

Robust Multicast Beamforming in Cognitive Radio Networks: Semidefinite Relaxation and Approximation Analysis

Anthony Man-Cho So

(Joint Work with Xiaoxiao Wu, Man-Chung Yue, and Wing-Kin Ma)

Department of Systems Engineering & Engineering Management
The Chinese University of Hong Kong (CUHK), Hong Kong

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Outline of Talk

- Robust multicast beamforming in cognitive radio networks
- Robust quadratically constrained quadratic optimization (QCQO) formulation
- Semidefinite relaxation of the robust QCQO formulation
- Approximation analysis of the semidefinite relaxation

Introduction

- The demand for wireless services has been ever increasing.
 - live mobile TV
 - multi-party video conferencing
 - multimedia streaming for groups of paid users



From <http://money.cnn.com>



From <http://www.telepresenceoptions.com>



From <http://www.slashgear.com>

- This creates a great demand for spectrum-based communication links.
- However, there is a shortage of available frequencies, as most have been allocated to licensed users.

Introduction

UNITED STATES FREQUENCY ALLOCATIONS THE RADIO SPECTRUM

RADIO SERVICES COLOR LEGEND

AERIAL MOBILE	INTER-SATELLITE	RADIO ASTRONOMY
AERIAL MOBILE SATELLITE	LAND MOBILE	RADIO DETERMINATION SATELLITE
AERIAL MOBILE SATELLITE	LAND MOBILE SATELLITE	RADIOLOCATION
AMATEUR	MARITIME MOBILE	RADIOLOCATION SATELLITE
AMATEUR SATELLITE	MARITIME MOBILE SATELLITE	RADIOCOMMUNICATION
BROADCASTING	MARITIME MOBILE SATELLITE	RADIOCOMMUNICATION SATELLITE
BROADCASTING SATELLITE	METEOROLOGICAL	SPACE OPERATION
EARLY EXPERIMENTAL SATELLITE	METEOROLOGICAL SATELLITE	SPACE RESEARCH
FIXED	MOBILE	STANDARD FREQUENCY AND TIME SIGNAL
FIXED SATELLITE	MOBILE SATELLITE	STANDARD FREQUENCY AND TIME SIGNAL SATELLITE

ACTIVITY CODE

FEDERAL EXCLUSIVE	FEDERAL NON-FEDERAL SHARED
NON-FEDERAL EXCLUSIVE	

ALLOCATION USAGE DESIGNATION

SERVICE	EXAMPLE	DESCRIPTION
Primary	Radio	Capable of one-way or two-way communication
Secondary	Radio	Not capable of one-way or two-way communication

U.S. DEPARTMENT OF COMMERCE
National Telecommunications and Information Administration
Office of Spectrum Management
JANUARY 2016

