurements for sparse recovery:

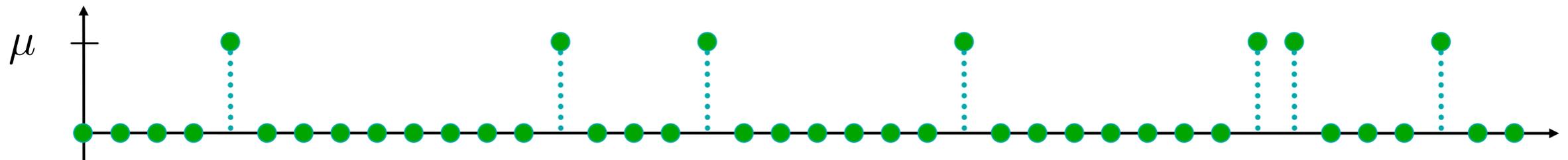
**linear sparsity**       $K \sim S \log n$

**multilinear sparsity**       $K \sim \min\{S^2 \log n, S \log^3(S) \log n, S^\alpha \log^\alpha n\}$

where  $\alpha \geq 1$  *depends* on pattern of sparsity

# Sparse Signal Model

Let  $x = (x_1, \dots, x_n) \in \mathbb{R}^n$  be an unknown sparse vector; most (or all) of its components  $x_i$  are equal to zero.



$$x_i = \begin{cases} \mu > 0, & i \in \mathcal{S} \\ 0, & i \notin \mathcal{S} \end{cases}, \text{ where } |\mathcal{S}| \ll n$$

signal **support set**

deterministic  
but unknown

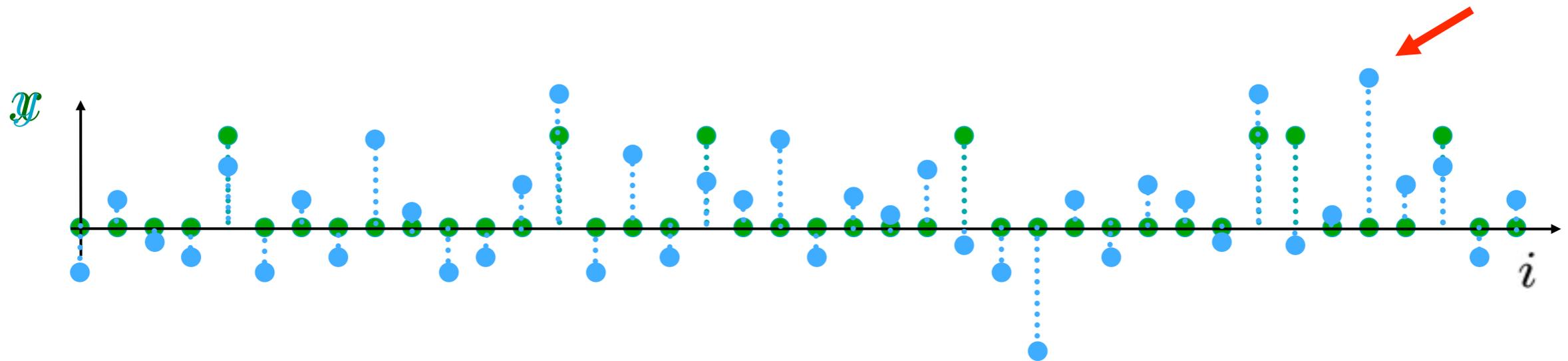
Assume sublinear sparsity level:  $|\mathcal{S}| \ll n$

number of signal  
components

# Noisy Observation Model

$$y_i = x_i + z_i, \quad i = 1, \dots, n$$

$$z_i \stackrel{\text{iid}}{\sim} \mathcal{N}(0, 1)$$



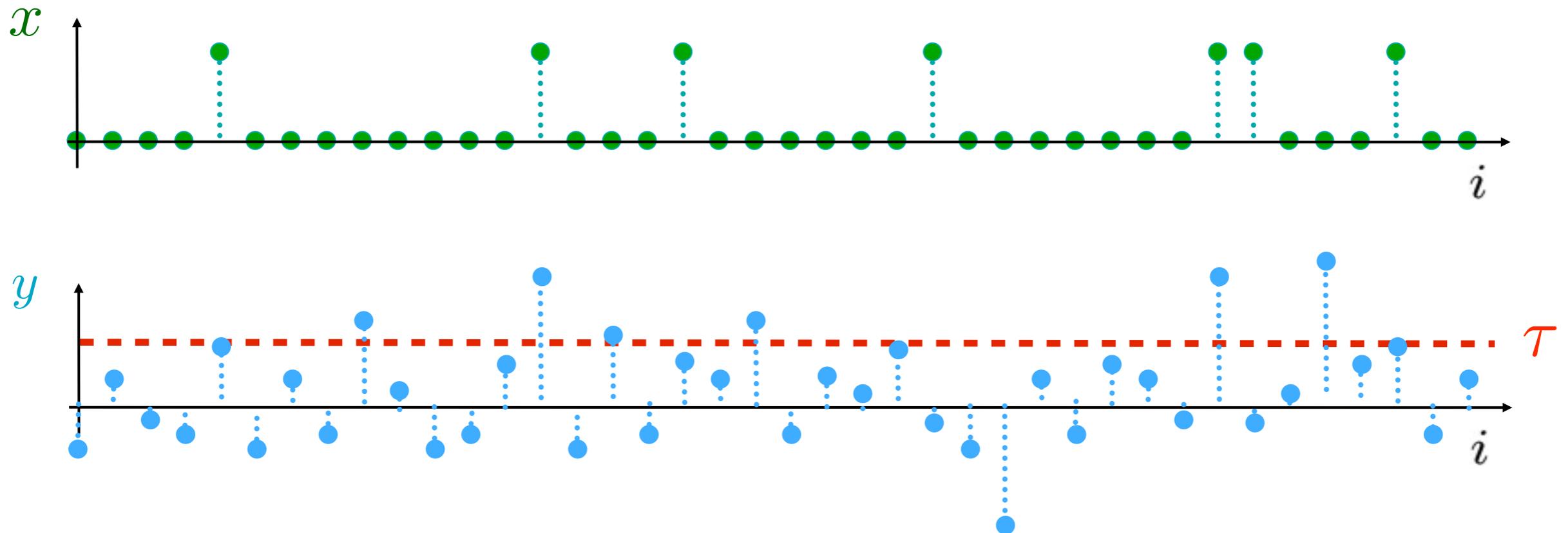
Suppose we want to locate just **one** signal component:  $\hat{i} = \arg \max_i y_i$

Even if no signal is present,  $\max_i y_i \sim \sqrt{2 \log n}$

It is *impossible* to reliably detect signal components weaker than  $O(\sqrt{\log n})$

# Threshold Tests

Our goal is to estimate the set of non-zero components:  $\mathcal{S} := \{i : x_i \neq 0\}$



