Transmit Strategies for the Gaussian Bidirectional Broadcast Channel & Latest Results

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Bidirectional Broadcast Channel



 $R_1 \le I(X_R; Y_1)$ $R_2 \le I(X_R; Y_2)$

Practically relevant since

- supports modularization
- gains are easily realized



(Some) Related Literature

• Early results: Decode & forward strategies based on

- superposition coding [Rankov et al. '05], [Oechtering et al.,'06], et al.
- XOR operation [Larsson et al. '04], [Wu et al. '04], [Yeung '05], et. al.
- optimal channel coding approach based on network coding idea (single information flow) found by many groups independently [Knopp '06], [Oechtering et al. '07], [Kim et al. '07], [Xie '07], [Wu '07]
- Closely related problems:
 - Common message BC (multicast) among others [Khisti '04]
 - Compound channel [Blackwell et al '59], [Wolfowitz '60], et al.
 - Slepian-Wolf coding over BC [Tuncel '06]
 - Physical-layer NC [Zhang et al. '06], [Popovski et al. '06], et al.
- Extensions: Compress or compute & forward strategies
 - [Schnurr et al '07], [Kim et al, '08], [Günduz et al.'08], [Wilson et al, '08], [Nam et al. '08], [Nazer et al. '08], [Ong et al. '10], [Lim et al. '10], et al.

Gaussian Multi-Antenna Bidirectional Relaying



Capacity Region of Bidirectional Broadcast Channel

$$C_{\rm BC} := \bigcup_{\text{tr} \, Q \le P, \, Q \ge 0} \left\{ [R_1, R_2] \in \mathbb{R}^2_+ : R_1 \le C_1(Q), R_2 \le C_2(Q) \right\}$$

with

$$C_i(\mathbf{Q}) := \log \det \left(\mathbf{I}_{N_i} + \frac{1}{\sigma^2} \mathbf{H}_i^H \mathbf{Q} \mathbf{H}_i \right), \quad i = 1, 2.$$

Transmit Covariance Optimization Problem

$$\arg \max_{\operatorname{tr} Q \leq P, Q \geq 0} \sum_{i=1}^{2} w_i \log \det \left(I_{N_i} + \frac{1}{\sigma^2} H_i^H Q H_i \right)$$

$$\underset{\operatorname{tr} Q \leq P, Q \geq 0}{\operatorname{tr} Q \leq P, Q \geq 0} \sum_{i=1}^{2} w_i \log \det \left(I_{N_i} + \frac{1}{\sigma^2} H_i^H Q H_i \right)$$

Optimal Transmit Strategies for the MISO case

- Optimal Transmit Strategies for the MIMO case
- Latest Results and Conclusion

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First Study – MISO Case $N_1 = N_2 = 1$

MISO optimization problem

$$Q_{\text{opt}}(w) = \underset{\text{tr} Q \leq P, Q \geq 0}{\operatorname{arg max}} \sum_{i=1}^{2} w_i \log \left(1 + \frac{1}{\sigma^2} \boldsymbol{h}_i^H \boldsymbol{Q} \boldsymbol{h}_i\right)$$

- Outline:
 - Subspace optimality and orthogonal channels
 - Single-beam optimality and its consequences
 - Optimal beamforming vector
- Results are published in [Trans SP '09].

First Observations

Proposition: Subspace optimality

An optimal transmit strategy transmits only into the subspace spanned by the channels, otherwise transmit power can be reduced while achieving the same rates.

••• Optimal transmit strategy Q_{opt} has always $rank(Q_{opt}) \le 2!$

Proposition: Orthogonal channels

For orthogonal channels any rate pair can be achieved with a single-beam as well as with a two-beam strategy.

For orthogonal channels the capacity region can be also achieved using the superposition encoding strategy.

Optimality of the Single-Beam Strategy

Theorem: Single-beam optimality

For the MISO case we can always find an optimal single-beam transmit strategy $(rank(Q_{opt}) = 1)$.