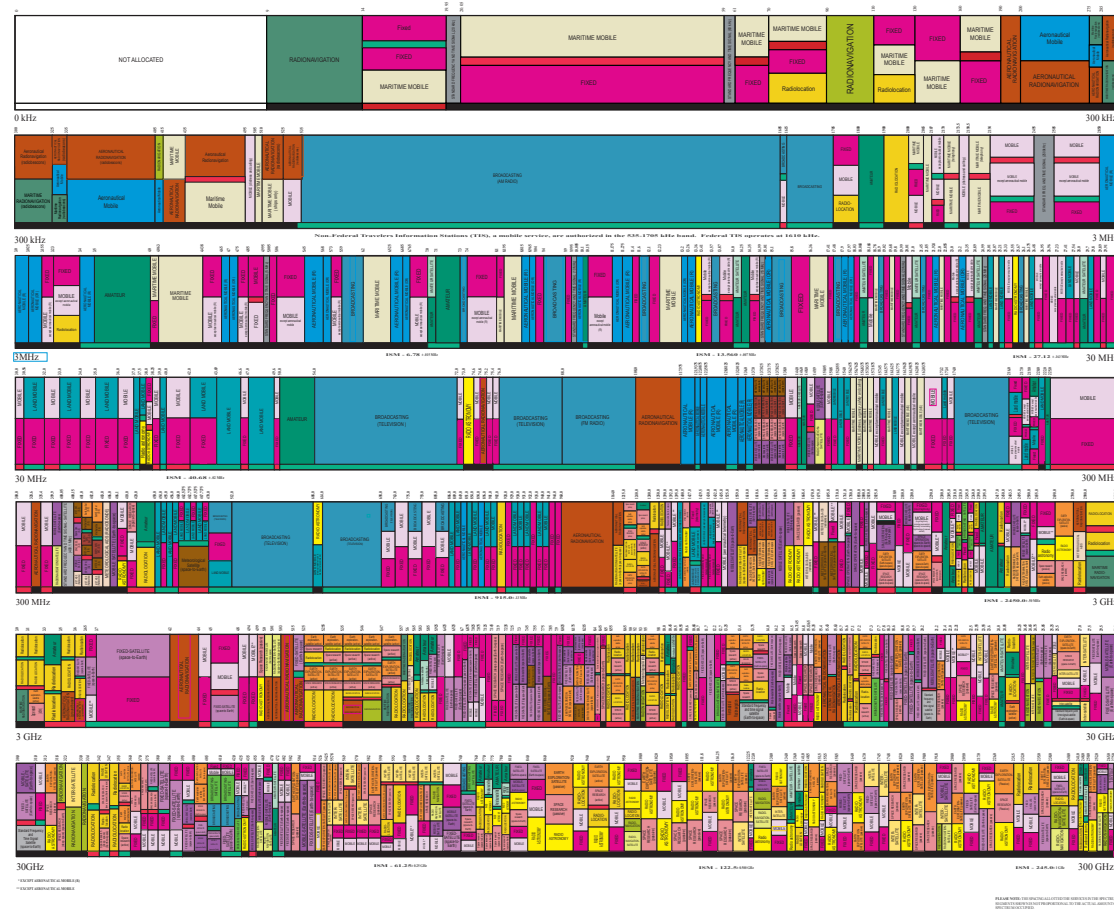


Figure 1: US Frequency Allocation Chart, as of January 2016. (Source: National Telecommunications and Information Administration, US Department of Commerce)

Introduction



香港無線電頻率劃分圖 Hong Kong Frequency Allocation Chart

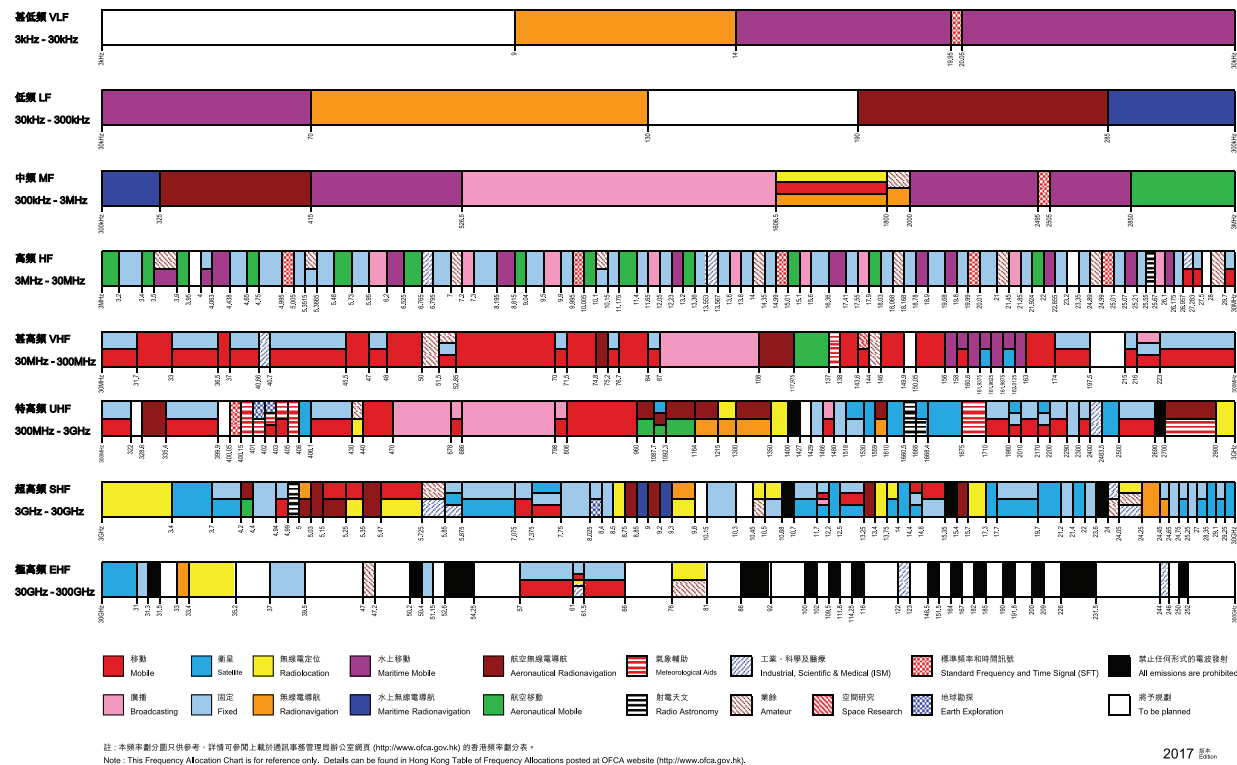


Figure 2: Hong Kong Frequency Allocation Chart, as of 2017. (Source: Office of the Communications Authority, The Government of the HKSAR)

Cognitive Radio Systems

- Cognitive radio (CR) technologies, which can adapt a radio's use of spectrum to real-time conditions of its operating environment, has emerged as a promising technology for improving spectrum utilization and bandwidth efficiency.
- In a CR network, secondary (unlicensed) users (SUs) are allowed to operate at the same frequency bands as the primary (licensed) users (PUs).
- We consider an underlay-CR network, in which both the PUs and SUs can use the frequency bands simultaneously.
- Naturally, the SUs should not cause excessive interference to the PUs.
- Current standards supporting CR: TV-broadcast bands (IEEE 802.22), LTE Advanced.

