

COLORING JORDAN REGIONS AND CURVES

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New trends in Graph Coloring, Banff
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PACKING CYCLES IN PLANAR DIGRAPHS

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Theorem (Reed & Shepherd 1996)

For every **planar** digraph G , $\nu^*(G) \leq 28 \nu(G)$.

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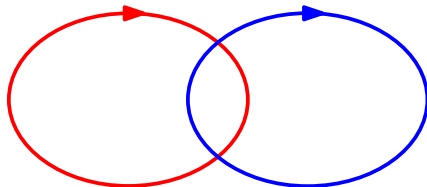
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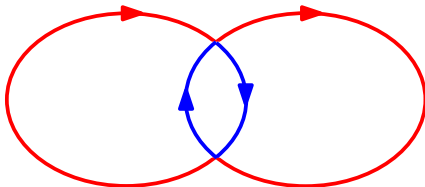
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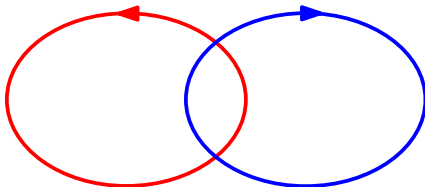
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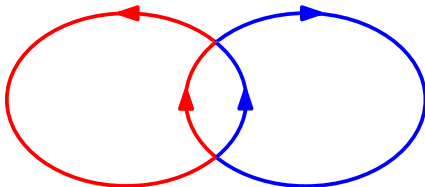
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In particular, if there are n curves, then there are at least $\frac{n}{6ek}$ pairwise disjoint curves. So $\nu(G) \geq \frac{n}{6ek} = \frac{\nu^*(G)}{6e}$, and

$$\nu^*(G) \leq 6e \nu(G) \approx 16.3 \nu(G)$$

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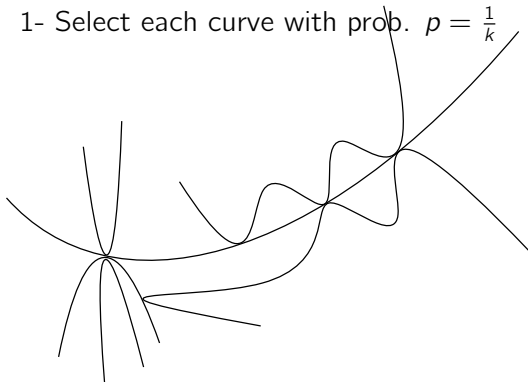
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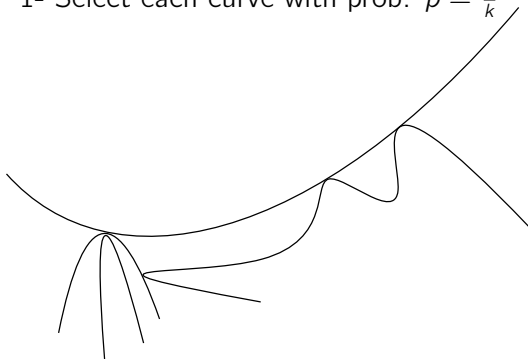
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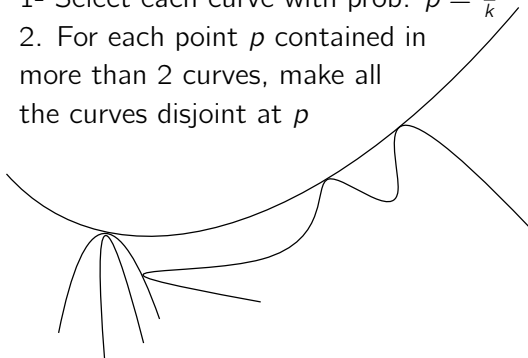
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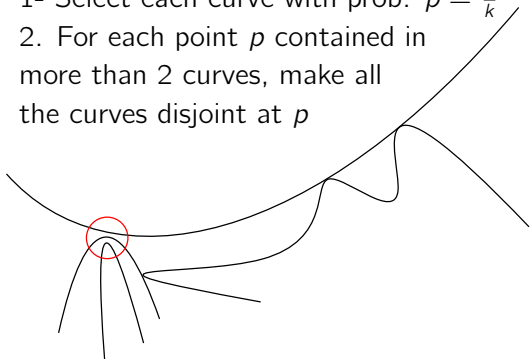
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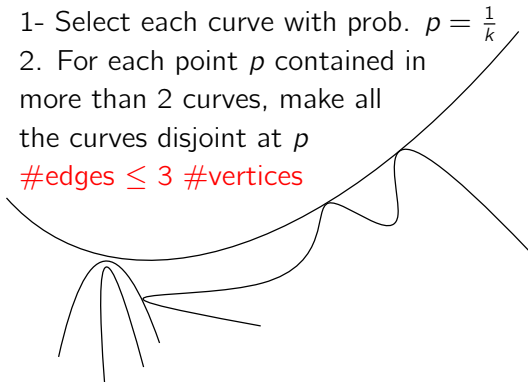
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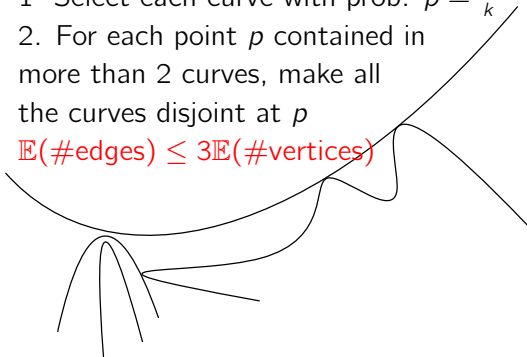
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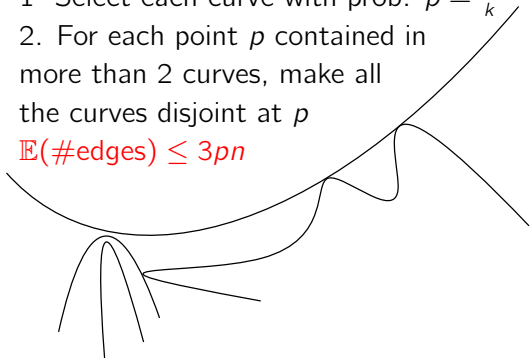
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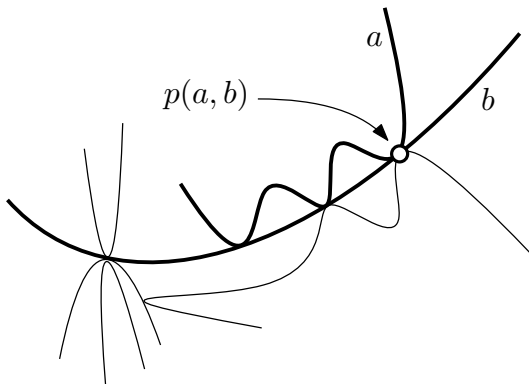