Proof outline:

- Orthogonal channels: Optimality follows immediately from previous propositions.
- *Non-orthogonal channels:* Any rank-two transmit strategy contradicts with the Karush-Kuhn-Tucker conditions so that the optimal transmit strategy has to have rank one.

Optimal strategy is to perform a single beam onto the subspace spanned by the channels!

Consequences of Single-Beam Optimality

For the bidirectional broadcast channel ...

Signal Processing

... the relay forms a single beam instead of individual beams for each user as for the classical MISO broadcast.

Correlated channels will be beneficial (result not shown).

Channel Coding

... it is sufficient to use an one-dimensional Gaussian codebook instead of a codebook with a dimension equal to the number of transmit antennas.



Optimal Beamforming Vector

Theorem: Property of Optimal Beamforming Vector

$$\boldsymbol{Q} := \boldsymbol{P} \boldsymbol{q} \boldsymbol{q}^{\boldsymbol{H}}, \qquad \boldsymbol{q} = a_1 \boldsymbol{u}_1 + a_2 \boldsymbol{u}_2, \qquad \boldsymbol{h}_i = |\boldsymbol{h}_i| \boldsymbol{u}_i, \ a_i \in \mathbb{C},$$

then

$$\arg(a_1) - \arg(a_2) = \varphi \qquad |\rho|e^{i\varphi} = u_1^H u_2$$

Normalized beamforming vector

$$q(t) = \frac{tu_1 + (1-t)e^{-j\varphi}u_2}{\|tu_1 + (1-t)e^{-j\varphi}u_2\|}, \qquad t \in [0,1]$$

- $[R_1(t), R_2(t)]$ with $R_i(t) := \log(1 + \frac{p}{\sigma^2}|h_k^H q(t)|^2), t \in [0, 1]$ parametrizes the curved section of the capacity region!
- Egalitarian solution easily calculated from $R_1(t_{eg}) = R_2(t_{eg})$.

Extended Study – MIMO Case

MIMO optimization problem

$$\mathbf{Q}_{\mathsf{opt}}(\boldsymbol{w}) = \underset{\operatorname{tr} \boldsymbol{Q} \leq \boldsymbol{P}, \boldsymbol{Q} \geq \boldsymbol{0}}{\operatorname{arg\,max}} \sum_{i=1}^{2} w_{i} \log \det \left(\mathbf{I}_{N_{i}} + \frac{1}{\sigma^{2}} \mathbf{H}_{i}^{H} \mathbf{Q} \mathbf{H}_{i} \right)$$

- Outline:
 - Subspace optimality and 'Orthogonal' channels
 - Karush-Kuhn-Tucker conditions Unsymmetric Riccati equation
 - Special case: Full rank transmission
 - Special case: Parallel channels
- Results are published in [Trans Com '09].

First Observations

Proposition: Subspace optimality

An optimal transmit strategy transmits only into the vector space spanned by the set of column vectors of H_1 and H_2 .

Proposition: 'Orthogonal channels'

 P_i projector onto the vector space spanned by the set of column vectors of H_i , i = 1, 2.

Any rate pair achievable with Q can be achieved with equivalent transmit strategies \hat{Q} with rank \hat{r} satisfying

 $\max\{r_1, r_2\} \le \hat{r} \le \min\{r_1 + r_2, N_R\}, \quad r_i := \operatorname{rank}(P_i Q P_i).$

• In general an optimal solution Q_{opt} will be not unique!

Makes analysis of the general optimization problem difficult.

Lagrangian

$$L(\mathbf{Q}, \mu, \Psi) = -\sum_{i=1}^{2} w_i C_i(\mathbf{Q}) - \mu \left(P - \operatorname{tr} \mathbf{Q} \right) - \operatorname{tr} \mathbf{Q} \Psi$$

Karush-Kuhn-Tucker conditions

$$\sum_{i=1}^{2} w_{i} H_{i} (\sigma^{2} I_{N_{i}} + H_{i}^{H} Q H_{i})^{-1} H_{i}^{H} = \mu I_{N_{R}} - \Psi$$
(1)

$$Q \geq 0, \qquad P \geq \operatorname{tr} Q,$$

$$\Psi \geq 0, \qquad \mu \geq 0,$$

$$\operatorname{tr} Q \Psi = 0, \qquad \mu (P - \operatorname{tr} Q) = 0$$

• If H_i^{-1} exists, then (1) can be expressed as...

