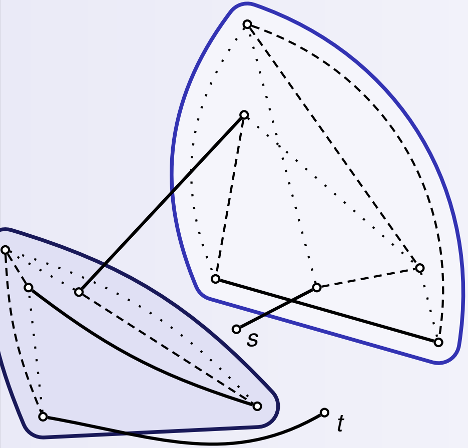
A 1.5-Approximation for Path TSP

Rico Zenklusen

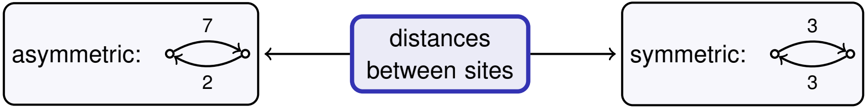
ETH Zurich

Presentation: Martin Nägele, ETH Zurich

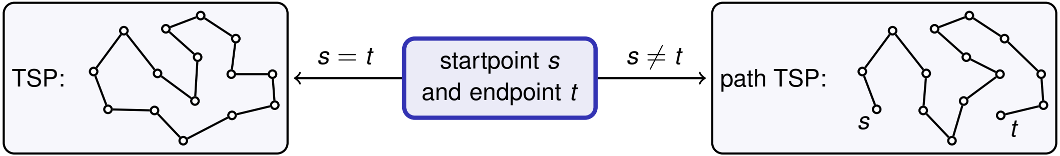


A brief intro to the Traveling Salesman Problem

Common Variations of TSP

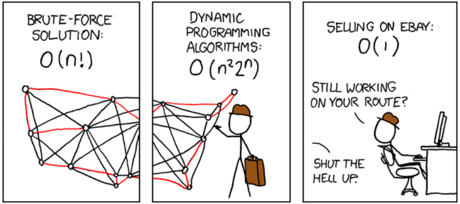


- Complete graph $G = (V, E)$.
- Metric length $\ell \colon E \to \mathbb{R}_{\geqslant 0}$.



What’s known?

- All variants are well-known to be APX-hard.



- Major open problem what efficient computation can achieve.

TSP	path TSP
	1.667 [Hoogeveen, 1991]
	1.618 [An, Kleinberg, Shmoys, 2012]
	1.6 [Sebő, 2013]
1.5 [Christofides, 1978]	1.599 [Vygen, 2016]
	1.566 [Gottschalk, Vygen, 2016]
	1.529 [Sebő, van Zuylen, 2016]
	$1.5 + \varepsilon$ [Traub, Vygen, 2018a]

Exciting progress for graph metrics:

[Oveis Gharan, Saberi, Singh, 2011]

[Mucha, 2014]

[Sebő, Vygen, 2014]

[Mömke, Svensson, 2016]

[Traub, Vygen, 2018b]

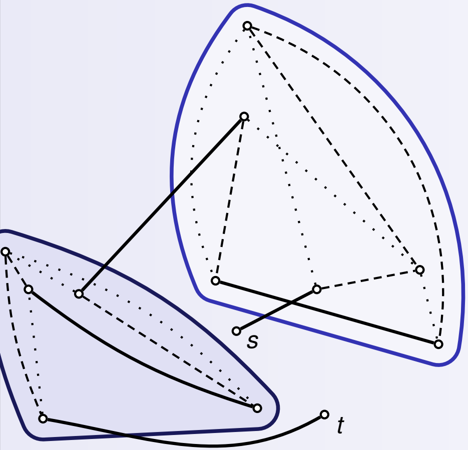
[. . .]

Our contribution

There is a 1.5-approximation for path TSP.

- We move away from prior approaches, which focussed on so-called *narrow cuts*.
- Technical ingredients: Obtain a strong Held-Karp solution z using
  - Karger's bound on the number of near-min cuts, and
  - Dynamic programming "à la Traub & Vygen".