System Model

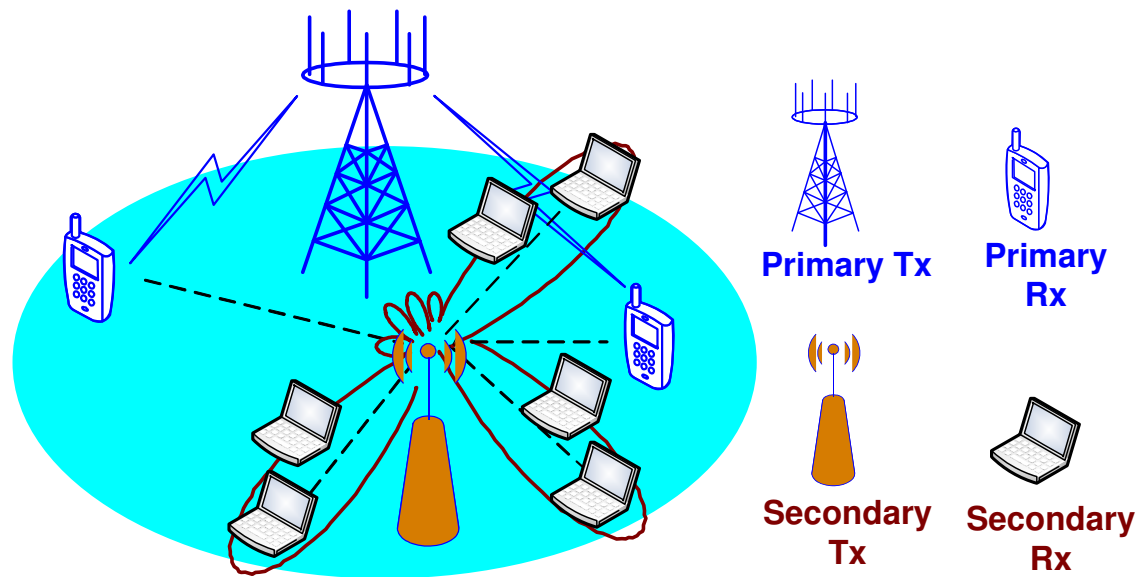


Figure 3: The multi-group multicast cognitive radio system model.

- Cognitive radio (CR) and multi-group multicast (MM) delivery play a significant role in supporting resource- and spectral-efficient data services in future communication systems.
- In general, one is interested in transmit designs in the MM-CR network.
- For simplicity, we focus on the scenario where there is only one group of SUs.

Problem Formulation

- Consider a secondary transmitter (ST) with N transmit antennae sending common information (i.e., multicasting) to M single-antenna SUs in the presence of J single-antenna PUs.
- By adopting the so-called **transmit beamforming** scheme [**Sidiropoulos-Davidson-Luo'06**], the received signal for user i ($i = 1, \dots, M$) is modeled as

$$y_i = \mathbf{h}_i^H \mathbf{w} s + n_i,$$

where

- $\mathbf{h}_i \in \mathbb{C}^N$: channel between the ST and user i ;
 - $\mathbf{w} \in \mathbb{C}^N$: ST's beamforming vector;
 - $s \in \mathbb{C}$: unit power data stream;
 - $n_i \sim \mathcal{CN}(0, \sigma_i^2)$: additive complex Gaussian noise at user i .
- The quality of service received by user i is then measured by the **signal-to-noise ratio** (SNR):

$$\text{SNR}_i(\mathbf{w}) = \frac{|\mathbf{h}_i^H \mathbf{w}|^2}{\sigma_i^2}.$$

Problem Formulation

- On the other hand, the interference power at the j -th PU is

$$\text{INT}_j(\mathbf{w}) = |\mathbf{g}_j^H \mathbf{w}|^2,$$

where

- $\mathbf{g}_j \in \mathbb{C}^N$: channel between the ST and PU j .
- Our goal is to design the beamforming vector $\mathbf{w} \in \mathbb{C}^N$ under the **max-min-fair (MMF)** criterion, subject to a **power constraint** at the ST and so-called **interference temperature** constraints at the PUs:

$$\begin{aligned} & \max_{\mathbf{w} \in \mathbb{C}^N} && \min_{i=1, \dots, M} \text{SNR}_i(\mathbf{w}) \\ & \text{subject to} && \|\mathbf{w}\|_2^2 \leq P, \\ & && \text{INT}_j(\mathbf{w}) \leq \beta_j \quad \text{for } j = 1, \dots, J. \end{aligned}$$

Here,

- $P > 0$: maximum allowable transmit power of the ST;
- $\beta_j > 0$: interference threshold of the PU j .