Lagrangian

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Karush-Kuhn-Tucker conditions

$$\sum_{i=1}^{2} w_{i} (\sigma^{2} \boldsymbol{H}_{i}^{-H} \boldsymbol{H}_{i}^{-1} + \boldsymbol{Q})^{-1} = \mu \boldsymbol{I}_{N_{R}} - \boldsymbol{\Psi}$$
$$\boldsymbol{Q} \geq 0, \qquad P \geq \operatorname{tr} \boldsymbol{Q},$$
$$\boldsymbol{\Psi} \geq 0, \qquad \mu \geq 0,$$
$$\operatorname{tr} \boldsymbol{Q} \boldsymbol{\Psi} = 0, \qquad \mu (P - \operatorname{tr} \boldsymbol{Q}) = 0$$

• with substitutions $A_i := \sigma^2 H_i^{-H} H_i^{-1}$ and $B := \mu I_{N_R} - \Psi \dots$

Lagrangian

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Karush-Kuhn-Tucker conditions

$$w_1(A_1 + Q)^{-1} + w_2(A_2 + Q)^{-1} = B$$

$$Q \ge 0, \qquad P \ge \operatorname{tr} Q,$$

$$\Psi \ge 0, \qquad \mu \ge 0,$$

$$\operatorname{tr} Q\Psi = 0, \qquad \mu (P - \operatorname{tr} Q) = 0$$

• multiplication with $(A_1 + Q)$ and $(A_2 + Q)$ we get...

Lagrangian

$$L(\mathbf{Q}, \boldsymbol{\mu}, \boldsymbol{\Psi}) = -\sum_{i=1}^{2} w_i C_i(\mathbf{Q}) - \boldsymbol{\mu} (P - \operatorname{tr} \mathbf{Q}) - \operatorname{tr} \mathbf{Q} \boldsymbol{\Psi}$$

Karush-Kuhn-Tucker conditions

 $QBQ + QBA_2 + A_1BQ - Q = w_1A_2 + w_2A_1 - A_1BA_2,$ $Q \ge 0, \qquad P \ge \operatorname{tr} Q,$ $\Psi \ge 0, \qquad \mu \ge 0,$ $\operatorname{tr} Q\Psi = 0, \qquad \mu (P - \operatorname{tr} Q) = 0$

 ... a quadratic matrix equation (also known as unsymmetric Riccati equation). A solution method exists, but further analytical results are not available so far (we do not know).

Lagrangian

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Karush-Kuhn-Tucker conditions

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$$Q \ge 0, \qquad P \ge \operatorname{tr} Q,$$

$$\Psi \ge 0, \qquad \mu \ge 0,$$

$$\operatorname{tr} Q\Psi = 0, \qquad \mu (P - \operatorname{tr} Q) = 0$$

$$(2)$$

Notice, at this step we can multiply (A₁ + Q) and (A₂ + Q) from the left or from the right!

Special Case: ... and Full Rank Transmission

Full Rank Transmission and Invertible Channels

Assume: H_i^{-1} exists & the optimal covariance matrix has rank Q = N

$$\Psi = \mathbf{0}$$
 and therefore $B = \mu I_N$

• Interchanging multiplications from the left and the right ... $w_1(A_2 + Q) + w_2(A_1 + Q) = \mu(A_1 + Q)(A_2 + Q) = \mu(A_2 + Q)(A_1 + Q)$ shows that matrices $(A_1 + Q)$ and $(A_2 + Q)$ commute!