  Run a Christofides-type algorithm with a spanning tree obtained from z.
- Analysis follows Wolsey's approach.
- Natural barrier 1.5: Any progress improves upon Christofides' 1.5-approximation for TSP.



Following in Christofides' footsteps

Why it works for TSP but fails for path TSP...

(Spoiler: ...and can be fixed.)

The general idea

- Find connected Eulerian graph with good total length, exploit metric lengths to shortcut.

The general idea

- Find connected Eulerian graph with good total length, exploit metric lengths to shortcut.

Start building a solution from a **spanning tree**

- Find connected Eulerian graph with good total length, exploit metric lengths to shortcut.

Start building a solution from a **spanning tree**

Add **edges** to correct degree parities

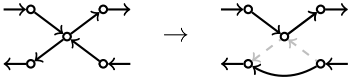
The general idea

- Find connected Eulerian graph with good total length, exploit metric lengths to shortcut.

Start building a solution from a **spanning tree**

Add **edges** to correct degree parities

obtain **Hamil-tonian** graph



The general idea

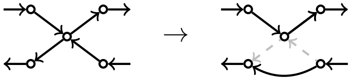
- Find connected Eulerian graph with good total length, exploit metric lengths to shortcut.

Start building a solution from a **spanning tree**

Add **edges** to correct degree parities

?

obtain **Hamiltonian** graph



Christofides' 1.5-approximation for TSP

1. Find a shortest spanning tree T.

$$\implies \ell(T) \leqslant \ell(\mathrm{OPT}) \ .$$

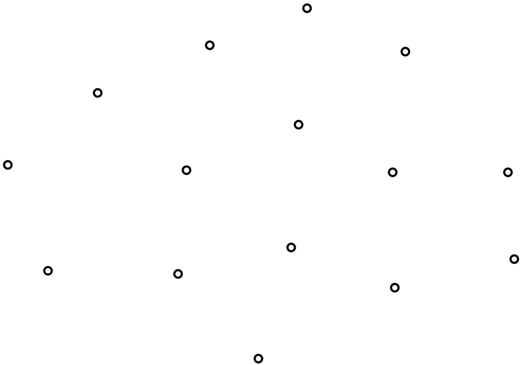
2. Find a shortest odd(T)-join J.

$$\implies \ell(J) \leqslant \frac{1}{2} \cdot \ell(\mathrm{OPT}) \ .$$

3. Find Eulerian tour in multiunion of T and J.

4. Return shortcutted Hamiltonian tour H.

$$\implies \ell(H) \leqslant \ell(T) + \ell(J) \leqslant \frac{3}{2} \cdot \ell(\mathrm{OPT}) \ .$$



Christofides' 1.5-approximation for TSP

1. Find a shortest spanning tree T.

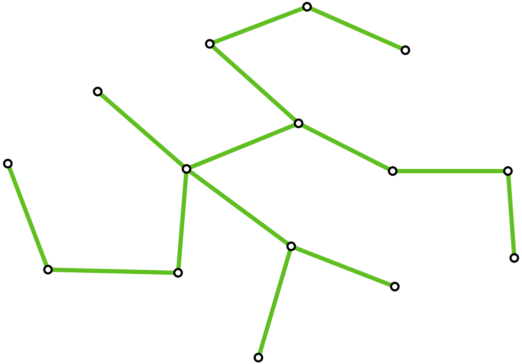
$$\implies \ell(T) \leqslant \ell(\mathrm{OPT}) \ .$$

2. Find a shortest odd(T)-join J.

$$\implies \ell(J) \leqslant \frac{1}{2} \cdot \ell(\mathrm{OPT}) \ .$$

3. Find Eulerian tour in multiunion of T and J.
4. Return shortcutted Hamiltonian tour H.

$$\implies \ell(H) \leqslant \ell(T) + \ell(J) \leqslant \frac{3}{2} \cdot \ell(\mathrm{OPT}) \ .$$



