

COLORING JORDAN REGIONS AND CURVES

Louis Esperet

CNRS, Laboratoire G-SCOP, Grenoble, France

New trends in Graph Coloring, Banff
October 19, 2016

PACKING CYCLES IN PLANAR DIGRAPHS

Let G be a directed graph.

$\nu(G)$ = the maximum number of **vertex-disjoint** (directed) cycles in G .

PACKING CYCLES IN PLANAR DIGRAPHS

Let G be a directed graph.

$\nu(G)$ = the maximum number of **vertex-disjoint** (directed) cycles in G .

$\nu^*(G)$ is the solution of the following linear program:

$$\begin{aligned} & \text{Maximize} && \sum_{C \in \mathcal{C}} x_C \\ & \text{s.t.} && \forall v \in V : \sum_{C: v \in C} x_C \leq 1 \\ & && \forall C \in \mathcal{C} : x_C \geq 0 \end{aligned}$$

PACKING CYCLES IN PLANAR DIGRAPHS

Let G be a directed graph.

$\nu(G)$ = the maximum number of **vertex-disjoint** (directed) cycles in G .

$\nu^*(G)$ is the solution of the following linear program:

$$\begin{aligned} & \text{Maximize} && \sum_{C \in \mathcal{C}} x_C \\ & \text{s.t.} && \forall v \in V : \sum_{C: v \in C} x_C \leq 1 \\ & && \forall C \in \mathcal{C} : x_C \geq 0 \end{aligned}$$

For every digraph G , $\nu(G) \leq \nu^*(G)$.

PACKING CYCLES IN PLANAR DIGRAPHS

Let G be a directed graph.

$\nu(G)$ = the maximum number of **vertex-disjoint** (directed) cycles in G .

$\nu^*(G)$ is the solution of the following linear program:

$$\begin{aligned} & \text{Maximize} && \sum_{C \in \mathcal{C}} x_C \\ & \text{s.t.} && \forall v \in V : \sum_{C: v \in C} x_C \leq 1 \\ & && \forall C \in \mathcal{C} : x_C \geq 0 \end{aligned}$$

For every digraph G , $\nu(G) \leq \nu^*(G)$.

Theorem (Reed & Shepherd 1996)

For every **planar** digraph G , $\nu^*(G) \leq 28 \nu(G)$.

PACKING CYCLES IN PLANAR DIGRAPHS

We can assume that $\nu^*(G) = \frac{n}{k}$, and G contains n cycles (repetitions allowed) such that every vertex is in at most k cycles.

PACKING CYCLES IN PLANAR DIGRAPHS

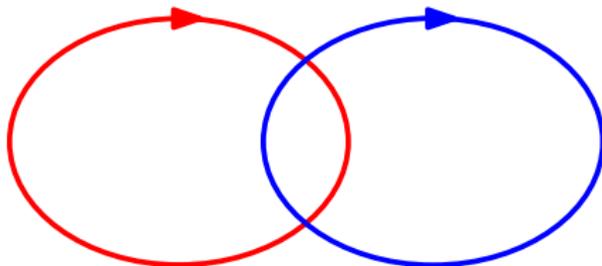
We can assume that $\nu^*(G) = \frac{n}{k}$, and G contains n cycles (repetitions allowed) such that every vertex is in at most k cycles.

We can also assume that the cycles are **pairwise non-crossing**.

PACKING CYCLES IN PLANAR DIGRAPHS

We can assume that $\nu^*(G) = \frac{n}{k}$, and G contains n cycles (repetitions allowed) such that every vertex is in at most k cycles.

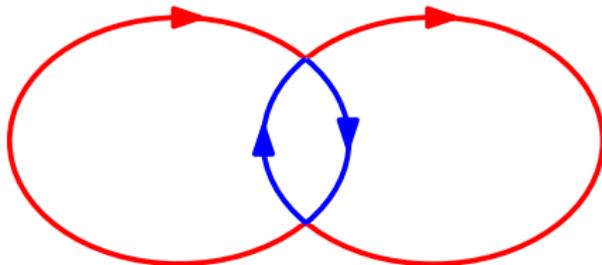
We can also assume that the cycles are **pairwise non-crossing**.