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Both have the same eigenspace, i.e., $(A_i + Q) = U\Sigma_i U^H$, i = 1, 2, which can be computed from

$$(A_2 + Q) - (A_1 + Q) = \underbrace{A_2 - A_1}_{=\sigma^2 ((H_2 H_2^H)^{-1} - (H_1 H_1^H)^{-1})} = U(\Sigma_2 - \Sigma_1) U^H$$

Optimal Eigenvalues

Proposition: Quadratic Equation Condition

 $\boldsymbol{U} = [\boldsymbol{u}_1, \boldsymbol{u}_2, \dots, \boldsymbol{u}_N]$ diagonalizes matrix equation

$$w_1(\delta_{2,k} + \epsilon_k) + w_2(\delta_{1,k} + \epsilon_k) = \mu(\delta_{1,k} + \epsilon_k)(\delta_{2,k} + \epsilon_k)$$

linear in ϵ_k

quadratic in ϵ_k

with
$$\delta_{i,k} = \boldsymbol{u}_k^H \boldsymbol{A}_i \boldsymbol{u}_k > 0$$
 and $\boldsymbol{\epsilon}_k = \boldsymbol{u}_k^H \boldsymbol{Q} \boldsymbol{u}_k > 0, k = 1, 2, \dots, N.$



- **Solution:** Intersection between line and parabola. (negative sol. can be excluded)
- Choose μ such that $\sum_{k=1}^{N} \epsilon_k = P$

Optimal Transmit Covariance Matrix

Optimal Transmit Covariance Matrix

The optimal transmit covariance for the case of invertible channels and a full rank transmission is given by

 $\boldsymbol{Q} = \boldsymbol{U} \text{diag} \left[\delta_{i,1} + \epsilon_1, \delta_{i,2} + \epsilon_2, \dots, \delta_{i,N} + \epsilon_N \right] \boldsymbol{U}^H - \boldsymbol{A}_i, \quad i = 1, 2,$

which can be completely calculated by the previous procedure.

• Possible extension to case where Q^{-1} , $(H_1H_1^H)^{-1}$, and $(H_2H_2^H)^{-1}$ exist (Sherman-Morrison-Woodbury formula).

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Open Problem: Generalization

Non-full rank transmission & channels with different subspaces **Problem:** $\Psi \neq 0$ **Potential difficulty:** Optimal solution may be not unique!

Complete Solution for Special Case: Parallel Channels

Definition: Parallel Channels (e.g. OFDM)

$$\boldsymbol{H}_{1}\boldsymbol{H}_{1}^{H} = \boldsymbol{W}\boldsymbol{S}_{1}\boldsymbol{W}^{H} \qquad \boldsymbol{H}_{2}\boldsymbol{H}_{2}^{H} = \boldsymbol{W}\boldsymbol{S}_{2}\boldsymbol{W}^{H}$$

with $S_i = \text{diag}(s_{i,1}, s_{i,2}, ..., s_{i,N}) \ge 0$, i = 1, 2, and W unitary.

Optimal eigenvectors

Hadamard Inequality: $\mathbf{Q} = \mathbf{W} \mathbf{\Sigma}_{Q} \mathbf{W}^{H}, \mathbf{\Sigma}_{Q} = \text{diag}(\lambda_{1}, \lambda_{2}, \dots, \lambda_{N})$

Optimal eigenvalues

Weighted rate sum maximization:

$$R_{\Sigma}(\boldsymbol{w}) = \max_{\lambda} \sum_{k=1}^{N} \sum_{i=1}^{2} w_i \log(1 + \frac{1}{\sigma^2} s_{i,k} \lambda_k) \text{ s.t. } \|\boldsymbol{\lambda}\|_1 \le P, \ \lambda_k \ge 0$$

• Previous procedure solves even non-full rank case.