Incorporating Robustness

- Right now, our problem formulation assumes that the channels $\{\mathbf{h}_i\}_{i=1}^M$ and $\{\mathbf{g}_j\}_{j=1}^J$ are perfectly known.
- This may be reasonable for the ST-SU channels $\{\mathbf{h}_i\}_{i=1}^M$ (as the ST and the SUs can exchange channel state information), but certainly not for the ST-PU channels $\{\mathbf{g}_j\}_{j=1}^J$ (as the PUs will not voluntarily reveal information to the ST).
- Thus, the channels between the ST and the PUs have to be estimated.
- A commonly adopted model for the ST-PU channels:

$$\mathbf{g}_j = \bar{\mathbf{g}}_j + \mathbf{\Delta}_j, \quad \|\mathbf{\Delta}_j\|_2 \leq \delta_j,$$

where

- $\bar{\mathbf{g}}_j \in \mathbb{C}^N$: estimated channel between the ST and PU j ;
- $\mathbf{\Delta}_j \in \mathbb{C}^N$: channel estimation error associated with PU j ;
- $\delta_j \geq 0$: error threshold associated with PU j .

Robust Multicast Beamforming in CR Networks

- This yields the following robust formulation:

$$\begin{aligned} v^* &= \max_{\mathbf{w} \in \mathbb{C}^N} \min_{i=1, \dots, M} \text{SNR}_i(\mathbf{w}) \\ &\text{subject to } \|\mathbf{w}\|_2^2 \leq P, \\ &\quad \max_{\|\Delta_j\|_2 \leq \delta_j} \text{INT}_j(\mathbf{w}, \Delta_j) \leq \beta_j \quad \text{for } j = 1, \dots, J. \end{aligned} \tag{R-MMF-BF}$$

Here,

$$\text{SNR}_i(\mathbf{w}) = \frac{|\mathbf{h}_i^H \mathbf{w}|^2}{\sigma_i^2}, \quad \text{INT}_j(\mathbf{w}, \Delta_j) = \left| (\bar{\mathbf{g}}_j + \Delta_j)^H \mathbf{w} \right|^2.$$

By re-defining \mathbf{h}_i if necessary, we may assume that $\sigma_i^2 = 1$ for $i = 1, \dots, M$.

- Even when $\delta_j = 0$ for $j = 1, \dots, J$, (R-MMF-BF) is still a **non-convex** QCQO problem and is **NP-hard** in general [**Luo-Sidiropoulos-Tseng-Zhang'07**] [**Karipidis-Sidiropoulos-Luo'08**].

Semidefinite Relaxation of (R-MMF-BF)

- A classic approach to tackling (R-MMF-BF) is to apply **semidefinite relaxation**. First, observe that

$$\mathbf{W} = \mathbf{w}\mathbf{w}^H \iff \mathbf{W} \succeq \mathbf{0}, \quad \text{rank}(\mathbf{W}) \leq 1.$$

Moreover, by the \mathcal{S} -lemma, we have

$$\text{INT}_j(\mathbf{w}, \mathbf{\Delta}_j) = \max_{\|\mathbf{\Delta}_j\|_2 \leq \delta_j} \left| (\bar{\mathbf{g}}_j + \mathbf{\Delta}_j)^H \mathbf{w} \right|^2 \leq \beta_j$$

if and only if there exists a $\kappa_j \geq 0$ such that

$$\begin{bmatrix} \kappa_j \mathbf{I} - \mathbf{w}\mathbf{w}^H & -\mathbf{w}^H \bar{\mathbf{g}}_j \mathbf{w} \\ -\bar{\mathbf{g}}_j^H \mathbf{w}\mathbf{w}^H & \beta_j - \left| \mathbf{w}^H \bar{\mathbf{g}}_j \right|^2 - \delta_j^2 \kappa_j \end{bmatrix} \succeq \mathbf{0}.$$

Semidefinite Relaxation of (R-MMF-BF)

- Putting everything together, we can reformulate (R-MMF-BF) as the following rank-constrained SDP:

$$\begin{aligned}
 & \max_{\mathbf{W}, t, \kappa_1, \dots, \kappa_J} && t \\
 \text{subject to} &&& \mathbf{h}_i^H \mathbf{W} \mathbf{h}_i \geq t \quad \text{for } i = 1, \dots, M, \\
 &&& \text{Tr}(\mathbf{W}) \leq P, \\
 &&& \begin{bmatrix} \kappa_j \mathbf{I} - \mathbf{W} & -\mathbf{W} \bar{\mathbf{g}} \\ -\bar{\mathbf{g}}_j^H \mathbf{W} & \beta_j - \bar{\mathbf{g}}_j^H \mathbf{W} \bar{\mathbf{g}}_j - \delta_j^2 \kappa_j \end{bmatrix} \succeq \mathbf{0}, \quad \kappa_j \geq 0 \\
 &&& \text{for } j = 1, \dots, J, \\
 &&& \mathbf{W} \succeq \mathbf{0}, \quad \text{rank}(\mathbf{W}) \leq 1.
 \end{aligned} \tag{SDR}$$

- By dropping the rank constraint, we obtain a semidefinite relaxation of the original problem.

