

Principal component analysis for the approximation of  
high-dimensional functions in tree-based tensor formats

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Supported by the ANR (CHORUS project)

We consider the problem of constructing an **approximation** of a random variable  $Y$  by a function of a set of random variables  $X = (X_1, \dots, X_d)$ , **using samples** of  $(X, Y)$ , when

- the samples of  $(X, Y)$  can be generated from **adaptively chosen samples** of  $X$  (**active learning**),
- there exists a **deterministic function**  $u$  such that

$$Y = u(X).$$

In practice,  $Y$  could be the output of a numerical model (computer code) and  $X$  a set of input parameters modelling uncertainties on the model, with known probability distribution.

The approximation can then be used as a **predictive surrogate model**.

When the generation of one sample requires a costly numerical simulation (or experiment), only a few samples are available.

**Low-dimensional structures of functions have to be exploited** (e.g., low effective dimensionality, anisotropy, sparsity, **low rank**).

- 1 Tree-based tensor formats
- 2 Principal component analysis in tree-based tensor formats

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# Tensor spaces of multivariate functions

Assume  $X_\nu$  has probability law  $\mu_\nu$  with support  $\mathcal{X}_\nu$ .

Let  $\mathcal{H}_\nu$  a Hilbert space of functions defined on  $\mathcal{X}_\nu$ , typically  $L^2_{\mu_\nu}(\mathcal{X}_\nu)$  or a reproducing kernel Hilbert space (RKHS) in  $L^2_{\mu_\nu}(\mathcal{X}_\nu)$ .

We consider multivariate functions defined on  $\mathcal{X} = \mathcal{X}_1 \times \dots \times \mathcal{X}_d$  that are elements of the **tensor Hilbert space**

$$\mathcal{H}_1 \otimes \dots \otimes \mathcal{H}_d := \mathcal{H}$$

equipped with the **canonical norm**.

## $\alpha$ -rank of higher-order tensors

For a non-empty subset  $\alpha$  of  $D = \{1, \dots, d\}$  and  $\alpha^c = D \setminus \alpha$ , a tensor  $u \in \mathcal{H}$  can be identified with an order-two tensor

$$\mathcal{M}_\alpha(u) \in \mathcal{H}_\alpha \otimes \mathcal{H}_{\alpha^c},$$

where

$$\mathcal{H}_\beta = \bigotimes_{\nu \in \beta} \mathcal{H}_\nu \subset \mathbb{R}^{\mathcal{X}_\beta}.$$

The  $\alpha$ -rank of  $u$  is defined by

$$\text{rank}_\alpha(u) = \text{rank}(\mathcal{M}_\alpha(u)),$$

which is the minimal integer  $r_\alpha$  such that

$$u(x) = \sum_{k=1}^{r_\alpha} v_k^\alpha(x_\alpha) w_k^{\alpha^c}(x_{\alpha^c})$$

## $\alpha$ -ranks and related tensor formats

- For  $T$  a collection of subsets of  $D$ , we define the  $T$ -rank of a tensor  $v$  as the tuple

$$\text{rank}_T(v) = \{\text{rank}_\alpha(v)\}_{\alpha \in T}.$$



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- The set of tensors with  $T$ -rank bounded by  $r = (r_\alpha)_{\alpha \in T}$  is

$$\mathcal{T}_r^T = \{v \in \mathcal{H} : \text{rank}_T(v) \leq r\} = \bigcap_{\alpha \in T} \{v \in \mathcal{H} : \text{rank}_\alpha(v) \leq r_\alpha\},$$

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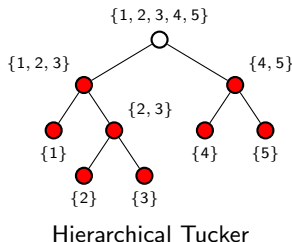
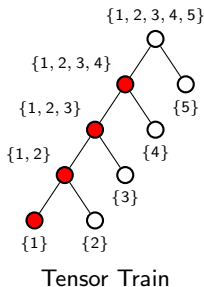
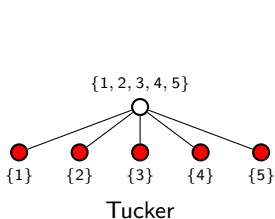
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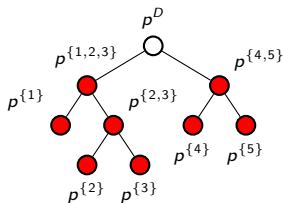
and is called a **tensor format**.