**Definition 1** A threshold test is an estimator of the form:

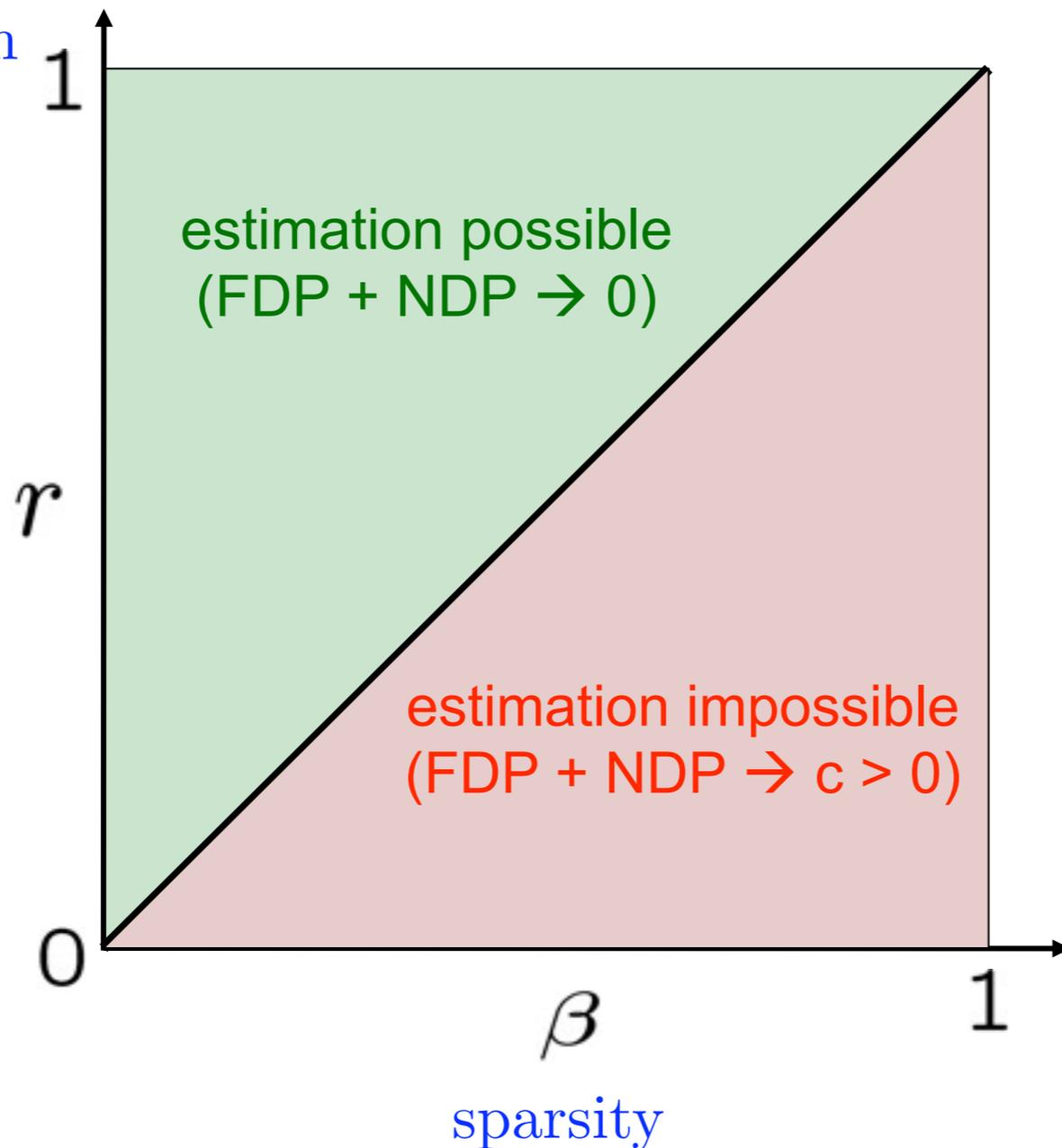
$$\hat{\mathcal{S}}_{\tau}(y) := \{i \in \{1, \dots, n\} : y_i \geq \tau > 0\}$$

**Bonferroni Correction:** To keep the error level small (e.g., less than 5%) the threshold must be on the order of  $\sqrt{\log n}$ .

# False Discovery Rate Control (Ingster '97, Jin & Donoho '03)

Assume sublinear sparsity level:  $|\mathcal{S}| = n^{1-\beta}$ ,  $\beta \in (0, 1)$

signal strength  
 $\mu = \sqrt{2r \log n}$



$$\mathbf{FDP}(\hat{\mathcal{S}}) := \frac{|\hat{\mathcal{S}} \setminus \mathcal{S}|}{|\hat{\mathcal{S}}|}$$

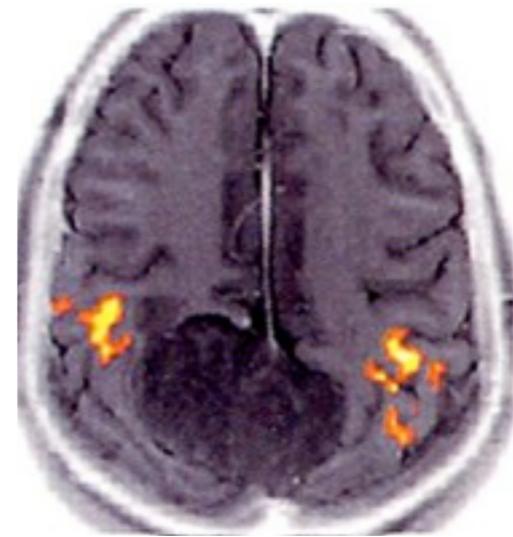
$$= \frac{\# \text{ false discoveries}}{\text{total } \# \text{ discoveries}}$$

$$\mathbf{NDP}(\hat{\mathcal{S}}) := \frac{|\mathcal{S} \setminus \hat{\mathcal{S}}|}{|\mathcal{S}|}$$

$$= \frac{\# \text{ missed non-zeros}}{\# \text{ true non-zeros}}$$

reliable detection iff  $\mu \sim \sqrt{\log n}$  !

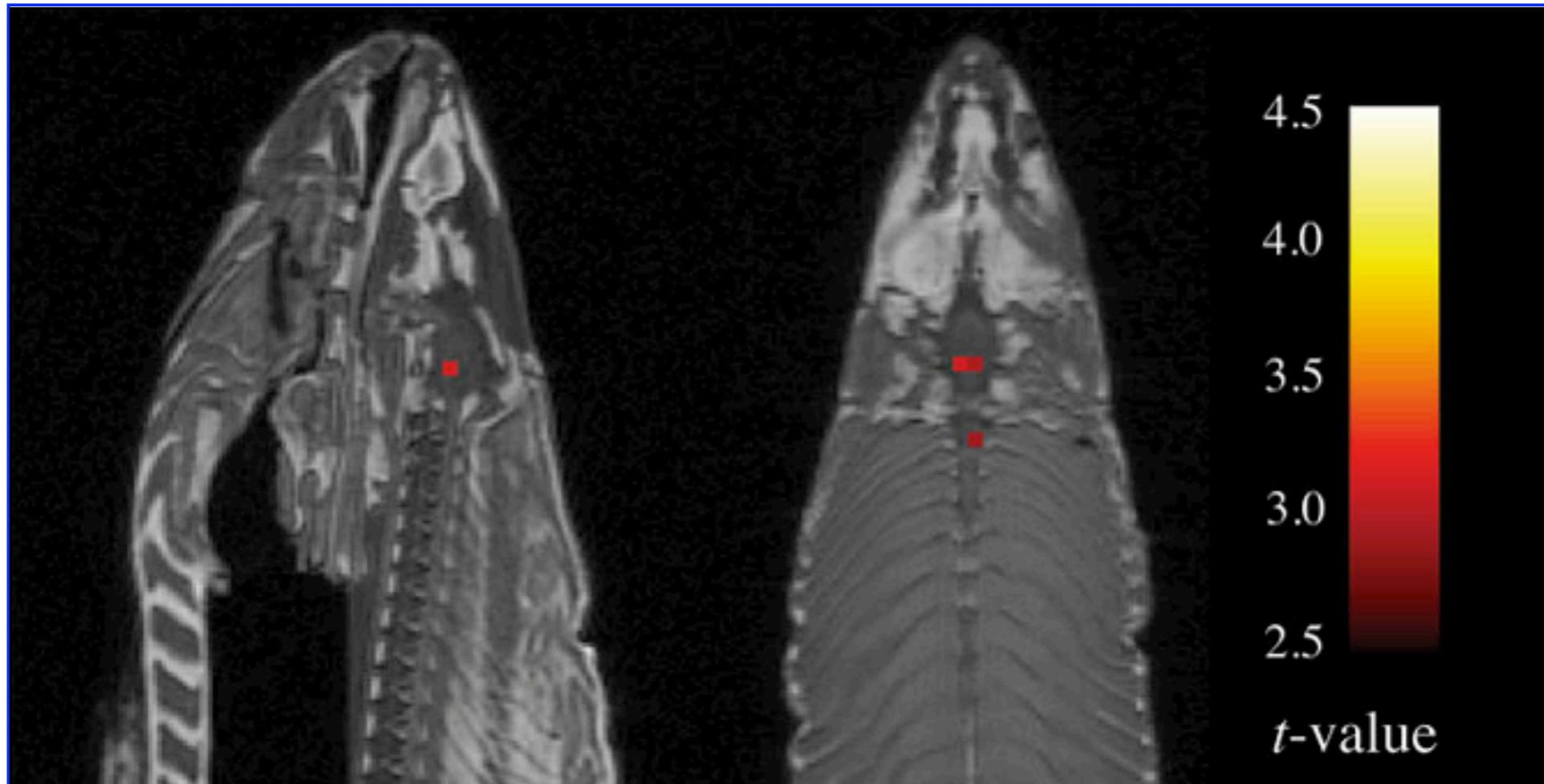
# Is there really a problem ?