PACKING CYCLES IN PLANAR DIGRAPHS

We can assume that $\nu^*(G) = \frac{n}{k}$, and G contains n cycles (repetitions allowed) such that every vertex is in at most k cycles.

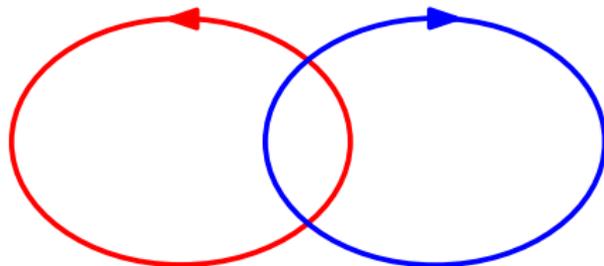
We can also assume that the cycles are **pairwise non-crossing**.



PACKING CYCLES IN PLANAR DIGRAPHS

We can assume that $\nu^*(G) = \frac{n}{k}$, and G contains n cycles (repetitions allowed) such that every vertex is in at most k cycles.

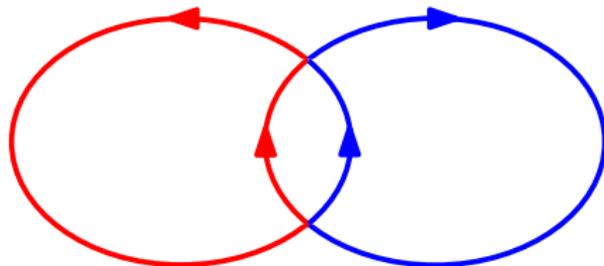
We can also assume that the cycles are **pairwise non-crossing**.



PACKING CYCLES IN PLANAR DIGRAPHS

We can assume that $\nu^*(G) = \frac{n}{k}$, and G contains n cycles (repetitions allowed) such that every vertex is in at most k cycles.

We can also assume that the cycles are **pairwise non-crossing**.



PACKING CYCLES IN PLANAR DIGRAPHS

We can assume that $\nu^*(G) = \frac{n}{k}$, and G contains n cycles (repetitions allowed) such that every vertex is in at most k cycles.

We can also assume that the cycles are **pairwise non-crossing**.

Theorem (Fox & Pach 2012)

Any collection of pairwise non-crossing curves in the plane, such that each point is in at most k curves, can be **properly colored with at most $6e k \approx 16.3 k$ colors**.

PACKING CYCLES IN PLANAR DIGRAPHS

We can assume that $\nu^*(G) = \frac{n}{k}$, and G contains n cycles (repetitions allowed) such that every vertex is in at most k cycles.

We can also assume that the cycles are **pairwise non-crossing**.

Theorem (Fox & Pach 2012)

Any collection of pairwise non-crossing curves in the plane, such that each point is in at most k curves, can be **properly colored with at most $6e k \approx 16.3 k$ colors**.

In particular, if there are n curves, then there are at least $\frac{n}{6ek}$ pairwise disjoint curves. So $\nu(G) \geq \frac{n}{6ek} = \frac{\nu^*(G)}{6e}$, and

$$\nu^*(G) \leq 6e \nu(G) \approx 16.3 \nu(G)$$

COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

To prove that these curves can be properly colored with $6k + 1$ colors, it is enough to show that their intersection graph is $6k$ -degenerate.

COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

To prove that these curves can be properly colored with $6k + 1$ colors, it is enough to show that their intersection graph is $6k$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3kn$.