- Tree-based tensor formats correspond to a **tree-structured  $T$**



# Tree-based tensor formats

- A tensor in  $\mathcal{T}_r^T$  admits a **multilinear parametrization** with parameters  $\{p^\alpha\}_{\alpha \in \text{TU}\{D\}}$  forming a **tree network of low order tensors**.



- **Storage complexity** scales as  $O(dR^{s+1})$  where  $R$  is the maximal  $\alpha$ -rank and  $s$  is the arity of the tree.
- Corresponds to a deep (sum-product) network with depth bounded by  $d - 1$ .
- $\mathcal{T}_r^T$  is weakly closed (and therefore proximal).

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- 2 Principal component analysis in tree-based tensor formats

# Principal subspaces

For a subset of variables  $\alpha$ , a multivariate function  $u(x_1, \dots, x_d)$  is identified with a bivariate function  $u \in \mathcal{H}_\alpha \otimes \mathcal{H}_{\alpha^c}$  which admits a singular value decomposition

$$u(x_\alpha, x_{\alpha^c}) = \sum_{k=1}^{\text{rank}_\alpha(u)} \sigma_k^\alpha v_k^\alpha(x_\alpha) v_k^{\alpha^c}(x_{\alpha^c})$$

The subspace of  $\alpha$ -principal components

$$U_\alpha = \text{span}\{v_1^\alpha, \dots, v_{r_\alpha}^\alpha\}$$

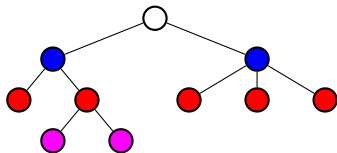
is solution of

$$\min_{\dim(U_\alpha)=r_\alpha} \|u - \mathcal{P}_{U_\alpha} u\| = \min_{\text{rank}_\alpha(v) \leq r_\alpha} \|u - v\|$$

where  $\mathcal{P}_{U_\alpha} = P_{U_\alpha} \otimes id_{\alpha^c}$  is the orthogonal projection onto  $U_\alpha \otimes \mathcal{H}_{\alpha^c}$ .

# Higher-order principal component analysis for tree-based formats

Let  $T$  be a tree-structured collection of subsets of  $2^D$



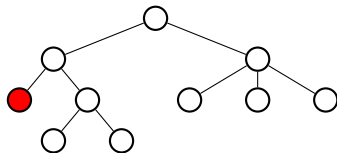
For all  $\alpha$  in  $T$ , we will determine subspaces  $U_\alpha$  that are **approximations of  $\alpha$ -principal subspaces** of  $u$  in **low-dimensional subspaces  $V_\alpha$**  of functions defined on  $\mathcal{X}_\alpha$ .

# Higher-order principal component analysis for tree-based formats

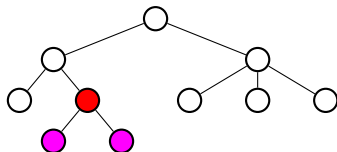
For each  $\alpha \in \mathcal{T}$ ,  $U_\alpha$  is defined as the  $r_\alpha$ -dimensional  $\alpha$ -principal subspace of an approximation of  $u$

$$u_\alpha = \mathcal{P}_{V_\alpha} u$$

- for  $S(\alpha) = \emptyset$  (**leaf node**),  $V_\alpha$  is a given approximation space in  $\mathcal{H}_\alpha$  (e.g., polynomials, wavelets, ...),



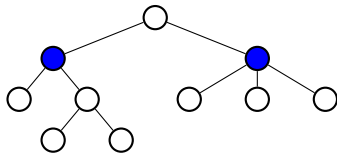
- for  $S(\alpha) \neq \emptyset$  (**interior node**),  $V_\alpha = \bigotimes_{\beta \in S(\alpha)} U_\beta$ .