[Wired Science](#)  
News for Your Neurons  
[Previous post](#)  
[Next post](#)

## Scanning Dead Salmon in fMRI Machine Highlights Risk of Red Herrings

By [Alexis Madrigal](#)  September 18, 2009 | 5:37 pm | Categories: [Brains and Behavior](#)



# An Alternative: Sequential Experimental Design

Instead of the usual non-adaptive observation model

$$y_i = x_i + z_i, \quad i = 1, \dots, n$$

suppose we are able to sequentially collect several **independent** measurements of each component of  $x$ , according to

$$y_{i,j} = x_i + \gamma_{i,j}^{-1/2} z_{i,j}, \quad i = 1, \dots, n, \quad j = 1, \dots, k$$

where

$j$  indexes the measurement steps

$k$  denotes the total number of steps

$$z_{i,j} \stackrel{\text{iid}}{\sim} \mathcal{N}(0, 1)$$

$\gamma_{i,j} \geq 0$  controls the precision of each measurement

**Total precision budget is constrained, but the choice of  $\gamma_{i,j}$  can depend on past observations  $\{y_{i,\ell}\}_{\ell < j}$ .**

# Experimental (Precision) Budget

sequential measurement model

$$y_{i,j} = x_i + \gamma_{i,j}^{-1/2} z_{i,j}, \quad i = 1, \dots, n, \quad j = 1, \dots, k$$

The precision parameters  $\{\gamma_{i,j}\}$  are required to satisfy

$$\sum_{j=1}^k \sum_{i=1}^n \gamma_{i,j} \leq n$$

For example, the usual non-adaptive, single measurement model corresponds to taking  $k = 1$ , and  $\gamma_{i,1} = 1$ ,  $i = 1, \dots, n$ . This baseline can be compared with adaptive procedures by allowing  $k > 1$  and variable  $\{\gamma_{i,j}\}$  satisfying budget.

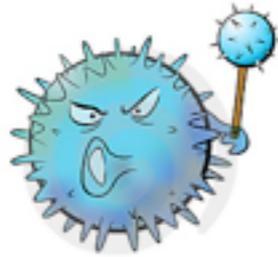
**Precision parameters control the SNR per component.**

SNR is increased/decreased by

- more/fewer repeated samples or
- longer/shorter observation times

# Fruit Fly Example

virus



fruit fly



assay



How to find genes involved in virus replication ?

## Sequential Design Idea

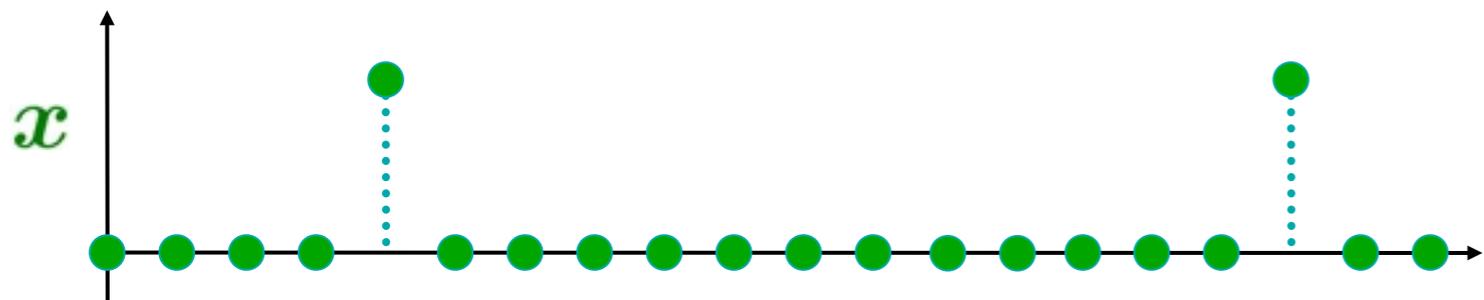
Budget: **k assays**, n tests/assay

**Assay 1**: measure fluorescence of all n genes; discard n/2 genes with weakest fluorescence.

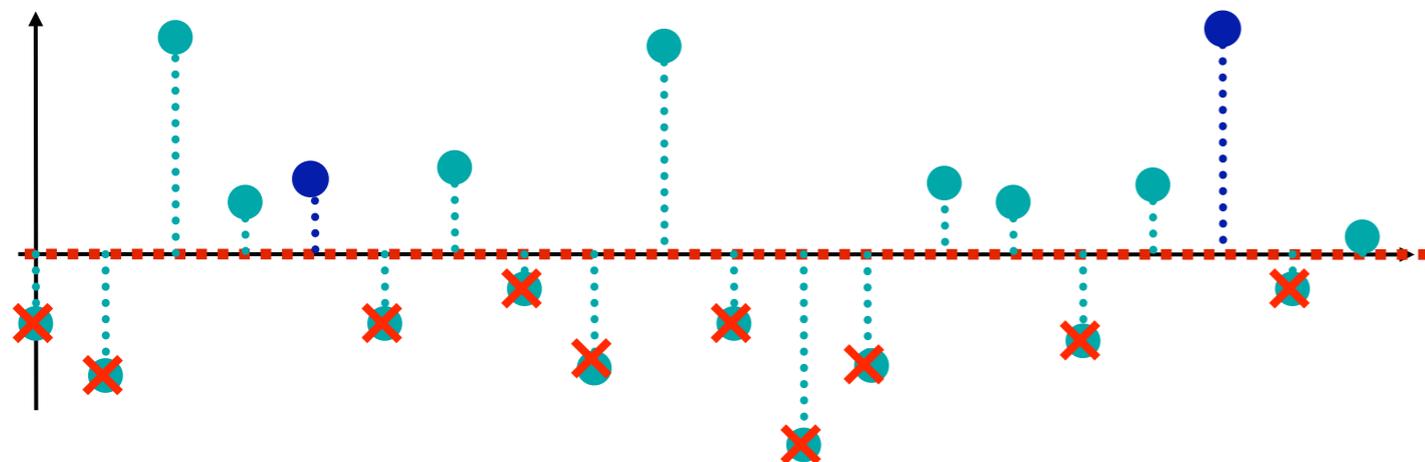
**Assay 2**: measure fluorescence for remaining n/2 genes, each tested twice (double SNR); discard n/4 genes with weakest fluorescence.

**Assay 3**: measure fluorescence for remaining n/4 genes, each tested four times (quadruple SNR); discard n/8 genes with weakest fluorescence.  
continue distilling....

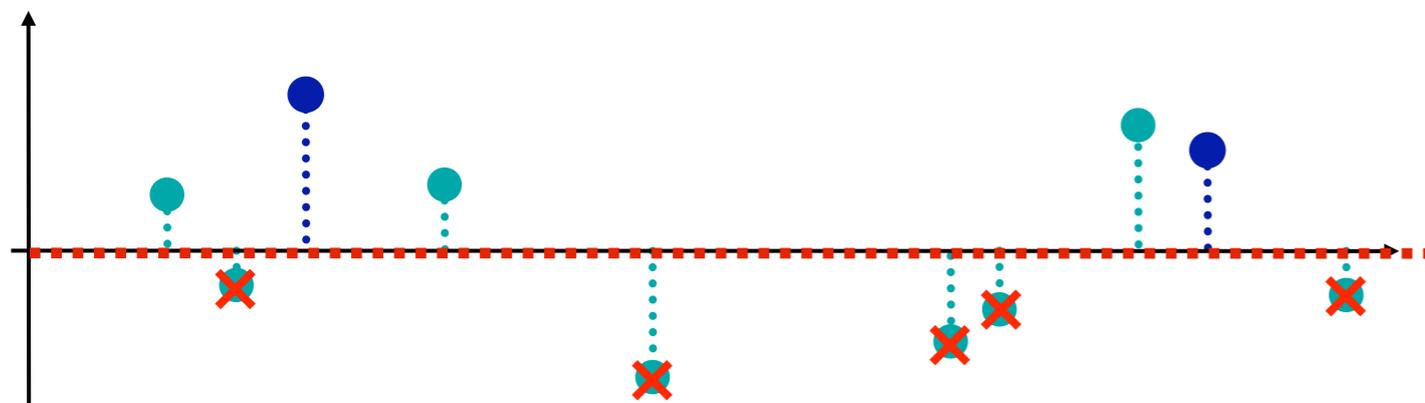
# Idealized Example



Take  $k = 3$  steps and split precision budget uniformly ( $n/3$  per step)



$$y_{i,1} = x_i + \mathcal{N}(0, 3)$$



$$y_{i,2} = x_i + \mathcal{N}\left(0, \frac{3}{2}\right)$$



$$y_{i,3} = x_i + \mathcal{N}\left(0, \frac{3}{4}\right)$$

# Distilled Sensing

## Simple Distilled Sensing

initialize:  $\mathcal{S}_0 = \{1, \dots, n\}$ ,  $\gamma_{i,j}^{-1} = 2 + \epsilon$ ,  $\epsilon > 0$

for  $j = 1, \dots, k$

1) measure:  $y_{i,j} \sim \mathcal{N}(x_i, 2 + \epsilon)$ ,  $i \in \mathcal{S}_{j-1}$