BiBC Under Channel Uncertainty

- Channel uncertainty is a ubiquitous phenomenon in practical systmes
 - Assume that it is only known that the exact channel realization belongs to a pre-specified set of channels *S*
 - We need universal strategies that work for all realizations simultaneously

Theorem: Capacity Region of Compound BiBC [TCom'10]

$$R_1 \leq \inf_{s \in \mathcal{S}} I(X_R; Y_{1,s})$$
 and $R_2 \leq \inf_{s \in \mathcal{S}} I(X_R; Y_{2,s})$

 $Y_{i,s}$ channel output at node *i* for channel realization $s \in S$.

• Results are extended to arbitrary varying channels.

Robust Transmit Strategies for the MISO case

• CSI uncertainty: h_{i,0} nominal (known) channel, [Loyka et al.'08]

$$y_i = (h_{i,0} + d_i)x + n_i, \qquad i = 1, 2$$

perturbation $d_i \in \mathcal{D}_i$

$$\mathcal{D}_i := \{ \boldsymbol{d}_i : \sigma_1(\boldsymbol{d}_i) = \| \boldsymbol{d}_i \| \le \epsilon_i \}$$

Optimal Robust Transmit Strategy

 Worst case capacity region of the MISO BiBC under channel uncertainty D_i is given by set of rate pairs

$$R_1 \leq \log\left(1 + \frac{1}{\sigma^2}(|\boldsymbol{h}_{1,0}^H\boldsymbol{q}| - \epsilon_1)^2\right), \quad R_2 \leq \log\left(1 + \frac{1}{\sigma^2}(|\boldsymbol{h}_{2,0}^H\boldsymbol{q}| - \epsilon_2)^2\right)$$

for some transmit strategy $Q = qq^H$ with $tr(Q) \le P$.

Worst-Case Perturbation



Worst-case perturbations can explicitly be characterized as

$$d_i(q) = -\epsilon_i e^{-j\varphi_i} u_q$$
 with $u_q = q/||q||$ and $\varphi_i = arg(h_{i,0}^H u_q)$

Worst-case d_i are anti-parallel to transmit strategy q

• Some extensions to MIMO case possible.

Conclusion

Bidirectional broadcast channel is an appealing problem which

- is practically relevant
 - modularization, realizes some network coding gains,
- is closely related to multicast, P2P channel,
- allows derivation of closed form results in the Gaussian case
- MISO: single-beam optimaity manifests single information flow view, favors correlated channels
- MIMO: closed-form procedure to find full rank solution \rightarrow extension open,
- has many interesting extensions, e.g.,
 - compound and arbitrary varying channel versions
 - \rightarrow robust transmit strategies.

Thank you for your attention!

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Optimal eigenvalues

• Lagrangian:
$$L(\lambda, \mu, \nu) = \mu(P - ||\lambda||_1) + \sum_{k=1}^{N} \left[\nu_k \lambda_k - \sum_{i=1}^{2} w_i \log(1 + \frac{1}{\sigma^2} s_{i,k} \lambda_k) \right],$$

Optimal λ_k follows from the Karush-Kuhn-Tucker conditions

 $w_1s_{1,k}(\sigma^2 + s_{2,k}\lambda_k) + w_2s_{2,k}(\sigma^2 + s_{1,k}\lambda_k) = (\mu - \nu_k)(\sigma^2 + s_{1,k}\lambda_k)(\sigma^2 + s_{2,k}\lambda_k)$

- Quadratic equation in λ_k ; negative solution can be excluded
- µ has to be chosen so that power constraint is fulfilled
- v_k controls positivity condition of λ_k

The case $\nu_k = \lambda_k = 0$ defines a thresholds $\mu_k(w)$ where mode *k* is activated, i.e., for larger $P \Rightarrow$ smaller μ , we will have $\lambda_k > 0$,

$$\mu_k(w) = \sigma^{-2}(w_1s_{1,k} + w_2s_{2,k}).$$

Modes Areas of Parallel Channels with $N_1 = N_2 = N_R = 3$



- Eigenvalue $s_{i,k}$ corresponds to the k-th eigenvector of $H_i H_i^H$
- The activated modes, $\lambda_k > 0$, change with the weights. At the weights corresponding to 'o' beamforming is never optimal.