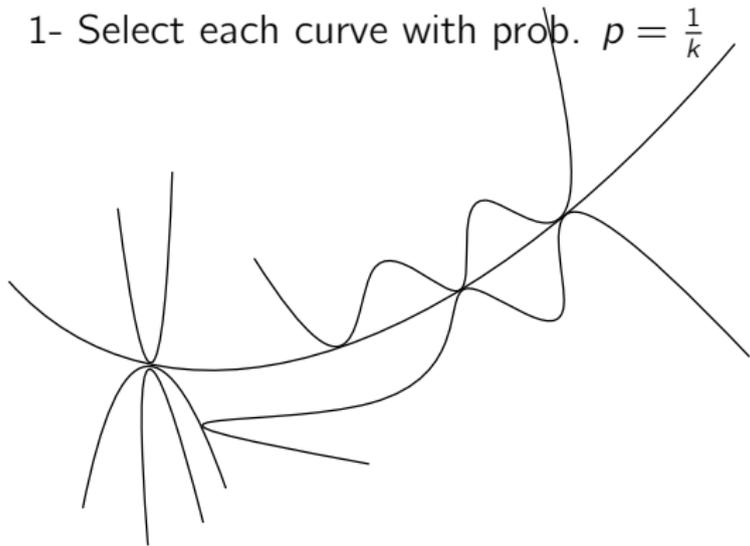
COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ek n$.

1- Select each curve with prob. $p = \frac{1}{k}$



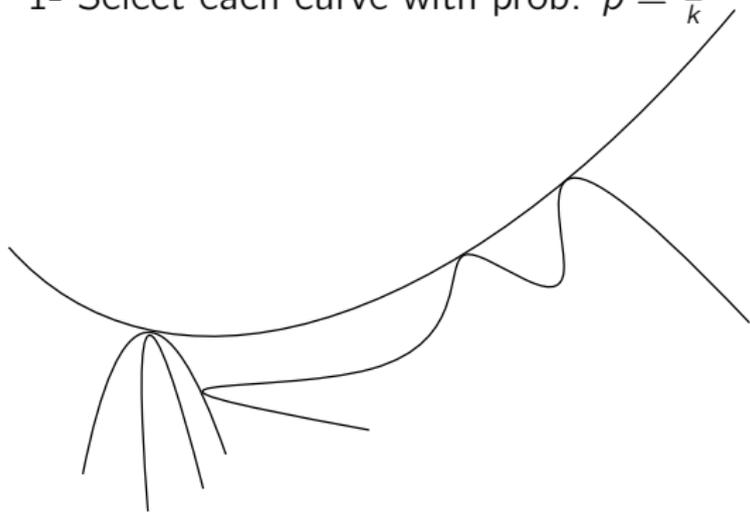
COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ek n$.

1- Select each curve with prob. $p = \frac{1}{k}$



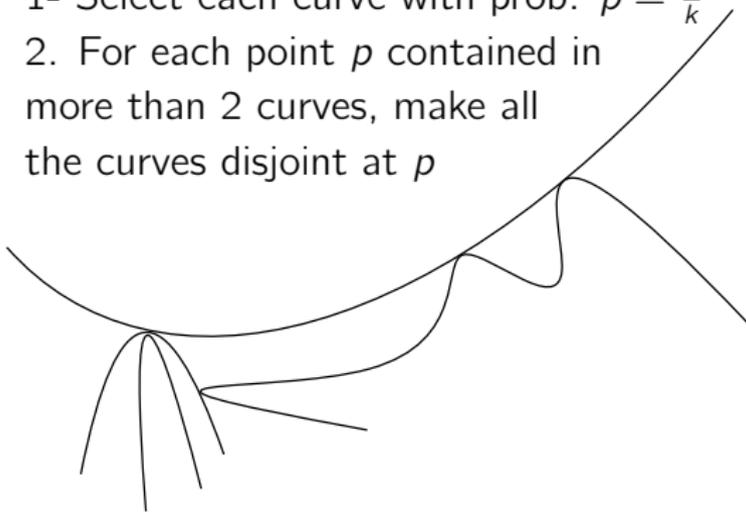
COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ekn$.