2) threshold:  $\mathcal{S}_j = \{i : y_{i,j} \geq 0\}$

end

output:  $\mathcal{S}_k = \{i : y_{i,k} > 0\}$

total precision budget:  $\mathbb{E} \left[ \sum_{i,j} \gamma_{i,j} \right]$

$$\begin{aligned} &= \frac{1}{2 + \epsilon} \sum_{j=1}^k \mathbb{E} |\mathcal{S}_{j-1}| \\ &\leq \frac{1}{2 + \epsilon} \sum_{j=1}^k \left( \frac{n - |\mathcal{S}|}{2^{j-1}} + |\mathcal{S}| \right) \\ &\leq \frac{2(n - |\mathcal{S}|)}{2 + \epsilon} + k|\mathcal{S}| \leq n \\ &\quad \text{(for } n \text{ large)} \end{aligned}$$

$$\begin{aligned} \text{probability of error: } \mathbb{P}(\mathcal{S}_k \neq \mathcal{S}) &= \mathbb{P}(\{\mathcal{S}^c \cap \mathcal{S}_k \neq \emptyset\} \cup \{\mathcal{S} \cap \mathcal{S}_k^c \neq \emptyset\}) \\ &\leq \mathbb{P}(\mathcal{S}^c \cap \mathcal{S}_k \neq \emptyset) + \mathbb{P}(\mathcal{S} \cap \mathcal{S}_k^c \neq \emptyset) \end{aligned}$$

# False Positives

$$\mathbb{P}(\mathcal{S}_k \neq \mathcal{S}) \leq \mathbb{P}(\mathcal{S}^c \cap \mathcal{S}_k \neq \emptyset) + \mathbb{P}(\mathcal{S} \cap \mathcal{S}_k^c \neq \emptyset)$$

$$\begin{aligned} \mathbb{P}(\mathcal{S}^c \cap \mathcal{S}_k \neq \emptyset) &= \mathbb{P}\left(\bigcup_{i \notin \mathcal{S}} \bigcap_{j=1}^k y_{i,j} > 0\right) \\ &\leq \sum_{i \notin \mathcal{S}} \mathbb{P}\left(\bigcap_{j=1}^k y_{i,j} > 0\right) \\ &= \sum_{i \notin \mathcal{S}} 2^{-k} = \frac{n - s}{2^k} \end{aligned}$$

# False Negatives

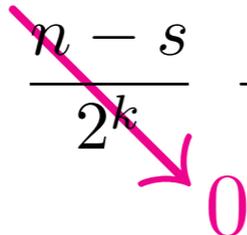
$$\mathbb{P}(\mathcal{S}_k \neq \mathcal{S}) \leq \mathbb{P}(\mathcal{S}^c \cap \mathcal{S}_k \neq \emptyset) + \mathbb{P}(\mathcal{S} \cap \mathcal{S}_k^c \neq \emptyset)$$

$$\begin{aligned} \mathbb{P}(\mathcal{S} \cap \mathcal{S}_k^c \neq \emptyset) &= \mathbb{P}\left(\bigcup_{j=1}^k \bigcup_{i \in \mathcal{S}} y_{i,j} < 0\right) \\ &\leq \frac{k|\mathcal{S}|}{2} \exp\left(-\frac{\mu^2}{2(2+\epsilon)}\right) \end{aligned}$$

# Probability of Error Bound

$$\begin{aligned}\mathbb{P}(\mathcal{S}_k \neq \mathcal{S}) &\leq \mathbb{P}(\mathcal{S}^c \cap \mathcal{S}_k \neq \emptyset) + \mathbb{P}(\mathcal{S} \cap \mathcal{S}_k^c \neq \emptyset) \\ &\leq \frac{n-s}{2^k} + \frac{k|\mathcal{S}|}{2} \exp\left(-\frac{\mu^2}{2(2+\epsilon)}\right) \\ &= \frac{n-s}{2^k} + \frac{1}{2} \exp\left(-\frac{(\mu^2 - 2(2+\epsilon)\log(k|\mathcal{S}|))}{2(2+\epsilon)}\right)\end{aligned}$$

Consider high-dimensional limit as  $n \rightarrow \infty$  and take  $k = \log_2 n^{1+\epsilon}$

$$\mathbb{P}(\mathcal{S}_k \neq \mathcal{S}) \leq \frac{n-s}{2^k} + \frac{1}{2} \exp\left(-\frac{(\mu^2 - 2(2+\epsilon)\log(|\mathcal{S}|(1+\epsilon)\log_2 n))}{2(2+\epsilon)}\right)$$


Second term tends to zero if

$$\mu \geq \sqrt{2(2+\epsilon)\log(|\mathcal{S}|(1+\epsilon)\log_2 n)}$$

# Gains of Sequential Design

non-adaptive threshold:

$$\mu \geq \sqrt{2 \log n}$$

DS threshold:

$$\mu \geq \sqrt{2(2 + \epsilon) \log(|\mathcal{S}|(1 + \epsilon) \log_2 n)}$$

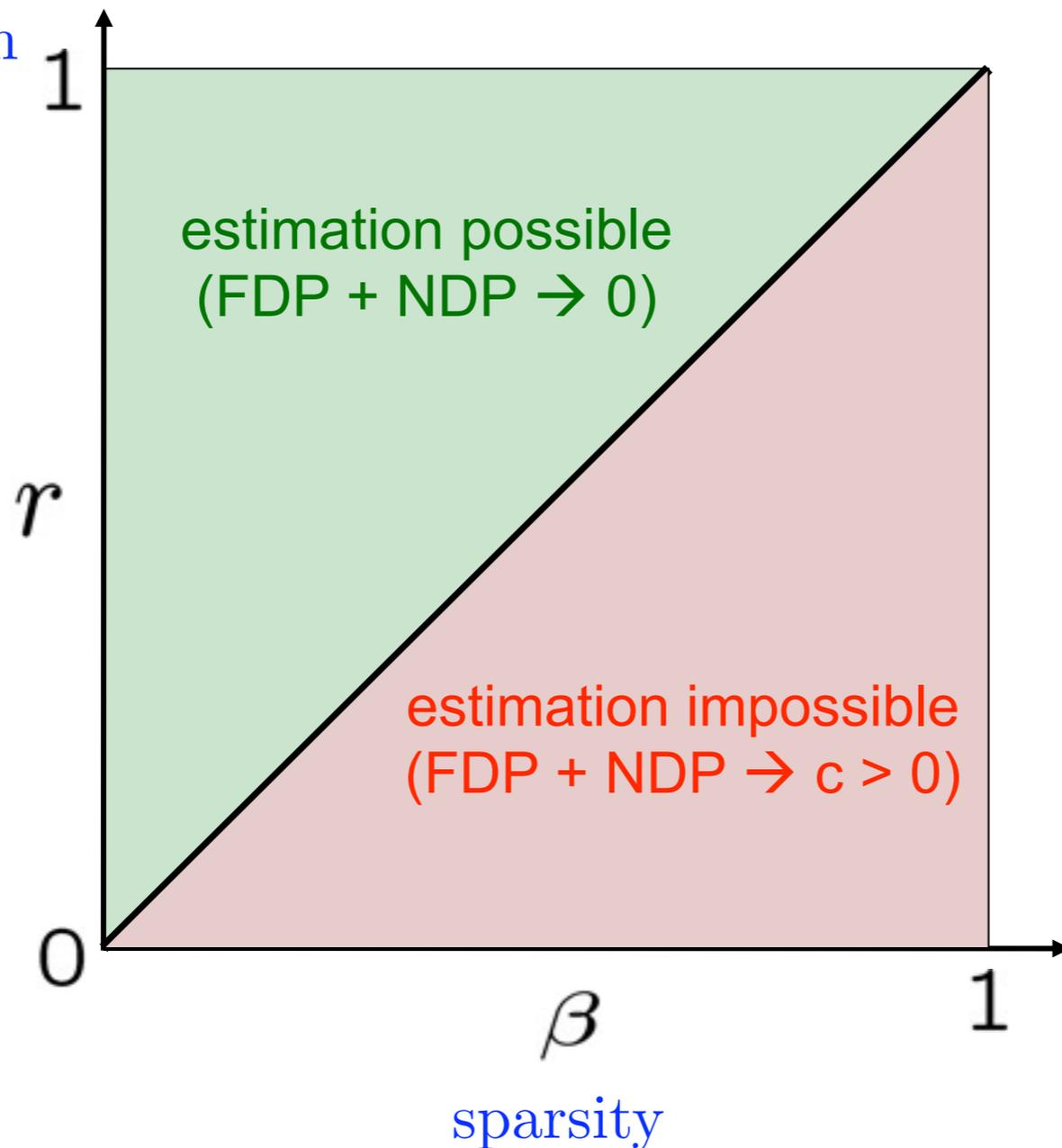
We get a gain whenever  $|\mathcal{S}| \preceq n^{1/2}$

**Punchline:** In ultra-sparse setting, say  $|\mathcal{S}| = C \log n$ , DS drives error to zero if  $\mu \geq \sqrt{(8 + \epsilon) \log \log n}$ , compared to the non-adaptive requirement  $\mu \geq \sqrt{2 \log n}$ .