# Higher-order principal component analysis for tree-based formats

We finally obtain an approximation  $u^*$  by projecting  $u$  onto the tensor space  $\bigotimes_{\alpha \in S(D)} U_\alpha$

$$u^* = \prod_{\alpha \in S(D)} \mathcal{P}_{U_\alpha} u$$





# Learning algorithm based on principal component analysis

For a feasible algorithm using samples,

- 1 orthogonal projections  $\mathcal{P}_{W_\alpha}$  on subspaces  $W_\alpha$  ( $U_\alpha$  or  $V_\alpha$ ) are replaced by **oblique projections  $\mathcal{I}_{W_\alpha}$  using samples** (e.g. **interpolation** or least-squares projection),
- 2 principal subspaces  $U_\alpha$  of  $u_\alpha = \mathcal{I}_{V_\alpha} u$  are **estimated using samples** of the  $V_\alpha$ -valued random variable

$$u_\alpha(\cdot, X_{\alpha^c})$$

With interpolation, this requires the evaluation of  $u$  at the  $\dim(V_\alpha) \times N_\alpha$  points

$$\{(x_\alpha, x_{\alpha^c}^k) : x_\alpha \in \Gamma_{V_\alpha}, 1 \leq k \leq N_\alpha\}$$

where  $\Gamma_{V_\alpha} \subset \mathcal{X}_\alpha$  is a unisolvent set of points for  $V_\alpha$  (magic points), and the  $x_{\alpha^c}^k$  are i.i.d. samples of  $X_{\alpha^c}$ .

# Properties of the algorithm

## Theorem (Prescribed rank)

For a given  $T$ -rank, if the subspaces  $U_\alpha$  are such that

$$\|\mathcal{P}_{U_\alpha} u_\alpha - u_\alpha\| \leq C \min_{\text{rank}_\alpha(v) \leq r_\alpha} \|v - u_\alpha\|$$

holds with probability higher than  $1 - \eta$ , then we obtain an approximation  $u^*$  such that

$$\|u^* - u\|^2 \leq \Lambda^2 C^2 \#T \min_{v \in \mathcal{T}_r^T} \|v - u\|^2 + \tilde{\Lambda}^2 \max_{1 \leq \nu \leq d} \|u - \mathcal{P}_{V_\nu} u\|^2$$

holds with probability higher than  $1 - \eta \#T$ , with  $\Lambda$  and  $\tilde{\Lambda}$  depending on the properties of the oblique projection operators.

**About complexity:** If  $N_\alpha = r_\alpha$  for all  $\alpha \in T$ , then the total number of evaluations  $N$  is equal to the storage complexity  $S$  of the resulting approximation  $u^* \in \mathcal{T}_r^T$ .

# Properties of the algorithm

## Theorem (Fixed precision)

Let  $\epsilon, \tilde{\epsilon} \geq 0$ . If the subspaces  $U_\alpha$  are determined such that

$$\|\mathcal{P}_{U_\alpha} u_\alpha - u_\alpha\| \leq \frac{\epsilon}{\sqrt{\#T}} \|u_\alpha\|$$

holds with probability higher than  $1 - \eta$ , and if the approximation spaces  $V_\nu$ ,  $1 \leq \nu \leq d$ , are such that

$$\|\mathcal{P}_{V_\nu} u - u\| \leq \tilde{\epsilon} \|u\|,$$

then we obtain an approximation  $u^*$  such that

$$\|u^* - u\|^2 \leq (\Lambda^2 \epsilon^2 + \tilde{\Lambda}^2 \tilde{\epsilon}^2) \|u\|^2$$

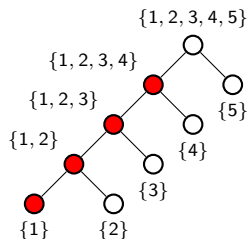
holds with probability higher than  $1 - \eta\#T$ , with  $\Lambda$  and  $\tilde{\Lambda}$  depending on the properties of the oblique projection operators.