1. Find a shortest spanning tree T.

$$\implies \ell(T) \leqslant \ell(\mathrm{OPT}) \ .$$

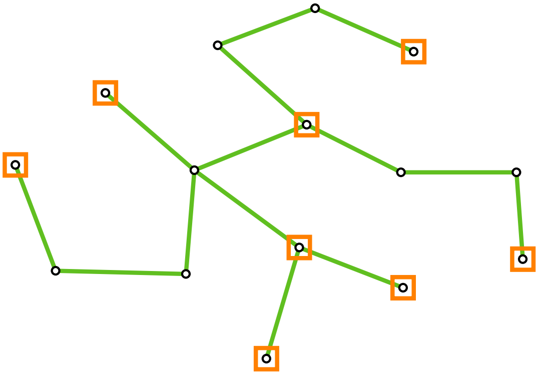
2. Find a shortest odd(T)-join J.

$$\implies \ell(J) \leqslant \frac{1}{2} \cdot \ell(\mathrm{OPT}) \ .$$

3. Find Eulerian tour in multiunion of T and J.

4. Return shortcutted Hamiltonian tour H.

$$\implies \ell(H) \leqslant \ell(T) + \ell(J) \leqslant \frac{3}{2} \cdot \ell(\mathrm{OPT}) \ .$$



□ $\mathrm{odd}(T) := \{v \in V \mid \deg(v) \text{ odd}\}$

Christofides' 1.5-approximation for TSP

1. Find a shortest spanning tree T.

$$\implies \ell(T) \leqslant \ell(\mathrm{OPT}) \ .$$

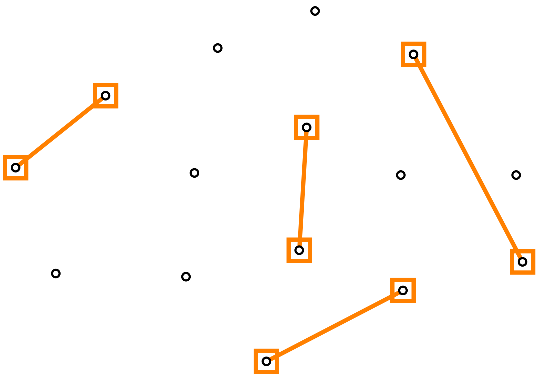
2. Find a shortest odd(T)-join J.

$$\implies \ell(J) \leqslant \frac{1}{2} \cdot \ell(\mathrm{OPT}) \ .$$

3. Find Eulerian tour in multiunion of T and J.

4. Return shortcutted Hamiltonian tour H.

$$\implies \ell(H) \leqslant \ell(T) + \ell(J) \leqslant \frac{3}{2} \cdot \ell(\mathrm{OPT}) \ .$$



$\square$ $\mathrm{odd}(T) := \{v \in V \mid \deg(v) \text{ odd}\}$

Christofides' 1.5-approximation for TSP

1. Find a shortest spanning tree T.

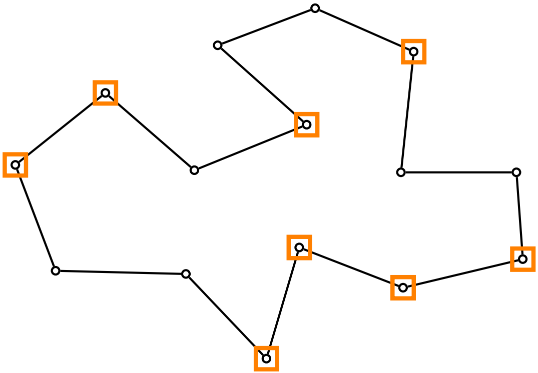
$$\implies \ell(T) \leqslant \ell(\mathrm{OPT}) \ .$$

2. Find a shortest odd(T)-join J.

$$\implies \ell(J) \leqslant \frac{1}{2} \cdot \ell(\mathrm{OPT}) \ .$$

3. Find Eulerian tour in multiunion of T and J.
4. Return shortcutted Hamiltonian tour H.

$$\implies \ell(H) \leqslant \ell(T) + \ell(J) \leqslant \frac{3}{2} \cdot \ell(\mathrm{OPT}) \ .$$