Quality of the Semidefinite Relaxation

- Let $(\mathbf{W}^*, t^*, \kappa_1^*, \dots, \kappa_J^*)$ be an optimal solution to (SDR). Clearly, we have $t^* \geq v^*$ because (SDR) is a relaxation of (R-MMF-BF).
- If we have $\text{rank}(\mathbf{W}^*) \leq 1$, then we have found an optimal solution to the original problem (R-MMF-BF).
- In general, there is no guarantee that \mathbf{W}^* will satisfy the rank constraint.

- **Questions**

Let \mathbf{W}^* be an optimal solution to (SDR). If $\text{rank}(\mathbf{W}^*) > 1$, how do we generate from \mathbf{W}^* a feasible solution $\hat{\mathbf{w}}$ to the original problem (R-MMF-BF)? Can we say something about the quality of the approximate solution $\hat{\mathbf{w}}$?

– In particular, can we find $\alpha \in (0, 1)$ (called the **approximation ratio**) such that

$$v^* \geq \min_{i=1, \dots, M} \text{SNR}_i(\hat{\mathbf{w}}) \geq \alpha \cdot v^*?$$

- A classic idea: Perform “**randomized rounding**”

Handling High-Rank SDP Solutions

- Suppose that $\text{rank}(\mathbf{W}^*) > 1$. Generate $\boldsymbol{\xi} \sim \mathcal{CN}(\mathbf{0}, \mathbf{W}^*)$.
- Why may this be a reasonable idea? Observe that

$$\mathbb{E} [\text{SNR}_i(\boldsymbol{\xi})] = \mathbb{E} \left[|\mathbf{h}_i^H \boldsymbol{\xi}|^2 \right] = \mathbf{h}_i^H \mathbf{W}^* \mathbf{h}_i \geq t^*,$$

$$\mathbb{E} [\|\boldsymbol{\xi}\|_2^2] = \text{Tr}(\mathbf{W}^*) \leq P.$$

In other words, **in expectation**, $\boldsymbol{\xi}$ satisfies the power constraint and achieves the optimal value of (SDR).

- However, without knowing the concentration properties of $\text{SNR}_i(\boldsymbol{\xi})$ and $\|\boldsymbol{\xi}\|_2^2$, the above may not mean much.
- Moreover, it is not clear whether the interference constraints are satisfied even in expectation:

$$\mathbb{E} \left[\max_{\|\boldsymbol{\Delta}_j\|_2 \leq \delta_j} \text{INT}_j(\boldsymbol{\xi}, \boldsymbol{\Delta}_j) \right] = \mathbb{E} \left[\max_{\|\boldsymbol{\Delta}_j\|_2 \leq \delta_j} \left| (\bar{\mathbf{g}}_j + \boldsymbol{\Delta}_j)^H \boldsymbol{\xi} \right|^2 \right] \stackrel{?}{\leq} \beta_j.$$

Analysis of the Randomization Procedure: Warm Up

- Consider first the case where $\delta_j = 0$ for $j = 1, \dots, J$; i.e., there is no channel estimation error. Then, all robust constraints reduce to quadratic constraints.
- This case has been well studied in the literature [**Nemirovski-Roos-Terlaky'99**], [**Luo-Sidiropoulos-Tseng-Zhang'07**], [**S-Ye-Zhang'08**]. The key lies in the following deviation inequalities for complex Gaussian quadratic forms:

Theorem [S-Ye-Zhang'08]: Let $\mathbf{Q}, \mathbf{W}^* \succeq \mathbf{0}$ be given. Let $\boldsymbol{\xi} \sim \mathcal{CN}(\mathbf{0}, \mathbf{W}^*)$. If $\mathbf{Q}\mathbf{W}^* = \mathbf{0}$, then $\boldsymbol{\xi}^H \mathbf{Q} \boldsymbol{\xi} = 0$ a.s. Otherwise, for any $\mu > 1$ and $\nu \in (0, 1)$,

$$\Pr(\boldsymbol{\xi}^H \mathbf{Q} \boldsymbol{\xi} \geq \mu \cdot \text{Tr}(\mathbf{Q}\mathbf{W}^*)) \leq \exp(1 - \mu + \ln \mu),$$

$$\Pr(\boldsymbol{\xi}^H \mathbf{Q} \boldsymbol{\xi} \leq \nu \cdot \text{Tr}(\mathbf{Q}\mathbf{W}^*)) \leq \exp(1 - \nu + \ln \nu).$$

Analysis of the Randomization Procedure: Warm Up

- By the theorem and the union bound, all the following inequalities

$$\text{SNR}_i(\boldsymbol{\xi}) = |\mathbf{h}_i^H \boldsymbol{\xi}|^2 \geq \nu \cdot \text{Tr}(\mathbf{h}_i \mathbf{h}_i^H \mathbf{W}^*) \quad \text{for } i = 1, \dots, M,$$

$$\|\boldsymbol{\xi}\|_2^2 \leq \mu \cdot \text{Tr}(\mathbf{W}^*),$$

$$\text{INT}_j(\boldsymbol{\xi}) = |\bar{\mathbf{g}}_j^H \boldsymbol{\xi}|^2 \leq \mu \cdot \text{Tr}(\bar{\mathbf{g}}_j \bar{\mathbf{g}}_j^H \mathbf{W}^*) \quad \text{for } j = 1, \dots, J$$

hold with probability at least $1 - (J + 1) \exp(1 - \mu + \ln \mu) - M \exp(1 - \nu + \ln \nu)$.