## Illustration of tensor recovery: Henon-Heiles potential

$$u(X) = \frac{1}{2} \sum_{i=1}^d X_i^2 + 0.2 \sum_{i=1}^{d-1} (X_i X_{i+1}^2 - X_i^3) + \frac{0.2^2}{16} \sum_{i=1}^{d-1} (X_i^2 + X_{i+1}^2)^2, \quad X_i \sim U(-1, 1),$$

$\text{rank}_\alpha(u) = 3$  for all  $\alpha$  in

$$T = \{\{1\}, \{1, 2\}, \dots, \{1, \dots, d-1\}\}$$



Then  $u$  can be exactly represented in the **tensor train format**  $\mathcal{T}_r^T$  with  $T$ -rank  $r = (3, \dots, 3)$

$$u = \sum_{k_1=1}^3 \sum_{k_2=1}^3 \dots \sum_{k_{d-1}=1}^3 v_{1,k_1}^{(1)}(x_1) v_{k_1,k_2}^{(1,2)}(x_2) v_{k_2,k_3}^{(1,2,3)}(x_3) \dots v_{k_{d-1},1}^{(1,\dots,d)}(x_d)$$

with **univariate polynomial functions** of degree 4.

## Illustration of tensor recovery: Henon-Heiles potential

Table: Approximation with prescribed  $T$ -rank  $r = (3, \dots, 3)$  and polynomial degree 4 for different values of  $d$  and  $\gamma = N_\alpha/r_\alpha$ .

| $\gamma = 1$                      |              |             |             |             |               |
|-----------------------------------|--------------|-------------|-------------|-------------|---------------|
| $d$                               | 5            | 10          | 20          | 50          | 100           |
| $\varepsilon(u^*) \times 10^{14}$ | [1.0; 234.2] | [1.5; 67.5] | [2.5; 79.9] | [6.6; 62.8] | [15.7; 175.1] |
| $S = N$                           | 165          | 390         | 840         | 2190        | 4440          |

| $\gamma = 10$                     |            |            |            |            |            |
|-----------------------------------|------------|------------|------------|------------|------------|
| $d$                               | 5          | 10         | 20         | 50         | 100        |
| $\varepsilon(u^*) \times 10^{14}$ | [0.1; 0.4] | [0.2; 0.4] | [0.3; 0.4] | [0.4; 0.7] | [0.6; 0.8] |
| $S$                               | 165        | 390        | 840        | 2190       | 4440       |
| $N$                               | 1515       | 3765       | 8265       | 21765      | 44265      |

## Illustration for approximation: Borehole function

The Borehole function models water flow through a borehole:

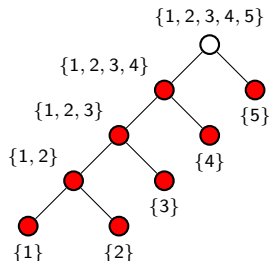
$$u(X) = \frac{2\pi T_u(H_u - H_l)}{\ln(r/r_w) \left( 1 + \frac{2LT_u}{\ln(r/r_w)r_w^2 K_w} + \frac{T_u}{T_l} \right)}, \quad X = (r_w, \log(r), T_u, H_u, T_l, H_l, L, K_w)$$

|       |  |                                     |
|-------|--|-------------------------------------|
| $r_w$ | radius of borehole (m)                               | $N(\mu = 0.10, \sigma = 0.0161812)$ |
| $r$   | radius of influence (m)                              | $LN(\mu = 7.71, \sigma = 1.0056)$   |
| $T_u$ | transmissivity of upper aquifer (m <sup>2</sup> /yr) | $U(63070, 115600)$                  |
| $H_u$ | potentiometric head of upper aquifer (m)             | $U(990, 1110)$                      |
| $T_l$ | transmissivity of lower aquifer (m <sup>2</sup> /yr) | $U(63.1, 116)$                      |
| $H_l$ | potentiometric head of lower aquifer (m)             | $U(700, 820)$                       |
| $L$   | length of borehole (m)                               | $U(1120, 1680)$                     |
| $K_w$ | hydraulic conductivity of borehole (m/yr)            | $U(9855, 12045)$                    |