# False Discovery Rate Control (Ingster '97, Jin & Donoho '03)

Assume sublinear sparsity level:  $|\mathcal{S}| = n^{1-\beta}$ ,  $\beta \in (0, 1)$

signal strength  
 $\mu = \sqrt{2r \log n}$



$$\mathbf{FDP}(\hat{\mathcal{S}}) := \frac{|\hat{\mathcal{S}} \setminus \mathcal{S}|}{|\hat{\mathcal{S}}|}$$

$$= \frac{\# \text{ false discoveries}}{\text{total } \# \text{ discoveries}}$$

$$\mathbf{NDP}(\hat{\mathcal{S}}) := \frac{|\mathcal{S} \setminus \hat{\mathcal{S}}|}{|\mathcal{S}|}$$

$$= \frac{\# \text{ missed non-zeros}}{\# \text{ true non-zeros}}$$

non-sequential methods require  $\mu \sim \sqrt{\log n}$

# FDR-type Control using DS

## FDR Distilled Sensing

**initialize:**  $\mathcal{S}_0 = \{1, \dots, n\}$  ,  $k = \lceil \log \log n \rceil$   
 $\gamma_{i,j} = \left(\frac{3}{4}\right)^j \frac{n}{8} / |\mathcal{S}_{j-1}|$ ,  $j = 1, \dots, k-1$   
 $\gamma_{i,k} = \frac{n}{2|\mathcal{S}_{k-1}|}$

for  $j = 1, \dots, k$

1) **measure:**  $y_{i,j} \sim \mathcal{N}(x_i, \gamma_{i,j}^{-1})$  ,  $i \in \mathcal{S}_{j-1}$

2) **threshold:**  $\mathcal{S}_j = \{i : y_{i,j} \geq 0\}$

end

**output:**  $\mathcal{S}_k = \{i : y_{i,k} \geq 4\}$

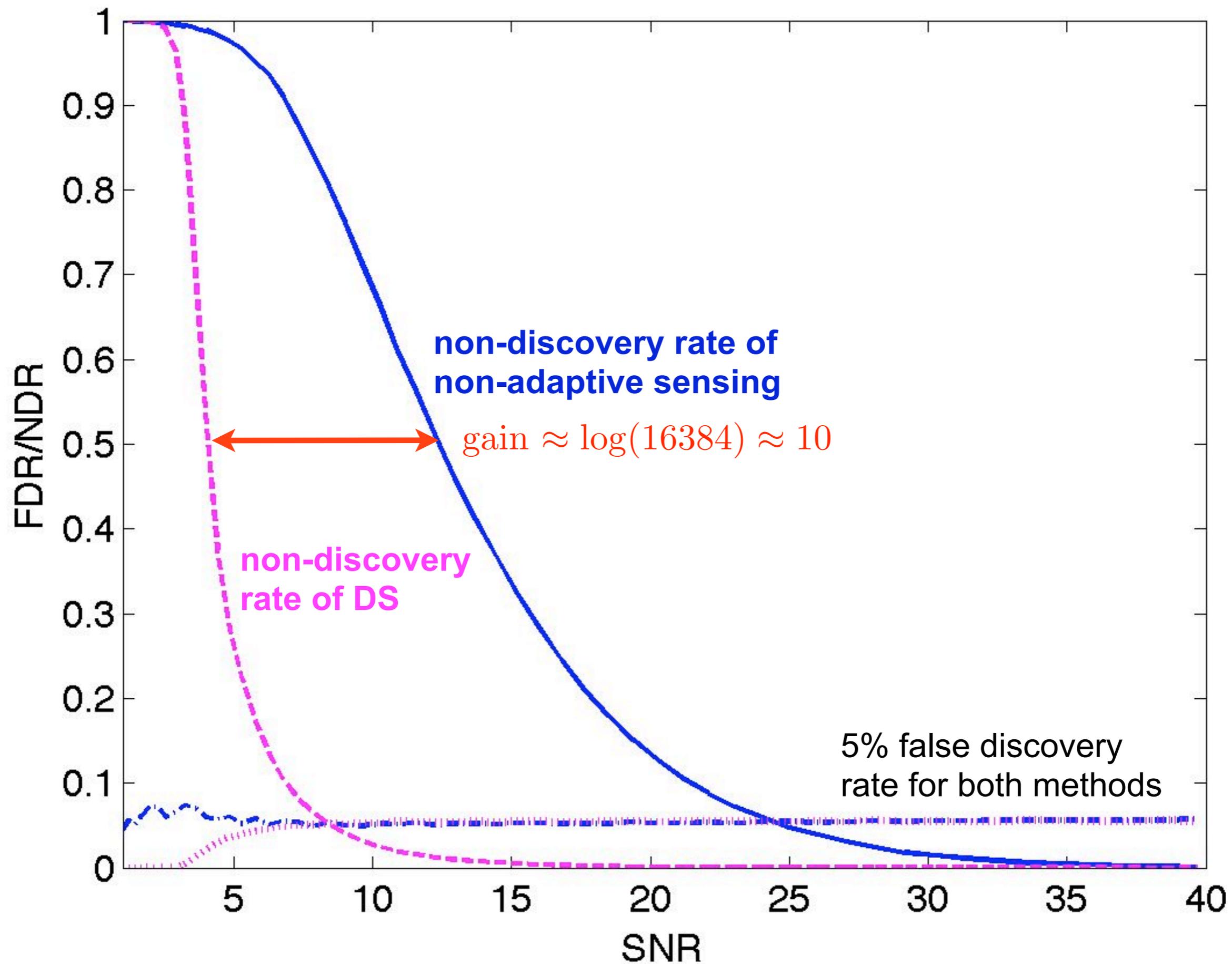
sublinear sparsity:  
 $|\mathcal{S}| = n^{1-\beta}$  ,  $\beta \in (0, 1)$

To guarantee that the proportions of FDP and NDP to zero as  $n \rightarrow \infty$

**Distilled Sensing**  $\mu \sim$  arbitrarily slowly growing function of  $n$   
**non-adaptive**  $\mu \sim \sqrt{\log n}$