□ $\mathrm{odd}(T) := \{v \in V \mid \deg(v) \text{ odd}\}$

Christofides' 1.5-approximation for TSP

1. Find a shortest spanning tree T.

$$\implies \ell(T) \leqslant \ell(\mathrm{OPT}) \ .$$

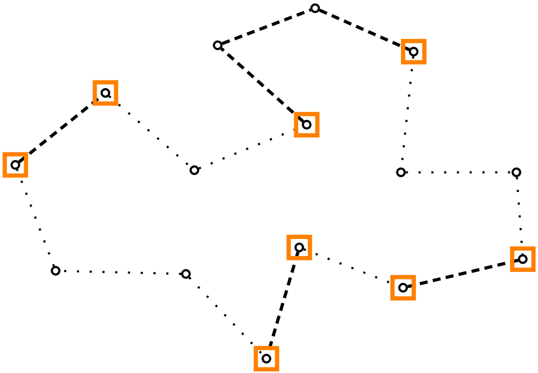
2. Find a shortest odd(T)-join J.

$$\implies \ell(J) \leqslant \frac{1}{2} \cdot \ell(\mathrm{OPT}) \ .$$

3. Find Eulerian tour in multiunion of T and J.

4. Return shortcutted Hamiltonian tour H.

$$\implies \ell(H) \leqslant \ell(T) + \ell(J) \leqslant \frac{3}{2} \cdot \ell(\mathrm{OPT}) \ .$$



$\square$ $\mathrm{odd}(T) := \{v \in V \mid \deg(v) \text{ odd}\}$

Christofides' 1.5-approximation for TSP

1. Find a shortest spanning tree T.

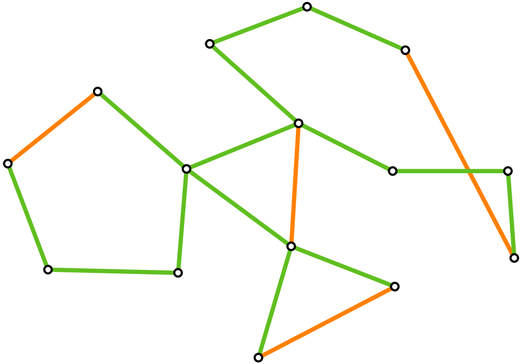
$$\Longrightarrow \ell(T) \leqslant \ell(\mathrm{OPT}) \ .$$

2. Find a shortest odd(T)-join J.

$$\Longrightarrow \ell(J) \leqslant \frac{1}{2} \cdot \ell(\mathrm{OPT}) \ .$$

3. Find Eulerian tour in multiunion of T and J.
4. Return shortcutted Hamiltonian tour H.

$$\Longrightarrow \ell(H) \leqslant \ell(T) + \ell(J) \leqslant \frac{3}{2} \cdot \ell(\mathrm{OPT}) \ .$$



Christofides' 1.5-approximation for TSP

1. Find a shortest spanning tree T.

$$\implies \ell(T) \leqslant \ell(\mathrm{OPT}) \; .$$

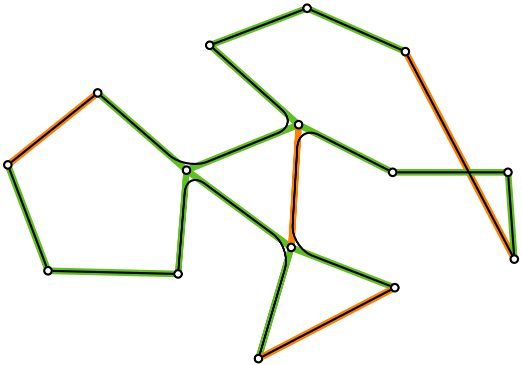
2. Find a shortest odd(T)-join J.

$$\implies \ell(J) \leqslant \frac{1}{2} \cdot \ell(\mathrm{OPT}) \; .$$

3. Find Eulerian tour in multiunion of T and J.

4. Return shortcutted Hamiltonian tour H.

$$\implies \ell(H) \leqslant \ell(T) + \ell(J) \leqslant \frac{3}{2} \cdot \ell(\mathrm{OPT}) \; .$$