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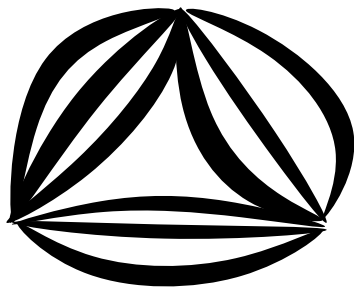
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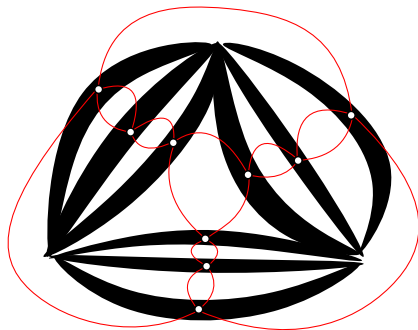
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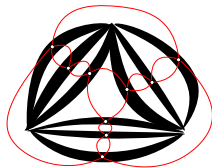
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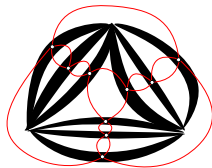
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Consequence

Let \mathcal{F} be a touching family of Jordan regions such that each point lies in at most k regions. Then the intersection graph G of \mathcal{F} satisfies $\chi(G) \leq \frac{3k}{2} + o(k)$.

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Theorem (Cames van Batenburg, E. & Müller 2016)

Let \mathcal{F} be a simple touching family of Jordan regions such that each point lies in at most k regions. Then the intersection graph G of \mathcal{F} satisfies $\chi(G) \leq k + 327$ (and $\chi(G) \leq k + 1$ if $k \geq 490$).

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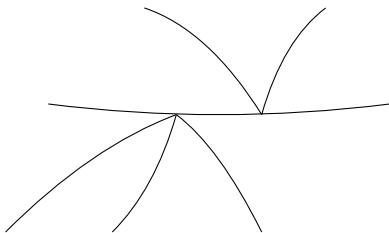
Let \mathcal{S} be a **one-sided contact system of strings**, such that any point of the plane is in **at most k strings**, and any two strings intersect **in at most one point**. Is it true that the intersection graph of \mathcal{S} has chromatic number at most **$k + o(k)$** ? (or even **$k + c$** , for some constant c ?)

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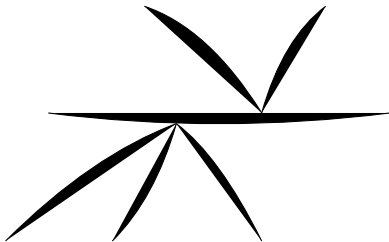


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- We can prove $k + 5$ for segments.
- $2k + c$ is also not hard to derive for closed curves.