- Since

$$\text{Tr}(\mathbf{h}_i \mathbf{h}_i^H \mathbf{W}^*) \geq t^*, \quad \text{Tr}(\mathbf{W}^*) \leq P, \quad \text{Tr}(\bar{\mathbf{g}}_j \bar{\mathbf{g}}_j^H \mathbf{W}^*) \leq \beta_j,$$

by choosing $\mu = O(\ln J)$ and $\nu = \Omega(1/M)$, we conclude that with constant probability,

– $\hat{\mathbf{w}} = \boldsymbol{\xi} / \sqrt{\mu}$ is **feasible** for the original problem (R-MMF-BF),

– $\text{SNR}_i(\hat{\mathbf{w}}) = |\mathbf{h}_i^H \hat{\mathbf{w}}|^2 = |\mathbf{h}_i^H \boldsymbol{\xi}|^2 / \mu \geq (\nu / \mu) t^* \geq (\nu / \mu) v^*$.

- Hence, $\hat{\mathbf{w}}$ is an **(ν / μ) -approximate solution** to the original problem (R-MMF-BF).

Summary of the Warm-Up Case

- Recall the problem we are trying to solve:

$$\begin{aligned} & \max_{\mathbf{w} \in \mathbb{C}^N} && \min_{i=1, \dots, M} \text{SNR}_i(\mathbf{w}) \\ & \text{subject to} && \|\mathbf{w}\|_2^2 \leq P, \\ & && \text{INT}_j(\mathbf{w}) \leq \beta_j \quad \text{for } j = 1, \dots, J. \end{aligned} \tag{MMF-BF}$$

- The proposed algorithm
 - Solve the semidefinite relaxation and get the optimal solution \mathbf{W}^* .
 - Generate $\boldsymbol{\xi} \sim \mathcal{CN}(\mathbf{0}, \mathbf{W}^*)$ and scale it down by

$$\mu = \max \left\{ \frac{\|\boldsymbol{\xi}\|_2^2}{P}, \max_{j=1, \dots, J} \frac{\text{INT}_j(\boldsymbol{\xi})}{\beta_j} \right\}$$

to get a feasible solution $\hat{\mathbf{w}}$ to (MMF-BF).

- Our analysis showed that with high probability, $\mu = O(\ln J)$ and $\hat{\mathbf{w}}$ is an $\Omega(1/M \ln J)$ -approximate solution.
- Note that the ratio degrades **linearly** in the number of SUs (M) but only **logarithmically** in the number of PUs (J).

Analysis of the Randomization Procedure: General Case

- Now, let us tackle the general case. Again, we generate $\boldsymbol{\xi} \sim \mathcal{CN}(\mathbf{0}, \mathbf{W}^*)$.
- Same as before, with probability at least $1 - \exp(1 - \mu + \ln \mu) - M \exp(1 - \nu + \ln \nu)$, all the following inequalities hold:

$$\begin{aligned} \text{SNR}_i(\boldsymbol{\xi}) &= |\mathbf{h}_i^H \boldsymbol{\xi}|^2 \geq \nu \cdot \text{Tr}(\mathbf{h}_i \mathbf{h}_i^H \mathbf{W}^*) \quad \text{for } i = 1, \dots, M, \\ \|\boldsymbol{\xi}\|_2^2 &\leq \mu \cdot \text{Tr}(\mathbf{W}^*). \end{aligned}$$

Analysis of the Randomization Procedure: General Case

- Recalling that $\text{INT}_j(\boldsymbol{\xi}, \boldsymbol{\Delta}_j) = |\mathbf{g}_j^H \boldsymbol{\xi}|^2$ with $\mathbf{g}_j = \bar{\mathbf{g}}_j + \boldsymbol{\Delta}_j$, consider

$$\max_{\|\boldsymbol{\Delta}_j\|_2 \leq \delta_j} \text{INT}_j(\boldsymbol{\xi}, \boldsymbol{\Delta}_j) = \max_{\|\boldsymbol{\Delta}_j\|_2 = \delta_j} \left| (\bar{\mathbf{g}}_j + \boldsymbol{\Delta}_j)^H \boldsymbol{\xi} \right|^2.$$

