# Reconstruction of ancestral gene orders

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Technische Fakultät Genominformatik

# Introduction







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Genomes



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#### Human X-Chromosome

(Pevzner and Tesler, 2003)

Genomes



Genomes









## **Rearrangement Scenario**



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## • **Distance**: Minimum # of rearrangements from A to B?

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- Scenario: Which rearrangements? (also called Sorting)
- Phylogeny: How did the genomes evolve?
- Ancestral Reconstruction: How do the ancestors look like?



# **Ancestral Reconstruction**

#### **Ancestral Reconstruction**

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**Input**: Tree and genomes A, B, C, D

**Ouput**: Ancestral genomes  $(M_1, M_2, M_3)$ 

• Distance-based Methods

• Homology-based Methods

#### **Distance-based Methods**



#### **Distance-based Methods**



#### **Distance-based Methods**

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Find ancestral genomes that minimize events on the tree

 $\rightarrow$  Small Parsimony Problem

#### Homology-based Methods





#### Homology-based Methods



#### Homology-based Methods



- Distance-based methods:
  - Assume a rearrangement model
  - Minimize branch lengths
- Homology-based methods:
  - Find conserved structures
  - Maximize some weight/probability function

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## **Intermediate Genomes**



## **Intermediate Genomes**



# Ancestral Reconstruction where internal nodes are Intermediate Genomes of its children.

# Definitions







$$A = \{\circ 1^t, 1^h 2^t, 2^h 3^h, 3^t 4^t, 4^h \circ\}$$

# The Double-Cut-and-Join (DCJ) operation

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# The Double-Cut-and-Join (DCJ) operation

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$$A = \{\circ 1^{t}, 1^{h} 2^{t}, 2^{h} 3^{t}, 3^{h} 4^{t}, 4^{h} \circ, \circ 5^{t}, 5^{h} 6^{t}, 6^{h} 7^{t}, 7^{h} \circ\}$$
$$B = \{1^{h} 2^{h}, 2^{t} 3^{h}, 3^{t} 4^{t}, 4^{h} 1^{t}, \circ 6^{t}, 6^{h} 5^{t}, 5^{h} 7^{h}, 7^{t} \circ\}$$


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A-edges are drawn in green, and B-edges in blue.

## $d_{\mathsf{DCJ}}(A,B) = N - C$

where N is the number of genes and C is the number of cycles in BP(A, B).

(Bergeron et al, 2006)

Genomes are **matchings** in the BP graph:



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Edges are non-crossing chords in the cycles of BP(A, B)

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 $\Rightarrow$ 

#### Intermediate Genome of A and B

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 $\Rightarrow$ 

Intermediate Genome of A and B



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  - 34,459,425 possible genomes

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- Reduces the search space. In the example:
  - 34,459,425 possible genomes
  - Only 40 intermediate genomes between A and B.

# Methods

#### Ancestral Reconstruction with IG



• Small parsimony with the restriction that internal nodes are IG's of the children.

#### Ancestral Reconstruction with IG



- Small parsimony with the restriction that internal nodes are IG's of the children.
- Still NP-hard

• Given adjacency weights, can we find an IG with maximum weight?

• Maximum Weight Independent Set: Polynomial Time

• DeClone (Chauve et al., 2015)

• New proposed algorithm based on InferCARs (Ma et al., 2006).

Example



Example



Example







+3,+4):0.9; (+1,+3):0.5; (-1,+3):0.4, ... Α В С D

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#### Genomes with unique genes

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DCJ InDel Model (Braga et al., 2010; Compeau, 2012)





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#### New components: AA-, BB-, AB- , A-, and B-paths.

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Find an optimal completion

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#### New components: AA-, BB-, AB- , A-, and B-paths.

Find an optimal completion

• AA- and BB- components are closed

• AB- are paired

• Sometimes A-, B- and AB- paths are joined in a triplet.

• A- and B- paths are paired, with opposing parity.





• How to find a completion with maximum weight?

• Calculate all possible pairings and solve a *Maximum Weight Matching* 

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## Maximum Weight Matching

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• Sometimes A-, AB-, B- triplets are possible.

Triple matching is usually NP-hard, but it is still open in this case.

## Results

• RINGO - ancestral Reconstruction with INtermediate GenOmes (Feijao and Araujo, 2016)

• MGRA2 (Avdeyev et al., 2016)

• MLGO (Hu et al., 2014)







Dataset	I = 1, unitary indels					
Diameter (D)	0.5n	1n	1.5n	2n	2.5n	
RINGO	3s	3s	5s	7s	7s	
MLGO	1 m 6 s	1 m 10 s	$1 \mathrm{m7s}$	1 m9 s	1 m 1 6 s	
MGRA	7s	$1 \mathrm{m} 46 \mathrm{s}$	$12\mathrm{m}12\mathrm{s}$	$56\mathrm{m}55\mathrm{s}$	2h2m41s	

• Duplicated genes

• Statistical models

• Elói Araújo (UFMS, Brazil)



• Jens Stoye (Bielefeld University, Germany)



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## Thanks!