- 1- Select each curve with prob. $p = \frac{1}{k}$
2. For each point p contained in more than 2 curves, make all the curves disjoint at p



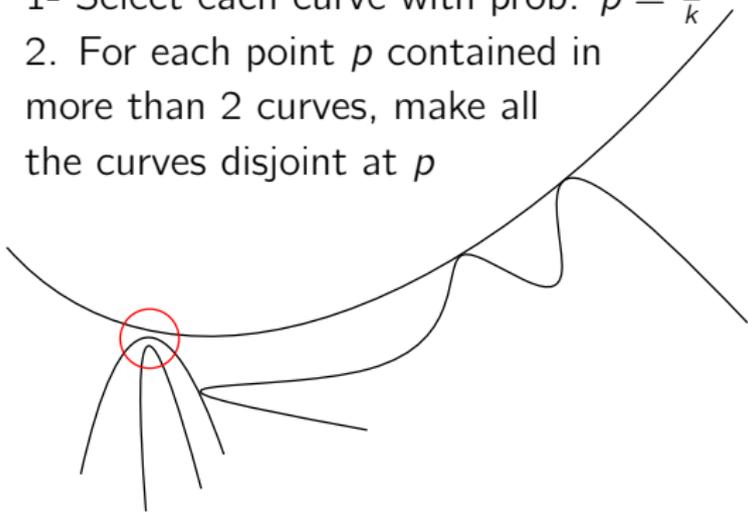
COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ek n$.

- 1- Select each curve with prob. $p = \frac{1}{k}$
2. For each point p contained in more than 2 curves, make all the curves disjoint at p



COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

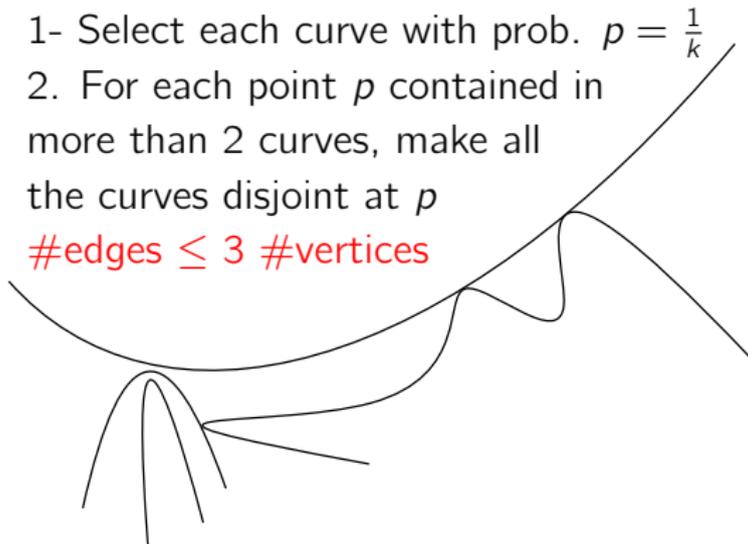
To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ek n$.

1- Select each curve with prob. $p = \frac{1}{k}$

2. For each point p contained in more than 2 curves, make all the curves disjoint at p

$\#edges \leq 3 \#vertices$



COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

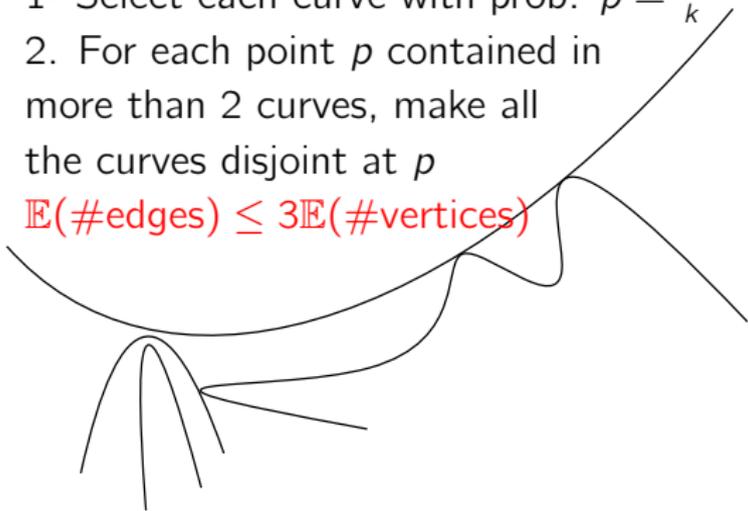
To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ek n$.

