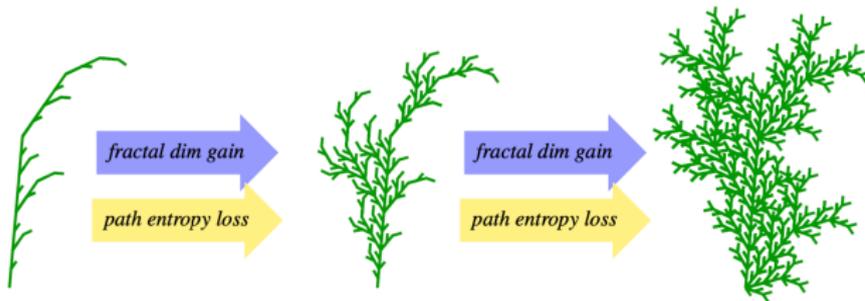


# Compactness and measure in second order arithmetic



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WKL :	Every infinite tree has a path
$P^+$ :	Every positive tree has a positive perfect subtree
P :	Every positive tree has a perfect subtree
$P^-$ :	Every positive tree has an infinite countable family of paths
WWKL :	Every positive tree has a path

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Table: Compactness principles derived by weakening weak König's lemma

Our paper: <http://arxiv.org/abs/2104.12066>

## References

- ▶ C. Chong, W. Li, W. Wang, and Y. Yang. On the computability of perfect subsets of sets with positive measure. *Proc. Amer. Math. Soc.*, 147:4021–4028, 2019.
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- ▶ L. A. Levin. Forbidden information. *J. ACM*, 60(2):9:1–9:9, 2013.
- ▶ Barmpalias and Wang: <http://arxiv.org/abs/2104.12066>

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Table: Compactness principles derived by weakening weak König's lemma

Trees of random reals are ubiquitous in measure-theoretic constructions in computability theory and, as recently suggested by Chong, Li, Wang, and Yang., essential in study of compactness in reverse mathematics.

Our goal is:

- ▶ to establish essential computational properties of the pathwise-random trees (formally defined below)
- ▶ to apply our analysis to the classification of compactness principles in second order arithmetic.

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Table: Compactness principles derived by weakening weak König's lemma

Provable in  $\text{RCA}_0$ :  $\text{WKL} \rightarrow P^+ \rightarrow P \rightarrow P^- \rightarrow \text{WWKL}$

### Theorem (Models)

Each of the following extensions of  $\text{RCA}$  has an  $\omega$ -model:

- $\text{WWKL} + \neg P^-$ : every positive tree has a path but some positive tree only has finitely many paths.
- $P + \neg P^+$ : every positive tree has a perfect subtree but some positive tree has no positive perfect subtree.
- $P^+ + \neg \text{WKL}$ : every positive tree has a positive perfect subtree but some infinite tree has no path.

## Definition

The deficiency of  $\sigma$  is  $|\sigma| - K(\sigma)$ ;

the deficiency of a real  $x$  is the supremum of the deficiencies of  $x \upharpoonright_n$ ,  $n \in \mathbb{N}$ .

The deficiency of a set of reals is the supremum of the deficiencies of its members.

## Definition (Pathwise randomness)

A tree  $T$  is:

- ▶ pathwise-random if the deficiency of each  $\sigma \in T$  is bounded above by a constant.
- ▶ weakly pathwise-random if all of its paths are random.
- ▶ proper if it has infinitely many paths.

## Theorem

If  $z$  is random and computes or enumerates a pathwise-random tree of unbounded width, then  $z \geq_T \emptyset'$ .

Hirschfeldt, Jockusch, and Schupp (2021) obtained a similar statement for perfect trees and 2-randoms.

## Corollary

The set of paths through the van Lambalgen array:

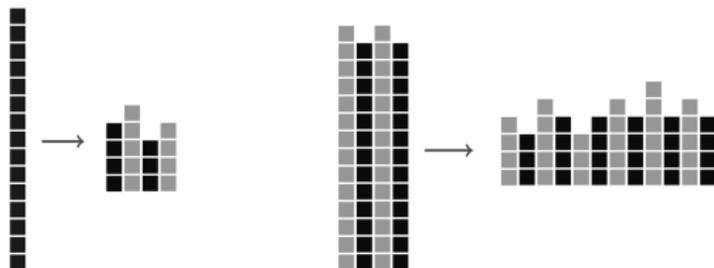
$$A_x := \{ x_{1^n * 0} : n \in \mathbb{N} \}$$

of any  $x \not\geq_T \emptyset'$  has infinite randomness deficiency.