For simplicity, let us take $\delta_j = 1$ and define

$$\boldsymbol{\Delta}_j(\boldsymbol{\xi}) = \arg \max_{\|\boldsymbol{\Delta}_j\|_2 = 1} \left| (\bar{\mathbf{g}}_j + \boldsymbol{\Delta}_j)^H \boldsymbol{\xi} \right|^2.$$

Analysis of the Randomization Procedure: General Case

- Two tempting ideas
 - Apply the previous theorem to claim

$$\left| (\bar{\mathbf{g}}_j + \Delta_j(\boldsymbol{\xi}))^H \boldsymbol{\xi} \right|^2 \leq \mu \cdot \text{Tr} \left((\bar{\mathbf{g}}_j + \Delta_j(\boldsymbol{\xi})) (\bar{\mathbf{g}}_j + \Delta_j(\boldsymbol{\xi}))^H \mathbf{W}^* \right)$$

with high probability. However, $\Delta_j(\boldsymbol{\xi})$ depends on $\boldsymbol{\xi}$ and hence we cannot apply the theorem directly.

- Let \mathbb{S}^{N-1} denote the unit $(N - 1)$ -sphere. Then, for a fixed $\Delta_j \in \mathbb{S}^{N-1}$, the previous theorem gives

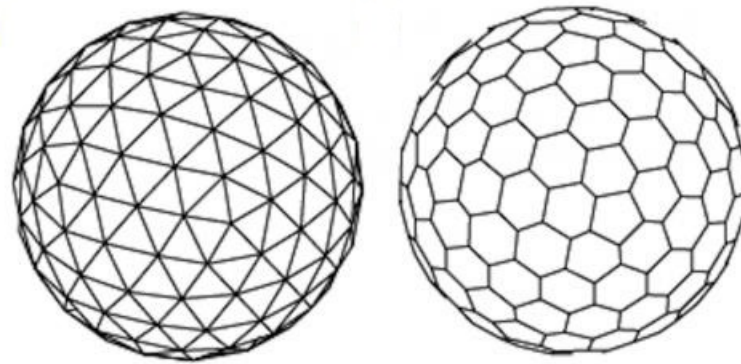
$$\left| (\bar{\mathbf{g}}_j + \Delta_j)^H \boldsymbol{\xi} \right|^2 \leq \mu \cdot \text{Tr} \left((\bar{\mathbf{g}}_j + \Delta_j) (\bar{\mathbf{g}}_j + \Delta_j)^H \mathbf{W}^* \right)$$

with high probability. However, we cannot take the union bound over \mathbb{S}^{N-1} !

- The notion of an ϵ -net comes to the rescue.

Analysis of the Randomization Procedure: General Case

- **Definition:** Let $\epsilon > 0$ be given. We say that $\mathcal{N} \subset \mathbb{S}^{N-1}$ is an ϵ -net of \mathbb{S}^{N-1} if for any $z \in \mathbb{S}^{N-1}$, there exists a $u \in \mathcal{N}$ such that $\|u - z\|_2 \leq \epsilon$.



- It is well known (and can be easily established via a volume argument) that given $\epsilon > 0$, there exists an ϵ -net \mathcal{N}_ϵ of \mathbb{S}^{N-1} of size at most $(1 + 2/\epsilon)^{2N}$.
 - More generally, there exists an ϵ -net $\mathcal{N}_\epsilon(\delta)$ of $\delta\mathbb{S}^{N-1}$ of size at most $(\delta(1 + 2/\epsilon))^{2N}$.

Analysis of the Randomization Procedure: General Case

- Now, take an ϵ -net \mathcal{N}_ϵ of \mathbb{S}^{N-1} with $\epsilon \in (0, 1)$. It can be shown that there exist sequences $\{\epsilon_k\}_{k \geq 0}$ and $\{\mathbf{u}^k\}_{k \geq 0}$ with $\epsilon_k \in [0, \epsilon^k]$ and $\mathbf{u}^k \in \mathcal{N}_\epsilon$ such that

$$\Delta_j(\boldsymbol{\xi}) = \sum_{k \geq 0} \epsilon_k \mathbf{u}^k.$$

- Let $S = \left(\sum_{k \geq 0} \epsilon_k \right)^{-1} \geq 1 - \epsilon$. Using the above, we can then show that

$$\begin{aligned} \max_{\|\Delta_j\|_2 \leq 1} \text{INT}_j(\boldsymbol{\xi}, \Delta_j) &= \left| (\bar{\mathbf{g}}_j + \Delta_j(\boldsymbol{\xi}))^H \boldsymbol{\xi} \right|^2 \\ &\leq \frac{2}{S^2} \max_{\mathbf{u} \in \mathcal{N}_\epsilon} \left| (\bar{\mathbf{g}}_j + \mathbf{u})^H \boldsymbol{\xi} \right|^2 + 2 \left(\frac{1}{S} - 1 \right)^2 \left| \bar{\mathbf{g}}_j^H \boldsymbol{\xi} \right|^2. \end{aligned}$$

The upshot is that both terms can be tackled by our previous theorem (the first entails taking a union bound over all points in \mathcal{N}_ϵ).