# Illustration for approximation: Borehole function

Approximation in hierarchical Tucker format with a linearly structured tree:

$$\mathcal{T} = \{\{1\}, \dots, \{d\}, \{1, 2\}, \dots, \{1, \dots, d-1\}\}$$



$$u^* = \sum_{i_1=1}^{r_1} \dots \sum_{i_d=1}^{r_d} \sum_{k_2=1}^{r_{1,2}} \dots \sum_{k_{d-1}=1}^{r_{1,\dots,d-1}} v_{i_1}^{(1)}(x_1) \dots v_{i_d}^{(d)}(x_d) C_{i_1, i_2, k_2}^{(1,2)} C_{k_2, i_3, k_3}^{(1,2,3)} \dots C_{k_{d-2}, i_{d-1}, k_{d-1}}^{(1,\dots,d-1)} C_{k_{d-1}, i_d}^{(1,\dots,d)}$$

with polynomial functions  $v_{i_\nu}^{(\nu)} \in V_\nu = \mathbb{P}_q$ .

## Illustration for approximation: Borehole function

Table: Approximation with **prescribed precision**  $\epsilon$ , **adaptive degree**  $p(\epsilon) = \log_{10}(\epsilon^{-1})$ , and  $N_\alpha = \dim(V_\alpha)$ . **Confidence intervals for relative error**  $\epsilon(u^*)$ , **storage complexity**  $S$  and **number of evaluations**  $M$  for different  $\epsilon$ , and **average ranks**.

| $\epsilon$ | $\epsilon(u^*)$             | $N$          | $S$        | $[r_{\{1\}}, \dots, r_{\{d\}}, r_{\{1,2\}}, \dots, r_{\{1,\dots,d-1\}}]$ |
|------------|-----------------------------|--------------|------------|--|
| $10^{-1}$  | $[1.8; 2.7] \times 10^{-1}$ | [39, 39]     | [23, 23]   | [1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1]                                  |
| $10^{-2}$  | $[0.3; 4.0] \times 10^{-2}$ | [88, 100]    | [41, 46]   | [1, 1, 1, 1, 1, 1, 1, 1, 1, 2, 1, 2, 1, 1]                               |
| $10^{-3}$  | $[0.8; 1.9] \times 10^{-3}$ | [159, 186]   | [61, 78]   | [2, 1, 1, 2, 2, 1, 1, 1, 1, 2, 2, 2, 1, 1]                               |
| $10^{-4}$  | $[2.5; 5.6] \times 10^{-5}$ | [328, 328]   | [141, 141] | [2, 2, 2, 3, 3, 2, 2, 2, 1, 2, 2, 2, 2, 2]                               |
| $10^{-5}$  | $[0.6; 1.6] \times 10^{-5}$ | [444, 472]   | [166, 178] | [2, 2, 2, 4, 4, 2, 2, 2, 1, 2, 2, 2, 2, 2]                               |
| $10^{-6}$  | $[3.1; 5.7] \times 10^{-6}$ | [596, 664]   | [204, 241] | [3, 2, 2, 4, 5, 3, 2, 2, 2, 2, 2, 2, 2, 2]                               |
| $10^{-7}$  | $[1.0; 6.3] \times 10^{-7}$ | [1042, 1267] | [374, 429] | [4, 3, 4, 6, 5, 3, 3, 3, 2, 2, 3, 2, 2, 2]                               |
| $10^{-8}$  | $[1.1; 7.1] \times 10^{-8}$ | [1567, 1567] | [512, 512] | [4, 3, 4, 7, 6, 3, 3, 3, 2, 2, 3, 2, 3, 3]                               |
| $10^{-9}$  | $[0.2; 4.9] \times 10^{-8}$ | [1719, 1854] | [534, 560] | [4, 4, 4, 8, 6, 3, 3, 3, 2, 2, 3, 2, 3, 3]                               |
| $10^{-10}$ | $[0.3; 1.9] \times 10^{-9}$ | [2482, 2828] | [774, 838] | [5, 4, 6, 10, 7, 4, 3, 3, 2, 2, 3, 2, 3, 3]                              |



The proposed algorithm

- provides an **approximation of a function in tree-based format using evaluations** of the function at a **structured and adapted set of points**,
- provides a **stable approximation with prescribed rank**, with a number of samples  $N$  equal to (or of the order of) the number of parameters,
- provides an approximation with **almost the desired precision**.