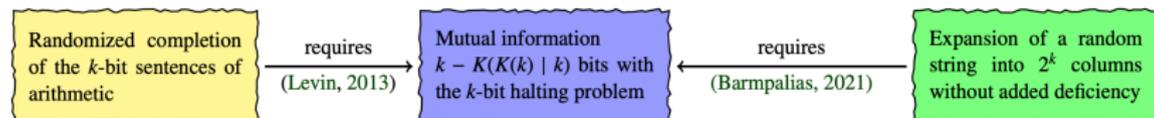
## Corollary

If  $F$  is a computable space of trees, for every computable measure  $\nu$  on  $F$  the class of proper pruned pathwise-random members of  $F$  is  $\nu$ -null.

# Finitary consequences



- Effectively splitting a random source into  $k$  many random sources without a significant increase in the randomness deficiency is about as hard as computing the  $k$ -bit halting problem.
- Levin (2013): randomly guessing a completion of the  $k$ -bit segment of PA is about as improbable as randomly guessing the  $k$ -bit halting problem.



## Digression on the proof of WWKL $\not\rightarrow$ P<sup>-</sup> (algorithms)

Proof base:

Randomness cannot be used in order to produce infinitely many random reals with a fixed upper bound on their deficiency.

We explore the finite version of this fact and its limits.

**Theorem** (Random production of incompressible strings I)

There exists no randomized algorithm and  $\epsilon > 0, c \in \mathbb{N}$  such that for each input  $k$ , with probability  $> \epsilon$  the output is a set of  $k$  strings  $\rho$  with  $K(\rho) > |\rho| - c$  of the same length  $\ell_k$  which depends only on  $k$ .

**Theorem** (Random production of incompressible strings II)

There exists a randomized algorithm which, almost surely and on almost all inputs  $k$ , outputs a set of  $k$  strings  $\rho$  of equal length  $\ell_k$  with  $K(\rho) > |\rho|$ , where  $\ell_k \in (2^k, 2^{k+1})$  is a random variable.

$P \not\rightarrow P^+$

## Theorem

The following hold for each  $z$ :

- (a) there exists a perfect pathwise-random tree  $T$  such that no  $T$ -c.e. positive tree is pathwise-random.
- (b) if no  $z$ -c.e. positive tree is pathwise-random, then there exists a perfect pathwise- $z$ -random tree  $T$  such that no  $(z \oplus T)$ -c.e. positive tree is pathwise-random.
- (c) the tree  $T$  of clause (b) can be found inside any given positive tree.

Forcing with sets of sets of positive measure.

Involving hitting sets and notions from Poisson point processes.

$P^+ \not\rightarrow WKL$

Patey has shown that every positive tree contains a perfect subtree which does not compute any PA degree.

Theorem (Patey)

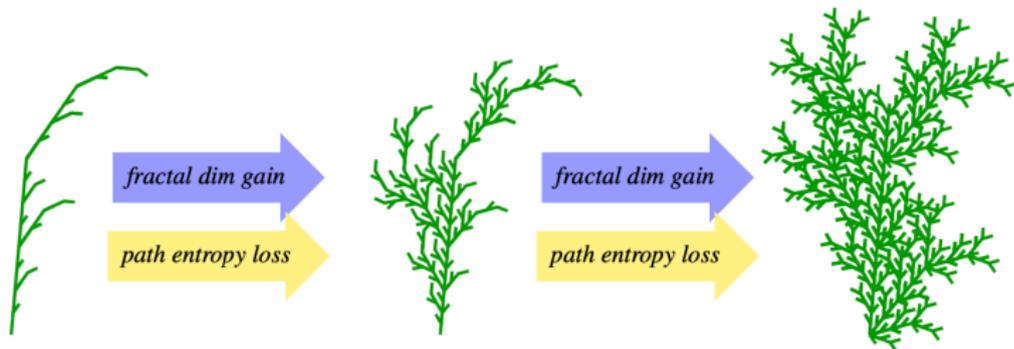
Every positive tree contains a perfect subtree which does not compute a complete extension of PA.

Theorem

There exists a positive perfect pathwise-random tree which does not compute any complete extension of Peano Arithmetic. In fact, given any non-computable  $z$ , every positive tree has a positive perfect subtree  $T \not\leq_T z$  which does not compute any complete extension of Peano Arithmetic.

Independently obtained by Greenberg, Miller, Nies (2021).

# Open problems



- ▶ build a model for  $P^- + \neg P$
- ▶ build measures  $\nu$  such that  $\nu$ -randomness gives the above separations.

Thanks for listening!