**Adaptivity effectively eliminates the fundamental statistical challenge in high-dimensional multiple testing.**

**Example**  $n = 2^{14}$ ,  $\|x\|_0 = \sqrt{n} = 128$



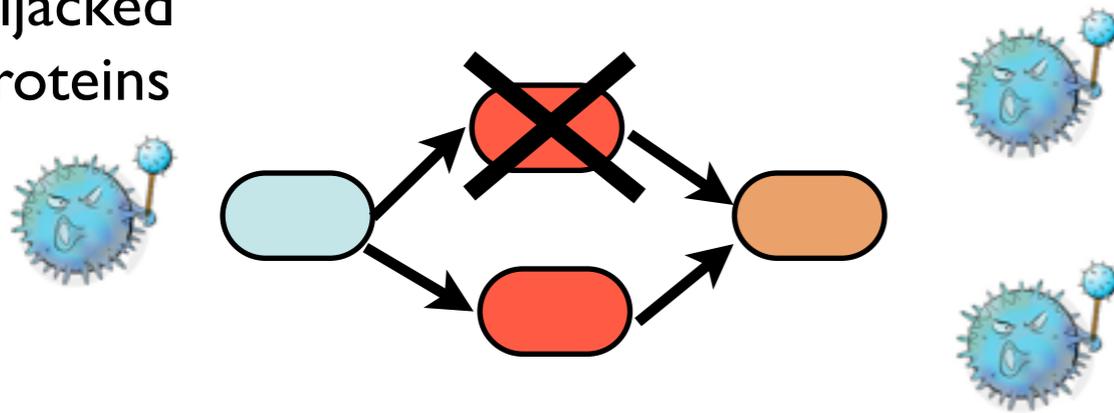
# Challenge 2: Nonlinearities

Genome



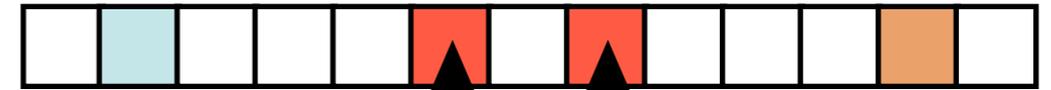
Knockdown

Hijacked Proteins



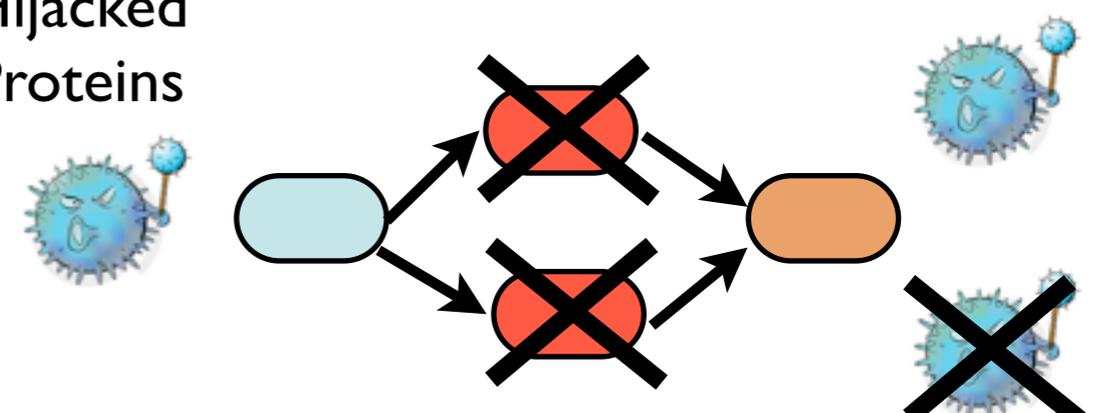
no effect

Genome



Knockdowns

Hijacked Proteins

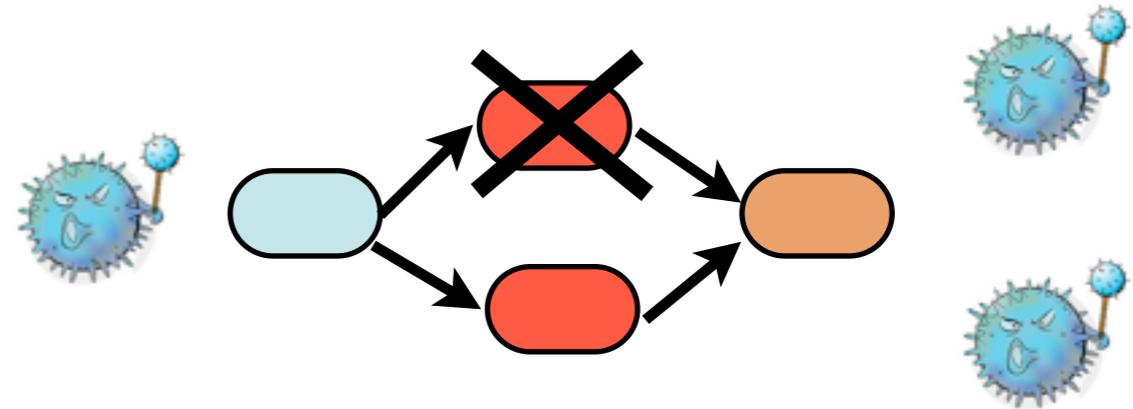
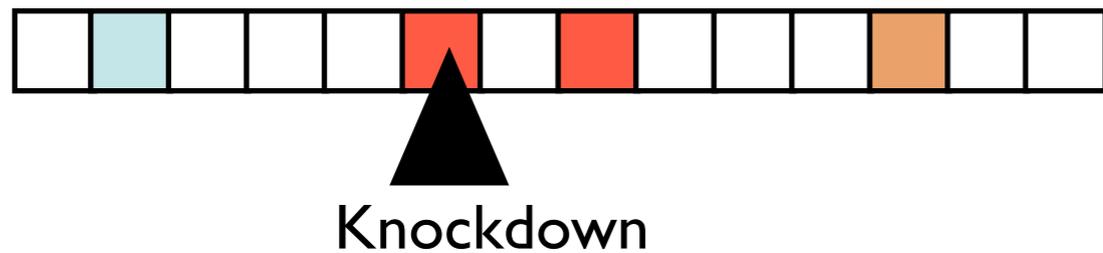


detectable effect

must knockdown both redundant genes to see an effect!

$\binom{13000}{2} \approx 85,000,000$  possible two-fold gene deletion strains !

# Sparse Interaction Models



$$y = \sum_i a_i x_i^{(1)} + \sum_{i < j} a_i a_j x_{ij}^{(2)}$$

**Approximate** output (virus reproduction) with a sparse bilinear system.

$x_i^{(1)}$  non-zero iff gene is critical to pathway

$x_{ij}^{(2)}$  non-zero iff gene pair is critical to pathway

$a_i$  1 if gene is knocked down; 0 otherwise

**sparsity =  
most are 0**

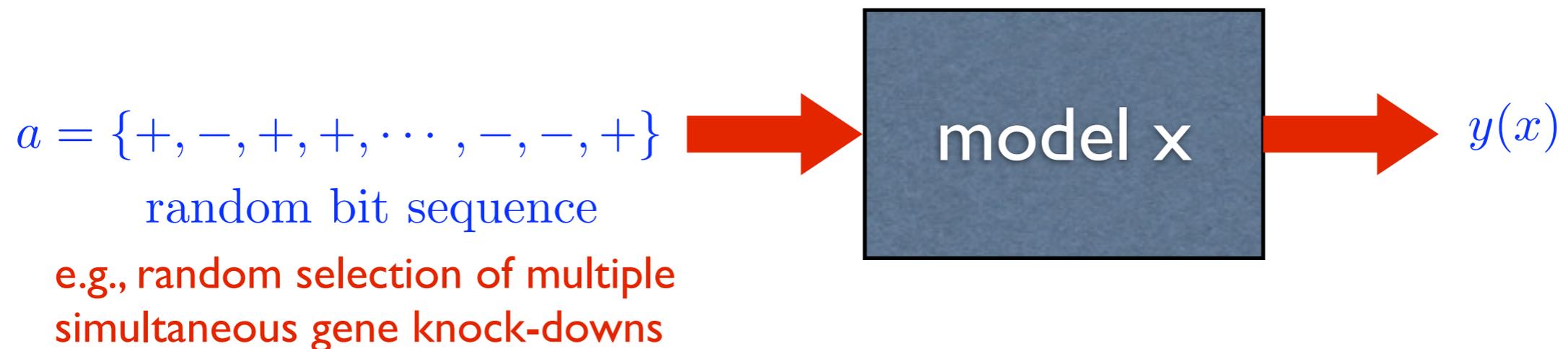
# Sensing Sparse Interactions

**Linear model:**  $y(x) = \sum_i x_i a_i$

**Bilinear model:**  $y(x) = \sum_i x_i^{(1)} a_i + \sum_{i < j} x_{i,j}^{(2)} a_i a_j$

$\binom{13000}{2} \approx 85,000,000$  possible pairwise interactions !

Since most coefficients,  $\{x_i\}$  or  $\{x_i^{(1)}, x_{ij}^{(2)}\}$ , are zero our goal is to identify critical components and interactions from using very few measurements



collect  $K \ll \text{size}(x)$  measurements  $y_1, y_2, \dots, y_K$  using  $K$  random inputs

# (Linear) Compressed Sensing

This is the conventional compressed sensing problem for the linear model.

$$y_k = \sum_i a_{ki} x_i, \quad k = 1, \dots, K$$

$\mathbf{x} \in \mathbb{R}^n, \quad \mathbf{a} \in \{-1, +1\}^n$   
sparse  $\mathbf{x}$ :  $\|\mathbf{x}\|_0 = S \ll n$

find sparse solution to  $\mathbf{y} = \mathbf{A} \mathbf{x}$

If the measurement matrix satisfies the restricted isometry property (RIP) with  $\delta_{2S} < \sqrt{2} - 1$  for all  $S$ -sparse vectors:

$$(1 - \delta_S) \|\mathbf{x}\|_2^2 \leq \|\mathbf{A} \mathbf{x}\|_2^2 \leq (1 + \delta_S) \|\mathbf{x}\|_2^2$$

then  $x$  can be recovered from  $y$  by convex optimization:

$$\min \|\mathbf{z}\|_1 \text{ subject to } \mathbf{A} \mathbf{z} = \mathbf{y}$$

RIP holds with high probability if  $K \geq c S \log(n/S)$

# Multilinear Compressed Sensing

Multilinear model:

$$y = \sum_{i_1 < i_2 < \dots < i_D} a_{i_1} a_{i_2} \dots a_{i_D} x_{i_1 i_2 \dots i_D} \quad \mathbf{x} \in \mathbb{R}^N, \quad \mathbf{a} \in \{-1, 1\}^n$$

$$\|x\|_0 = S \ll N \quad N = \binom{n}{D}$$

$y$  is called a *Rademacher chaos* of order  $D$

Compressed sensing problem for multilinear model

$K$  measurements of this form:  $\mathbf{y} = [y_1 \dots y_K]^T$

find sparse solution to  $\mathbf{y} = \mathbf{A} \mathbf{x}$

matrix  $\mathbf{A}$  now composed of monomials in  $a_{i,j}$

Does it satisfy the RIP property?