Christofides' 1.5-approximation for TSP

1. Find a shortest spanning tree T.

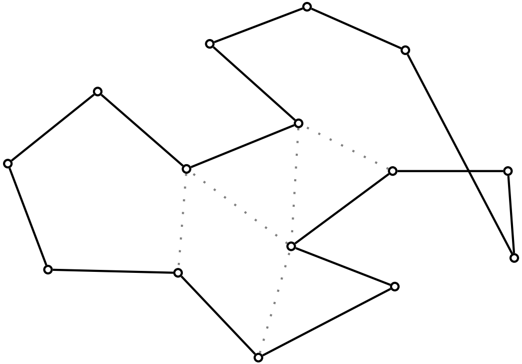
$$\implies \ell(T) \leqslant \ell(\mathrm{OPT}) \ .$$

2. Find a shortest odd(T)-join J.

$$\implies \ell(J) \leqslant \frac{1}{2} \cdot \ell(\mathrm{OPT}) \ .$$

3. Find Eulerian tour in multiunion of T and J.
4. Return shortcutted Hamiltonian tour H.

$$\implies \ell(H) \leqslant \ell(T) + \ell(J) \leqslant \frac{3}{2} \cdot \ell(\mathrm{OPT}) \ .$$

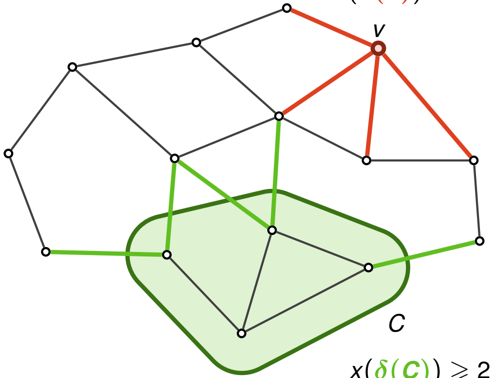


- Held-Karp polytope

$$P_{\mathrm{HK}} := \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(\delta(v)) = 2 & \forall v \in V \\ x(\delta(C)) \geqslant 2 & \forall C \subsetneq V,\ C \neq \emptyset \end{array} \right\} .$$

- Held-Karp relaxation

$$\min\{\ell^\top x \mid x \in P_{\mathrm{HK}}\} .$$



- Let $x^* \in \operatorname{argmin}\{\ell^\top x \mid x \in P_{\mathrm{HK}}\}$.

$$P_{\mathrm{HK}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(\delta(v)) = 2 & \forall v \in V \\ x(\delta(C)) \geqslant 2 & \forall C \subsetneq V,\ C \neq \emptyset \end{array} \right\}$$

Claim

If T is a shortest spanning tree, and J is a shortest odd(T)-join, then

(a) $\ell(T) \leqslant \ell^\top x^*$, and ***(b)*** $\ell(J) \leqslant \frac{1}{2} \cdot \ell^\top x^*$.

Wolsey's analysis

- Let $x^* \in \operatorname{argmin}\{\ell^\top x \mid x \in P_{\mathrm{HK}}\}$.

$$P_{\mathrm{HK}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(\delta(v)) = 2 & \forall v \in V \\ x(\delta(C)) \geqslant 2 & \forall C \subsetneq V,\ C \neq \emptyset \end{array} \right\}$$

Claim

If T is a shortest spanning tree, and J is a shortest odd(T)-join, then

(a) $\ell(T) \leqslant \ell^\top x^*$, and ***(b)*** $\ell(J) \leqslant \frac{1}{2} \cdot \ell^\top x^*$.

(a) $\frac{n-1}{n} \cdot x^* \in P_{\mathrm{ST}}$.

$$P_{\mathrm{ST}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(E) = |V| - 1 & \\ x(E[S]) \leqslant |S| - 1 & \forall S \subsetneq V,\ S \neq \emptyset \end{array} \right\}$$

- Let $x^* \in \operatorname{argmin}\{\ell^\top x \mid x \in P_{\text{HK}}\}$.

$$P_{\text{HK}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(\delta(v)) = 2 & \forall v \in V \\ x(\delta(C)) \geqslant 2 & \forall C \subsetneq V,\ C \neq \emptyset \end{array} \right\}$$

Claim

If T is a shortest spanning tree, and J is a shortest odd(T)-join, then

(a) $\ell(T) \leqslant \ell^\top x^*$, and ***(b)*** $\ell(J) \leqslant \frac{1}{2} \cdot \ell^\top x^*$.