Analysis of the Randomization Procedure: General Case

- Carrying out the necessary calculations, we get the following:

Theorem: Fix $j \in \{1, \dots, J\}$. For any $\mu > 1$ and $\epsilon \in (0, 1)$, the inequality

$$\max_{\|\Delta_j\|_2 \leq 1} \text{INT}_j(\boldsymbol{\xi}, \Delta_j) \leq \frac{2\mu(1 + \epsilon^2)}{(1 - \epsilon)^2} \max_{\|\Delta_j\|_2 \leq 1} \text{Tr} \left((\bar{\mathbf{g}}_j + \Delta_j) (\bar{\mathbf{g}}_j + \Delta_j)^H \mathbf{W}^* \right)$$

holds with probability at least $1 - (|\mathcal{N}_\epsilon| + 1) \exp(1 - \mu + \ln \mu)$.

- Summing up, with probability at least $1 - (J(|\mathcal{N}_\epsilon| + 1) + 1) \exp(1 - \mu + \ln \mu) - M \exp(1 - \nu + \ln \nu)$, all the following inequalities hold:

$$\text{SNR}_i(\boldsymbol{\xi}) = |\mathbf{h}_i^H \boldsymbol{\xi}|^2 \geq \nu \cdot \text{Tr}(\mathbf{h}_i \mathbf{h}_i^H \mathbf{W}^*) \quad \text{for } i = 1, \dots, M,$$

$$\|\boldsymbol{\xi}\|_2^2 \leq \mu \cdot \text{Tr}(\mathbf{W}^*),$$

$$\max_{\|\Delta_j\|_2 \leq 1} \text{INT}_j(\boldsymbol{\xi}, \Delta_j) \leq \frac{2\mu(1 + \epsilon^2)}{(1 - \epsilon)^2} \max_{\|\Delta_j\|_2 \leq 1} \text{Tr} \left((\bar{\mathbf{g}}_j + \Delta_j) (\bar{\mathbf{g}}_j + \Delta_j)^H \mathbf{W}^* \right)$$

for $j = 1, \dots, J$.

Analysis of the Randomization Procedure: General Case

- Hence, by taking

$$\epsilon = 1/2, \quad \mu = O(N \ln J), \quad \nu = \Omega(1/M)$$

and arguing as before, we conclude that $\hat{\mathbf{w}} = \boldsymbol{\xi}/\sqrt{\mu}$ is an (ν/μ) -approximate solution to (R-MMF-BF) with constant probability, assuming $\delta_1 = \cdots = \delta_J = 1$.

- Note that $\nu/\mu = \Omega(1/MN \ln J)$. Compared with the non-robust case, the ratio also degrades **linearly** in the number of transmit antennae at the ST (N).

Summary of the General Case

- Recall the problem we are trying to solve:

$$\begin{aligned} & \max_{\mathbf{w} \in \mathbb{C}^N} \min_{i=1, \dots, M} \text{SNR}_i(\mathbf{w}) \\ & \text{subject to } \|\mathbf{w}\|_2^2 \leq P, \\ & \max_{\|\Delta_j\|_2 \leq \delta_j} \text{INT}_j(\mathbf{w}, \Delta_j) \leq \beta_j \quad \text{for } j = 1, \dots, J. \end{aligned} \quad (\text{R-MMF-BF})$$

- The proposed algorithm
 - Solve the semidefinite relaxation and get the optimal solution \mathbf{W}^* .
 - Generate $\boldsymbol{\xi} \sim \mathcal{CN}(\mathbf{0}, \mathbf{W}^*)$ and scale it down by

$$\mu = \max \left\{ \frac{\|\boldsymbol{\xi}\|_2^2}{P}, \max_{j=1, \dots, J} \max_{\|\Delta_j\|_2 \leq \delta_j} \frac{\text{INT}_j(\boldsymbol{\xi}, \Delta_j)}{\beta_j} \right\}$$

to get a feasible solution $\hat{\mathbf{w}}$ to (R-MMF-BF).

- Question:** How to compute μ ?

Summary of the General Case

- The computation of μ involves computing

$$\max_{\|\Delta_j\|_2 \leq \delta_j} \text{INT}_j(\xi, \Delta_j) = \max_{\|\Delta_j\|_2 \leq \delta_j} \left| (\bar{g}_j + \Delta_j)^H \xi \right|^2,$$

which is a non-convex QCQO problem.

- Nevertheless, it is a trust region-type problem, for which there are polynomial-time algorithms **[Ye'92, Ye'94, Adachi-Iwata-Nakatsukasa-Takeda'17]**.

Final Remarks

- We analyzed the approximation quality of a semidefinite relaxation-based algorithm for solving the robust multicast beamforming problem in cognitive radio networks.
- Our techniques can be extended to develop approximation analysis of semidefinite relaxation-based algorithms for a class of **rank-constrained robust fractional SDPs**.
- There are not many results concerning the approximability of NP-hard robust optimization problems in the literature. This would be an interesting direction for future study.

Thank You!