# On Average, Things Look Good

$$\text{RIP: } (1 - \delta_S) \|\mathbf{x}\|_2^2 \leq \|\mathbf{Ax}\|_2^2 \leq (1 + \delta_S) \|\mathbf{x}\|_2^2$$

$$\text{isotropic measurements: } \mathbb{E} [\|\mathbf{Ax}\|_2^2] = \|\mathbf{x}\|_2^2$$

$$\text{symmetric binary random inputs: } \mathbb{P}(a_{i,j} = +1) = \mathbb{P}(a_{i,j} = -1) = 1/2$$

Linear CS:  $n = 3$  inputs,  $K = 2$  measurements and  $D = 1$ ,

$$\mathbf{A} = \frac{1}{\sqrt{2}} \begin{bmatrix} a_{1,1} & a_{1,2} & a_{1,3} \\ a_{2,1} & a_{2,2} & a_{2,3} \end{bmatrix} \Rightarrow \mathbb{E}[\mathbf{A}^T \mathbf{A}] = \text{Identity}$$

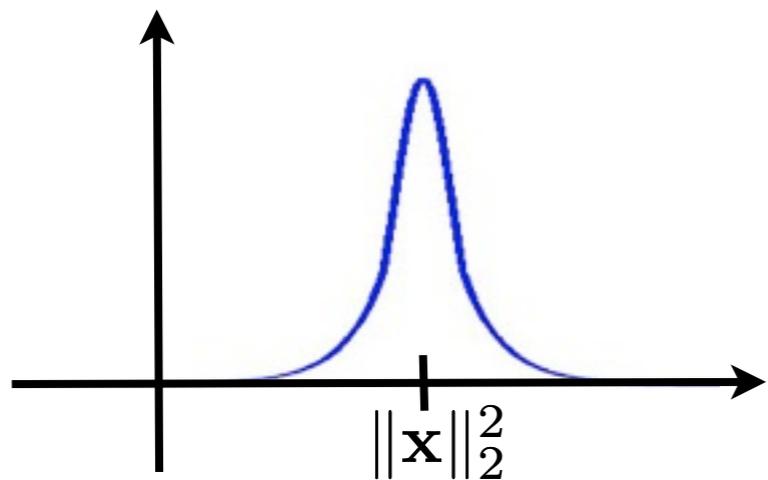
Bilinear CS:  $n = 3$  inputs,  $K = 2$  measurements and  $D = 2$ ,

$$\mathbf{A} = \frac{1}{\sqrt{2}} \begin{bmatrix} a_{1,1}a_{1,2} & a_{1,1}a_{1,3} & a_{1,2}a_{1,3} \\ a_{2,1}a_{2,2} & a_{2,1}a_{2,3} & a_{2,2}a_{2,3} \end{bmatrix} \Rightarrow \mathbb{E}[\mathbf{A}^T \mathbf{A}] = \text{Identity}$$

# What about the distributions?

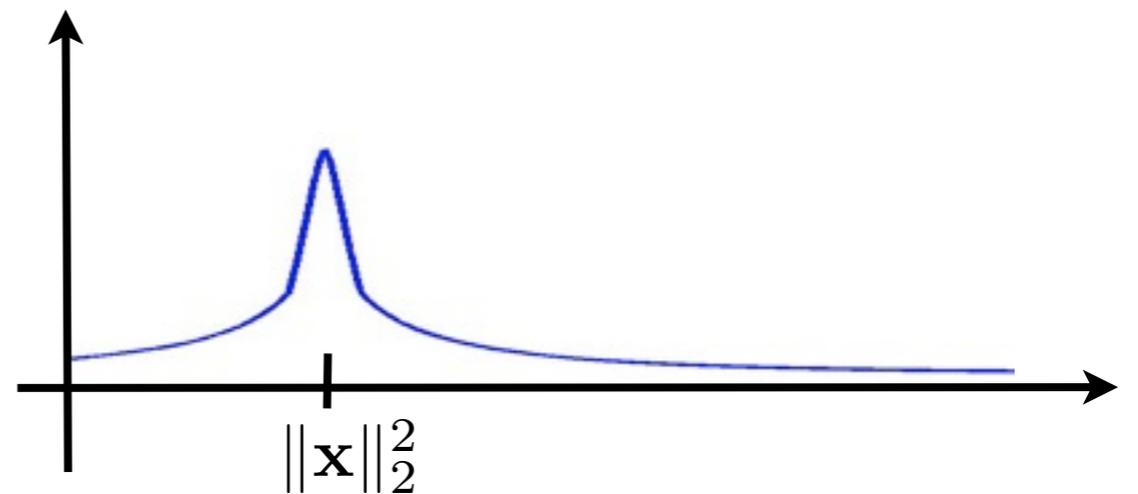
$$\mathbf{y} = \mathbf{A} \mathbf{x}, \quad \mathbb{E} \|\mathbf{y}\|_2^2 = \|\mathbf{x}\|_2^2$$

$$\mathbb{P}(\left| \|\mathbf{y}\|_2^2 - \|\mathbf{x}\|_2^2 \right| > t) \sim \exp(-\text{poly}(t))$$



Gaussian tails  
(as in linear CS)

or



heavy tails?

# Tail Behavior

**Best case:** decoupled chaos

$$\begin{aligned} y &= a_1 a_2 x_{1,2} + a_3 a_4 x_{3,4} + \cdots + a_{2k-1} a_{2k} x_{2k-1,2k} \\ &\equiv \tilde{a}_1 x_{1,2} + \tilde{a}_2 x_{3,4} + \cdots + \tilde{a}_k x_{2k-1,2k} \end{aligned}$$

equivalent to iid binary symmetric sensing

$\Rightarrow$  subgaussian tails independent of  $D$ :  $\mathbb{P}(y^2 - \|x\|_2^2 > t) \leq \exp(-ct)$

**Worst case:** strongly coupled chaos

$$y = \sum_{1 \leq i < j \leq \ell} a_i a_j x_{i,j}, \quad k := \binom{\ell}{2}$$

significant probability of large deviations from mean:

$$\text{if } x_{i,j} = 1/\sqrt{k}, \text{ then } \mathbb{P}(y^2 \geq k) = 2^{-\ell} = 2^{-ck^{1/2}}$$

$\Rightarrow$  heavy tails depending on  $D$ :  $\mathbb{P}(y^2 - \|x\|_2^2 > t) \geq \exp(-ct^{1/D})$

# Combinatorial Dimension of Rademacher Chaos

The **combinatorial dimension**  $1 \leq \alpha \leq D$  measures the level of dependence introduced by a particular pattern of sparsity.

$$y = \sum_{i_1 < i_2 < \dots < i_D} a_{i_1} a_{i_2} \dots a_{i_D} x_{i_1 i_2 \dots i_D}$$

