## Reconstruction of ancestral gene orders

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Banff, February 17, 2017

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Introduction


## Mouse X-Chromosome



## Human X-Chromosome

(Pevzner and Tesler, 2003)







## Genome Rearrangement Problems

- Distance: Minimum \# of rearrangements from $A$ to $B$ ?


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- Scenario: Which rearrangements? (also called Sorting)
- Phylogeny: How did the genomes evolve?
- Ancestral Reconstruction: How do the ancestors look like?


Ancestral Reconstruction


Input: Tree and genomes $A, B, C, D$
Ouput: Ancestral genomes ( $M_{1}, M_{2}, M_{3}$ )

## Ancestral Reconstruction

- Distance-based Methods
- Homology-based Methods


$$
(-2+1+3+4) \quad(+2-1+3+4) \quad(+1+3-2+4) \quad(+1-3-2+4)
$$

## Distance-based Methods



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## Distance-based Methods



Find ancestral genomes that minimize events on the tree $\rightarrow$ Small Parsimony Problem


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## Ancestral Reconstruction methods

- Distance-based methods:
- Assume a rearrangement model
- Minimize branch lengths
- Homology-based methods:
- Find conserved structures
- Maximize some weight/probability function



Ancestral Reconstruction where internal nodes are

Intermediate Genomes of its children.

## Definitions

2
$2 \quad-3$
4

## Genome model



| 1 | 2 | -3 | 4 |
| :---: | :---: | :---: | :---: |
| $1^{t}$ | $1^{h} 2^{t}$ | $2^{h} 3^{h}$ | $3^{t} 4^{t}$ |

## The Double-Cut-and-Join (DCJ) operation

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$$
\begin{gathered}
A=\left\{\circ 1^{t}, 1^{h} 2^{t}, 2^{h} 3^{t}, 3^{h} 4^{t}, 4^{h} \circ, \circ 5^{t}, 5^{h} 6^{t}, 6^{h} 7^{t}, 7^{h} \circ\right\} \\
B=\left\{1^{h} 2^{h}, 2^{t} 3^{h}, 3^{t} 4^{t}, 4^{h} 1^{t}, \circ 6^{t}, 6^{h} 5^{t}, 5^{h} 7^{h}, 7^{t} \circ\right\}
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$$

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\end{gathered}
$$


$A$-edges are drawn in green, and $B$-edges in blue.

## $d_{\mathrm{DCJ}}(A, B)=N-C$

where $N$ is the number of genes and $C$ is the number of cycles in $B P(A, B)$.
(Bergeron et al, 2006)

Genomes are matchings in the BP graph:


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\Rightarrow
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Intermediate Genome of $A$ and $B$

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Intermediate Genome of $A$ and $B$


- Very easy to detect (linear time)
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- Reduces the search space. In the example:
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- $34,459,425$ possible genomes
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- Reduces the search space. In the example:
- $34,459,425$ possible genomes
- Only 40 intermediate genomes between A and B .


## Methods

## Ancestral Reconstruction with IG



- Small parsimony with the restriction that internal nodes are IG's of the children.


## Ancestral Reconstruction with IG



- Small parsimony with the restriction that internal nodes are IG's of the children.
- Still NP-hard
- Given adjacency weights, can we find an IG with maximum weight?
- Maximum Weight Independent Set: Polynomial Time
- DeClone (Chauve et al., 2015)
- New proposed algorithm based on InferCARs (Ma et al., 2006).












## Genomes with unique genes



## Genomes with unique genes



DCJ InDel Model (Braga et al., 2010; Compeau, 2012)



## Breakpoint graph with Unique genes



## Breakpoint graph with Unique genes



New components: AA-, BB-, AB- , A-, and B-paths.

## Breakpoint graph with Unique genes



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Find an optimal completion

## Breakpoint graph with Unique genes



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## Optimal completion

- $A A$ - and $B B$ - components are closed

- $A B$ - are paired

- $A$ - and $B$ - paths are paired, with opposing parity.
- Sometimes $A$-, $B$ - and $A B$ - paths are joined in a triplet.
- How to find a completion with maximum weight?
- Calculate all possible pairings and solve a Maximum Weight Matching

- Sometimes $A-, A B-, B$ - triplets are possible.
- Triple matching is usually NP-hard, but it is still open in this case.

Results

- RINGO - ancestral Reconstruction with INtermediate GenOmes (Feijao and Araujo, 2016)
- MGRA2 (Avdeyev et al., 2016)
- MLGO (Hu et al., 2014)



| Dataset | $I=1$, unitary indels |  |  |  |  |
| :--- | ---: | ---: | ---: | ---: | ---: |
| Diameter (D) | $0.5 n$ | $1 n$ | $1.5 n$ | $2 n$ | $2.5 n$ |
| RINGO | 3 s | 3 s | 5 s | 7 s | 7 s |
| MLGO | 1 m 6 s | 1 m 10 s | 1 m 7 s | 1 m 9 s | 1 m 16 s |
| MGRA | 7 s | 1 m 46 s | 12 m 12 s | 56 m 55 s | 2 h 2 m 41 s |

## Current Challenges

- Duplicated genes
- Statistical models


## Acknowledgments

- Elói Araújo (UFMS, Brazil)

- Jens Stoye (Bielefeld University, Germany)

|  | Technische Fakultät |
| :--- | :--- |
| Universität Bielefeld |  | | Genominformatik |
| :--- |

Thanks!