1- Select each curve with prob. $p = \frac{1}{k}$

2. For each point p contained in more than 2 curves, make all the curves disjoint at p

$$\mathbb{E}(\#edges) \leq 3\mathbb{E}(\#vertices)$$



COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

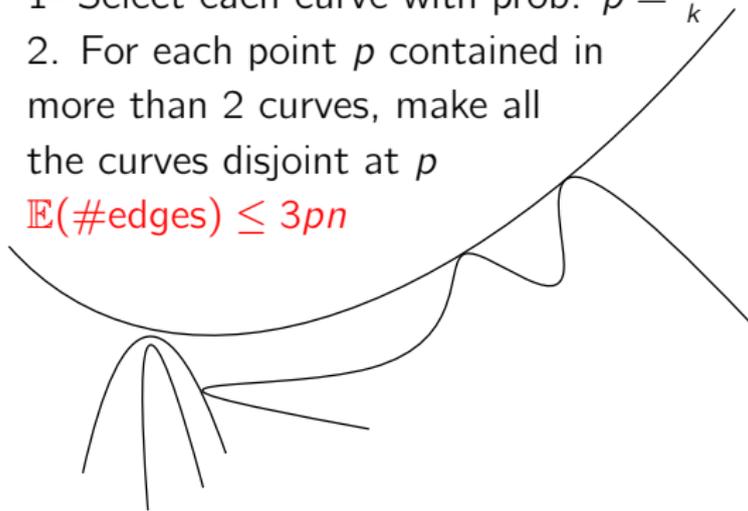
To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ek n$.

1- Select each curve with prob. $p = \frac{1}{k}$

2. For each point p contained in more than 2 curves, make all the curves disjoint at p

$$\mathbb{E}(\#edges) \leq 3pn$$

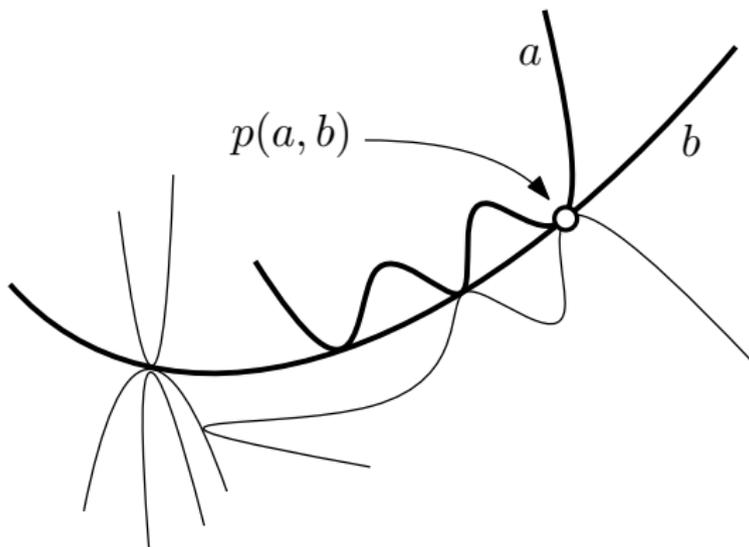


COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ekn$.



COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ek n$.

For two intersecting curves a, b , let $E_{a,b}$ be the event that a and b were selected, and no other curve containing $p(a, b)$ was selected.

COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ek n$.

For two intersecting curves a, b , let $E_{a,b}$ be the event that a and b were selected, and no other curve containing $p(a, b)$ was selected.

$\mathbb{P}(E_{a,b}) \geq p^2(1-p)^{k-2}$, so the expected number of pairs a, b for which $E_{a,b}$ holds is at least $mp^2(1-p)^{k-2}$.

COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ekn$.

For two intersecting curves a, b , let $E_{a,b}$ be the event that a and b were selected, and no other curve containing $p(a, b)$ was selected.

$\mathbb{P}(E_{a,b}) \geq p^2(1-p)^{k-2}$, so the expected number of pairs a, b for which $E_{a,b}$ holds is at least $mp^2(1-p)^{k-2}$.

$$mp^2(1-p)^{k-2} \leq \mathbb{E}(\#edges) \leq 3\mathbb{E}(\#vertices) = 3pn$$

COLORING NON-CROSSING CURVES

We consider a set of n pairwise non-crossing curves in the plane, such that each point is in at most k curves.

To prove that these curves can be properly colored with $6ek + 1$ colors, it is enough to show that their intersection graph is $6ek$ -degenerate.

To prove this, it is enough to show that if G is such an intersection graph, with m edges, then $m \leq 3ek n$.

For two intersecting curves a, b , let $E_{a,b}$ be the event that a and b were selected, and no other curve containing $p(a, b)$ was selected.

$\mathbb{P}(E_{a,b}) \geq p^2(1-p)^{k-2}$, so the expected number of pairs a, b for which $E_{a,b}$ holds is at least $mp^2(1-p)^{k-2}$.

$$mp^2(1-p)^{k-2} \leq \mathbb{E}(\#edges) \leq 3\mathbb{E}(\#vertices) = 3pn$$

$$m \leq \frac{3n}{p(1-p)^{k-2}} \leq 3ekn$$

COLORING JORDAN REGIONS

A **Jordan region** is a region of the plane bounded by a Jordan curve.

COLORING JORDAN REGIONS

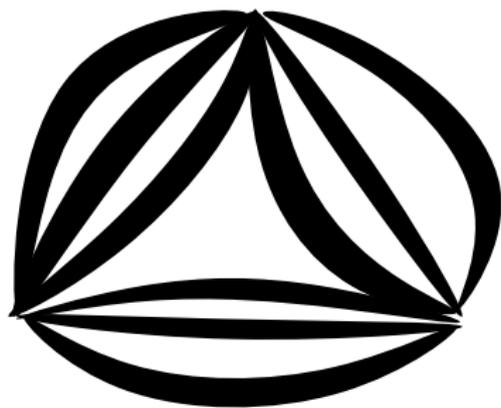
A **Jordan region** is a region of the plane bounded by a Jordan curve.

A collection of Jordan regions is **touching** if their interiors are pairwise disjoint.

COLORING JORDAN REGIONS

A **Jordan region** is a region of the plane bounded by a Jordan curve.

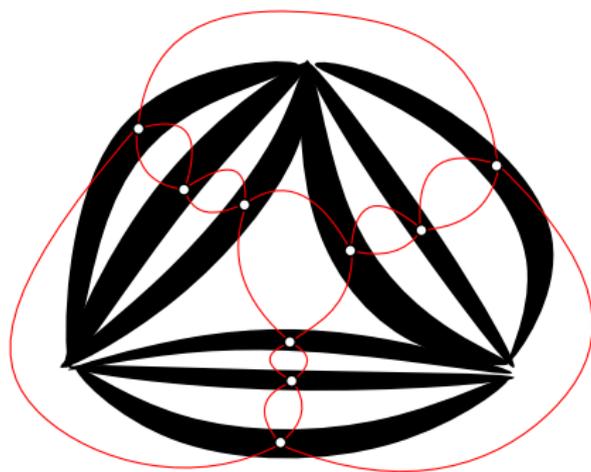
A collection of Jordan regions is **touching** if their interiors are pairwise disjoint.