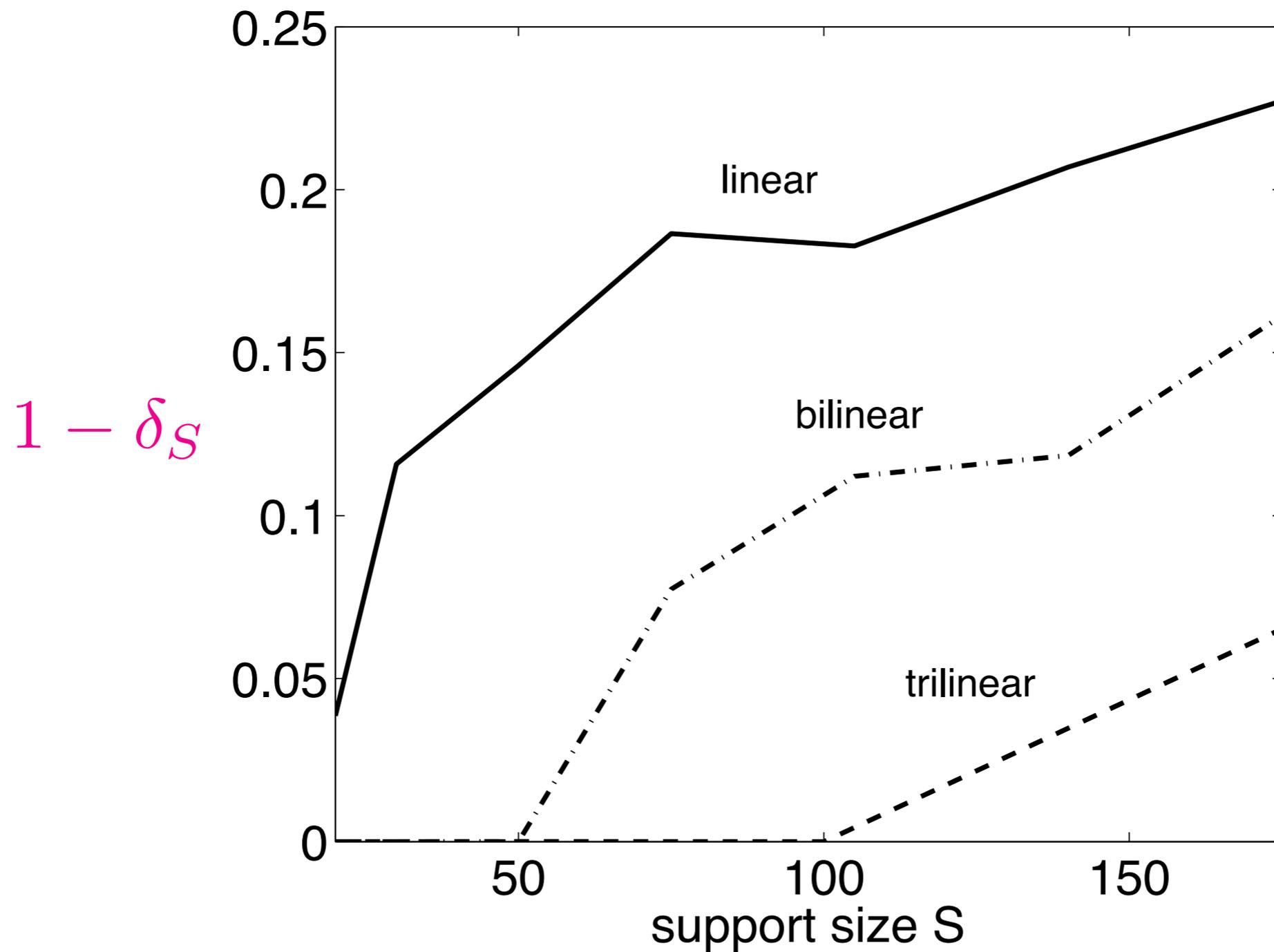
**Blei-Janson '04:** A Rademacher chaos with combinatorial dimension  $\alpha$  satisfies

$$\exp\left(-c_1 t^{1/\alpha}\right) \leq \mathbb{P}\left(|y|^2 > t\right) \leq \exp\left(-c_2 t^{1/\alpha}\right)$$

tails are light to heavy, depending on  $1 \leq \alpha \leq D$

# Dependencies Matter (in practice)

$$\text{RIP: } (1 - \delta_S) \|\mathbf{x}\|_2^2 \leq \|\mathbf{Ax}\|_2^2 \leq (1 + \delta_S) \|\mathbf{x}\|_2^2$$



$1 - \delta_S$

$$K = 5S$$

# Dependencies Matter (in theory)

$K$  = number of measurements needed to recovery  $S$ -sparse multilinear forms

proof technique

bound on  $K$

ingredients

Gershgorin	$S^2 \log N$	empirical 2nd moment bounds, union bound
Rudelson-Vershynin	$S(\log^3 S)(\log N)$	empirical 2nd moment bounds, no union bound
Rademacher chaos	$S^\alpha \log^\alpha(N/S)$ $1 \leq \alpha \leq D$	tail bounds union bound

compare with linear CS bound:  $K \geq S \log(N/S)$

# Gershgorin Bound

i) Control each element of (partial) Gram matrix  $\mathbf{G}_{\mathcal{T}} = \mathbf{A}_{\mathcal{T}}^T \mathbf{A}_{\mathcal{T}}$  using Hoeffding's inequality and bound probability that  $\mathbf{G}_{\mathcal{T}}$  is approximately diagonal.

$$\mathbf{G}_{\mathcal{T}} = \begin{bmatrix} 1 & \frac{\delta_S}{S} & \cdots & \frac{\delta_S}{S} \\ \frac{\delta_S}{S} & 1 & \cdots & \frac{\delta_S}{S} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\delta_S}{S} & \frac{\delta_S}{S} & \cdots & 1 \end{bmatrix}$$

ii) Gershgorin's Disc Theorem guarantees that eigenvalues lie in the range

$$g_{ii} - \sum_{j \neq i} |g_{ij}| \leq \lambda_i(\mathbf{G}_{\mathcal{T}}) \leq g_{ii} + \sum_{j \neq i} |g_{ij}|$$

iii) union bound over all  $\binom{N}{S}$  sparsity patterns. RIP holds if

$$K \geq c S^2 \log N$$

# Heavy-Tailed Restricted Isometries

**Theorem 1 (Vershynin)** Let  $\tilde{\mathbf{A}}$  be a  $K \times N$  measurement matrix whose rows  $\mathbf{a}_i^T$  are independent isotropic random vectors in  $\mathbb{R}^N$ . Let  $B$  be a number such that all entries  $|a_{ij}| \leq B$  almost surely. Then the normalized matrix  $\mathbf{A} = \frac{1}{\sqrt{K}} \tilde{\mathbf{A}}$  satisfies the following for  $K \leq N$ , for every sparsity level  $S \leq N$  and  $0 < \epsilon < 1$ :

if the number of measurements satisfies

$$K \geq C \epsilon^{-2} S \log N \log^3(S)$$

then the RIP constant  $\delta_S$  of  $\mathbf{A}$  satisfies  $\mathbb{E}[\delta_S] \leq \epsilon$ .

Check conditions:

- isotropy:  $\mathbf{G} := \mathbf{A}^T \mathbf{A}$  ,  $\mathbb{E}[\mathbf{G}] = \text{Identity}$
- elements of  $\tilde{\mathbf{A}}$  bounded by 1.

# Chaos Tail Bound

**Lemma 1** Assume that  $y_k$ ,  $k = 1, \dots, K$ , are i.i.d. Rademacher variables of order  $D$  with combinatorial dimension  $1 \leq \alpha \leq D$  and  $\mathbb{E}y_k^2 = 1$ . There exist constants  $c, C > 0$  such that

$$\mathbb{P} \left( \left| \frac{1}{K} \sum_{k=1}^K y_k^2 - 1 \right| > t \right) \leq C \exp(-c \min(Kt^2, K^{1/\alpha} t^{1/\alpha}))$$

proof technique:

- Blei-Jansen chaos tail bounds
- moment bound for sums of symmetric i.i.d. variables due to R. Latała
- apply lemma and union bound over  $\epsilon$ -net for sparse vectors (technique from Baraniuk-Devore-Davenport-Wakin '08)

RIP holds if  $K \geq C S^\alpha \log^\alpha(N/S)$

# Conclusions

## I. Sequential Experimental Designs for High-Dimensional Testing

thresholds for recovery in high-dimensional limit:

**non-adaptive designs**       $\text{SNR} \sim \log n$

**sequential designs**       $\text{SNR} \sim$  arbitrarily slowly growing function of  $n$

Distilled Sensing: Adaptive Sampling for Sparse Detection and Estimation  
J. Haupt, R. Castro, and RN, arXiv:1001.5311v2

## 2. Compressed Sensing of Sparse Multilinear Functions

number of compressed sensing measurements for sparse recovery:

**linear sparsity**       $K \sim S \log n$

**multilinear sparsity**       $K \sim \min\{S^2 \log n, S \log^3(S) \log n, S^\alpha \log^\alpha n\}$

where  $\alpha \geq 1$  *depends* on pattern of sparsity

Sparse Interactions: Identifying High-Dimensional. Multilinear  
Systems via Compressed Sensing, B. Nazer and RN, Allerton 2010