What should be done:

- Control norms of projections and statistical estimations of principal subspaces for **obtaining a certified approximation**.
- Provide **a priori estimations of the complexity** for certain classes of functions.



A. Nouy.

Higher-order principal component analysis for the approximation of tensors in tree-based low-rank formats.

*ArXiv e-prints*, 2017.



A. Nouy.

Low-rank methods for high-dimensional approximation and model order reduction.

In P. Benner, A. Cohen, M. Ohlberger, and K. Willcox (eds.), *Model Reduction and Approximation: Theory and Algorithms*. SIAM, Philadelphia, PA, 2016.

# Tree-based tensor formats I

The **minimal subspace**  $U_\alpha^{min}(u)$  of  $u$  is the smallest subspace such that

$$\mathcal{M}_\alpha(u) \in U_\alpha^{min}(u) \otimes \mathcal{H}_\alpha^c$$

and  $\text{rank}_\alpha(u) = \dim(U_\alpha^{min}(u))$ .

- Any tensor  $v$  is such that

$$v \in \bigotimes_{\alpha \in S(D)} U_\alpha^{min}(v)$$

with  $S(D)$  a partition of  $D$ , and

$$U_\alpha^{min}(v) \subset \bigotimes_{\beta \in S(\alpha)} U_\beta^{min}(v)$$

for any  $\alpha \subsetneq D$  with non trivial partition  $S(\alpha)$ .

## Tree-based tensor formats II

- For a tensor  $v \in \mathcal{T}_r^T$  with  $\text{rank}_T(v) = r$ , let  $\{\varphi_{k_\alpha}^{(\alpha)}\}_{k_\alpha=1}^{r_\alpha}$  be bases of the minimal subspace  $U_\alpha^{\min}(v)$ . The tensor  $v$  then admits a **hierarchical representation**

$$v = \sum_{\substack{1 \leq k_\alpha \leq r_\alpha \\ \alpha \in S(D)}} p_{(k_\alpha)_{\alpha \in S(D)}}^D \bigotimes_{\alpha \in S(D)} \varphi_{k_\alpha}^\alpha,$$

with

$$\varphi_{k_\alpha}^\alpha = \sum_{\substack{1 \leq k_\beta \leq r_\beta \\ \beta \in S(\alpha)}} p_{k_\alpha, (k_\beta)_{\beta \in S(\alpha)}}^\alpha \bigotimes_{\beta \in S(\alpha)} \varphi_{k_\beta}^\beta$$

## Partial interpolation of tensors

For a subspace  $W_\alpha$  in  $\mathcal{H}_\alpha$ , we define a **unisolvant set of points**  $\Gamma_{W_\alpha}$  in  $\mathcal{X}_\alpha$  (**magic points**) and the associated **interpolation operator**  $I_{W_\alpha}$  onto  $W_\alpha$  defined for  $v \in \mathbb{R}^{\mathcal{X}_\alpha}$  by

$$I_{W_\alpha} v(x_\alpha) = v(x_\alpha) \quad \forall x_\alpha \in \Gamma_{W_\alpha}.$$

We then define the corresponding partial interpolation operator  $\mathcal{I}_{W_\alpha} = I_{W_\alpha} \otimes id_{\alpha^c}$  defined for  $u \in \mathbb{R}^{\mathcal{X}}$  by

$$\mathcal{I}_{W_\alpha} u(x_\alpha, \cdot) = u(x_\alpha, \cdot) \quad \forall x_\alpha \in \Gamma_{W_\alpha}.$$

If the minimal subspace  $U_\alpha^{min}(u)$  is a RKHS, then  $I_{W_\alpha}$  is continuous from  $U_\alpha^{min}(u)$  to  $W_\alpha$  and  $\mathcal{I}_{W_\alpha}$  is continuous from  $U_\alpha^{min}(u) \otimes \mathcal{H}_{\alpha^c}$  to  $W_\alpha \otimes \mathcal{H}_{\alpha^c}$ , so that

$$\mathcal{I}_{W_\alpha} u \in W_\alpha \otimes \mathcal{H}_{\alpha^c}.$$

# Statistical estimation of principal components

For  $\alpha \in T$ , consider  $u_\alpha = \mathcal{I}_{V_\alpha} u$ .