COLORING JORDAN REGIONS

A **Jordan region** is a region of the plane bounded by a Jordan curve.

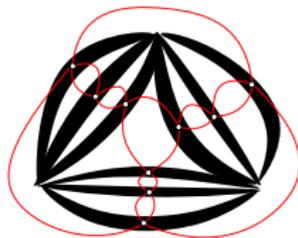
A collection of Jordan regions is **touching** if their interiors are pairwise disjoint.



COLORING JORDAN REGIONS

A **Jordan region** is a region of the plane bounded by a Jordan curve.

A collection of Jordan regions is **touching** if their interiors are pairwise disjoint.



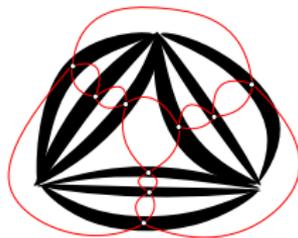
Theorem (Amini, E. & van den Heuvel 2016)

Any planar graph with **maximum face degree** k has a cyclic coloring with $\frac{3k}{2} + o(k)$ colors.

COLORING JORDAN REGIONS

A **Jordan region** is a region of the plane bounded by a Jordan curve.

A collection of Jordan regions is **touching** if their interiors are pairwise disjoint.



Theorem (Amini, E. & van den Heuvel 2016)

Any planar graph with **maximum face degree** k has a cyclic coloring with $\frac{3k}{2} + o(k)$ colors.

Consequence

Let \mathcal{F} be a touching family of Jordan regions such that each point lies in at most k regions. Then the intersection graph G of \mathcal{F} satisfies $\chi(G) \leq \frac{3k}{2} + o(k)$.

COLORING SIMPLE JORDAN REGIONS

A collection of Jordan region is **simple** if any two regions intersect in at most one point.

COLORING SIMPLE JORDAN REGIONS

A collection of Jordan region is **simple** if any two regions intersect in at most one point.

Question (Reed & Shepherd 1996)

Is there a constant C such that the intersection graph of any simple touching collection of Jordan regions satisfies $\chi(G) \leq \omega(G) + C$? Can we take $C = 1$?

COLORING SIMPLE JORDAN REGIONS

A collection of Jordan region is **simple** if any two regions intersect in at most one point.

Question (Reed & Shepherd 1996)

Is there a constant C such that the intersection graph of any simple touching collection of Jordan regions satisfies $\chi(G) \leq \omega(G) + C$? Can we take $C = 1$?

Theorem (Cames van Batenburg, E. & Müller 2016)

Let \mathcal{F} be a simple touching family of Jordan regions such that each point lies in at most k regions. Then the intersection graph G of \mathcal{F} satisfies $\chi(G) \leq k + 327$ (and $\chi(G) \leq k + 1$ if $k \geq 490$).

COLORING CONTACT SYSTEMS OF STRINGS

A collection of strings is a **contact system** if the interiors of any two strings have empty intersection.

COLORING CONTACT SYSTEMS OF STRINGS

A collection of strings is a **contact system** if the interiors of any two strings have empty intersection. A contact system of strings is **one-sided** if at any intersection point, all strings leave from the same side.

COLORING CONTACT SYSTEMS OF STRINGS

A collection of strings is a **contact system** if the interiors of any two strings have empty intersection. A contact system of strings is **one-sided** if at any intersection point, all strings leave from the same side.

Question (Hliněný 1998)

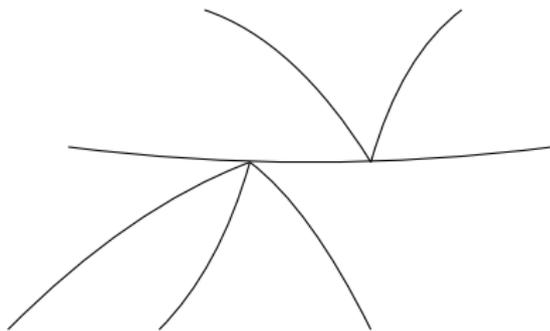
Let \mathcal{S} be a **one-sided contact system of strings**, such that any point of the plane is in **at most k strings**, and any two strings intersect **in at most one point**. Is it true that the intersection graph of \mathcal{S} has chromatic number at most **$k + o(k)$** ? (or even **$k + c$** , for some constant c ?)