For  $\|\cdot\|$  the  $L^2_\mu(\mathcal{X})$ -norm, the  $\alpha$ -principal subspace of  $u_\alpha$  is solution of

$$\min_{\dim(U_\alpha)=r_\alpha} \mathbb{E} \left( \|u_\alpha(\cdot, X_{\alpha^c}) - \mathcal{P}_{U_\alpha} u_\alpha(\cdot, X_{\alpha^c})\|_{L^2_\mu(\mathcal{X}_\alpha)}^2 \right),$$

where  $u_\alpha(\cdot, X_{\alpha^c})$  is interpreted as a  $V_\alpha$ -valued random variable.

It can be estimated by the solution of

$$\min_{\dim(U_\alpha)=r_\alpha} \frac{1}{N_\alpha} \sum_{k=1}^{N_\alpha} \|u_\alpha(\cdot, x_{\alpha^c}^k) - \mathcal{P}_{U_\alpha} u_\alpha(\cdot, x_{\alpha^c}^k)\|_{\mathcal{H}_\alpha}^2$$

where the  $x_{\alpha^c}^k$  are i.i.d. samples of  $X_{\alpha^c}$ .

- The **storage complexity** (number of parameters) of a tensor in  $\mathcal{T}_r^T \cap \mathcal{V}$  is

$$S = \sum_{\alpha \in (T \cup \{D\}) \setminus \mathcal{L}(T)} r_\alpha \prod_{\beta \in S(\alpha)} r_\beta + \sum_{\alpha \in \mathcal{L}(T)} r_\alpha \dim(V_\alpha).$$

- The **total number of evaluations** of the function required by the algorithm is

$$N = \sum_{\alpha \in \mathcal{L}(T)} N_\alpha \dim(V_\alpha) + \sum_{\alpha \in T \setminus \mathcal{L}(T)} N_\alpha \prod_{\beta \in S(\alpha)} r_\beta + \prod_{\beta \in S(D)} r_\beta,$$

where  $N_\alpha$  is the number of samples used for estimating the  $r_\alpha$   $\alpha$ -principal components of  $u_\alpha$ , taken such that

$$r_\alpha \leq N_\alpha$$

- If  $N_\alpha = r_\alpha$  for all  $\alpha$ , then

$$N = S.$$

## About the constants

If oblique projections  $I_{U_\alpha}$  and  $I_{V_\alpha}$  were orthogonal projections, the constants  $\Lambda$  and  $\tilde{\Lambda}$  would be equal to 1.

These constants  $\Lambda$  and  $\tilde{\Lambda}$  depend on

$$\|I_{V_\alpha}\|_{U_\alpha^{\min}(u) \rightarrow \mathcal{H}_\alpha} \quad \text{and} \quad \|I_{U_\alpha} - P_{U_\alpha}\|_{U_\alpha^{\min}(u) \rightarrow \mathcal{H}_\alpha}$$

that depend on the properties of **oblique projection operators restricted to minimal subspaces of  $u$** .

### Case of tensor recovery

Assume that  $U_\alpha^{\min}(u) \subset V_\alpha$  for all leaves  $\alpha$  (no discretization error).

If for all  $\alpha \in T$ , the set of  $N_\alpha$  samples  $u(\cdot, x_{\alpha^c}^k)$  contains  $\text{rank}_\alpha(u)$  linearly independent functions, then  $U_\alpha = U_\alpha^{\min}(u)$ .

The constants  $\Lambda = 1$ , and  $\tilde{\Lambda} = 1$  (i.e. same stability than the ideal algorithm).