COLORING CONTACT SYSTEMS OF STRINGS

A collection of strings is a **contact system** if the interiors of any two strings have empty intersection. A contact system of strings is **one-sided** if at any intersection point, all strings leave from the same side.

Question (Hliněný 1998)

Let \mathcal{S} be a **one-sided contact system of strings**, such that any point of the plane is in **at most k strings**, and any two strings intersect **in at most one point**. Is it true that the intersection graph of \mathcal{S} has chromatic number at most $k + o(k)$? (or even $k + c$, for some constant c ?)

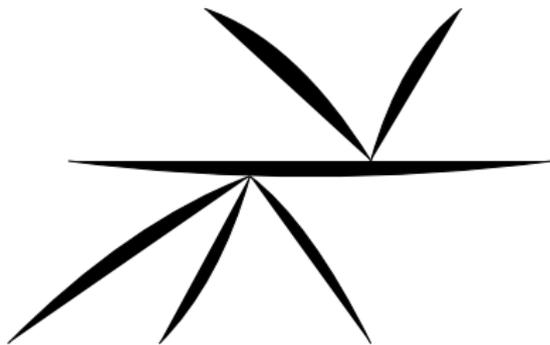


COLORING CONTACT SYSTEMS OF STRINGS

A collection of strings is a **contact system** if the interiors of any two strings have empty intersection. A contact system of strings is **one-sided** if at any intersection point, all strings leave from the same side.

Question (Hliněný 1998)

Let \mathcal{S} be a **one-sided contact system of strings**, such that any point of the plane is in **at most k strings**, and any two strings intersect **in at most one point**. Is it true that the intersection graph of \mathcal{S} has chromatic number at most $k + o(k)$? (or even $k + c$, for some constant c ?)



A CONJECTURE

Let us say that a collection of strings is *k-touching* if the strings are pairwise non-crossing, and any point of the plane is contained in at most k strings.

A CONJECTURE

Let us say that a collection of strings is *k-touching* if the strings are pairwise non-crossing, and any point of the plane is contained in at most k strings.

We also say that the collection is *simple* if any two strings intersect in at most one point.

A CONJECTURE

Let us say that a collection of strings is ***k*-touching** if the strings are pairwise non-crossing, and any point of the plane is contained in at most k strings.

We also say that the collection is **simple** if any two strings intersect in at most one point.

Conjecture (E., Gonçalves, & Labourel 2009)

There a constant C such that the intersection graph of any simple k -touching collection of strings satisfies $\chi(G) \leq k + C$.

A CONJECTURE

Let us say that a collection of strings is ***k*-touching** if the strings are pairwise non-crossing, and any point of the plane is contained in at most k strings.

We also say that the collection is **simple** if any two strings intersect in at most one point.

Conjecture (E., Gonçalves, & Labourel 2009)

There a constant C such that the intersection graph of any simple k -touching collection of strings satisfies $\chi(G) \leq k + C$.

- We can prove $k + 5$ for segments.

A CONJECTURE

Let us say that a collection of strings is ***k*-touching** if the strings are pairwise non-crossing, and any point of the plane is contained in at most k strings.

We also say that the collection is **simple** if any two strings intersect in at most one point.

Conjecture (E., Gonçalves, & Labourel 2009)

There a constant C such that the intersection graph of any simple k -touching collection of strings satisfies $\chi(G) \leq k + C$.

- We can prove $k + 5$ for segments.
- $2k + c$ is also not hard to derive for closed curves.