(a) $\frac{n-1}{n} \cdot x^* \in P_{\text{ST}}$.

$$P_{\text{ST}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(E) = |V| - 1 & \\ x(E[S]) \leqslant |S| - 1 & \forall S \subsetneq V,\ S \neq \emptyset \end{array} \right\}$$

(b) $\frac{1}{2} \cdot x^* \in P^{\uparrow}_{Q\text{-join}}$ for any $Q \subseteq V$, $|Q|$ even.

$$P^{\uparrow}_{Q\text{-join}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; x(\delta(C)) \geqslant 1 \quad \forall C \subseteq V,\ |C \cap Q| \text{ odd} \right\}$$

- Let $x^* \in \operatorname{argmin}\{\ell^\top x \mid x \in P_{\mathrm{HK}}\}$.

$$P_{\mathrm{HK}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(\delta(v)) = 2 & \forall v \in V \\ x(\delta(C)) \geqslant 2 & \forall C \subsetneq V,\ C \neq \emptyset \end{array} \right\}$$

Claim

If T is a shortest spanning tree, and J is a shortest odd(T)-join, then

(a) $\ell(T) \leqslant \ell^\top x^*$, and ***(b)*** $\ell(J) \leqslant \frac{1}{2} \cdot \ell^\top x^*$.

(a) $\frac{n-1}{n} \cdot x^* \in P_{\mathrm{ST}}$.

$$P_{\mathrm{ST}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(E) = |V| - 1 & \\ x(E[S]) \leqslant |S| - 1 & \forall S \subsetneq V,\ S \neq \emptyset \end{array} \right\}$$

(b) $\frac{1}{2} \cdot x^* \in P^{\uparrow}_{Q\text{-join}}$ for any $Q \subseteq V$, $|Q|$ even.

$$P^{\uparrow}_{Q\text{-join}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; x(\delta(C)) \geqslant 1 \quad \forall C \subseteq V,\ |C \cap Q| \text{ odd} \right\}$$

- Shows 1.5-approximation and upper bound on integrality gap.

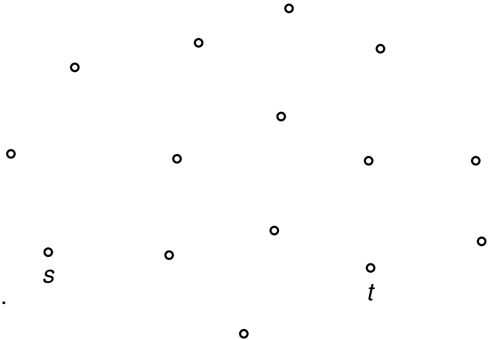
Christofides' approach for path TSP [Hoogeveen, 1991]

- Shortest spanning tree ***T***: $\ell(T) \leqslant \ell(\mathrm{OPT})$.
- But: OPT does not contain two disjoint Q_T-joins.
- Still, shortest Q_T-join ***J*** satisfies

$$\ell(J) \leqslant \tfrac{2}{3} \cdot \ell(\mathrm{OPT}).$$ [Hoogeveen, 1991]

Proof: Together, OPT and T contain three Q_T-joins.

- This algorithm is only $\frac{5}{3}$-approximate on some instances.

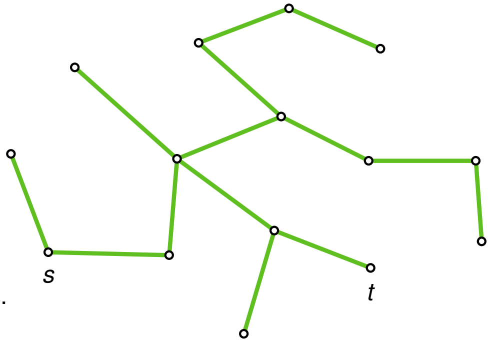


Christofides' approach for path TSP [Hoogeveen, 1991]

- Shortest spanning tree T: $\ell(T) \leqslant \ell(\mathrm{OPT})$.
- But: OPT does not contain two disjoint Q_T-joins.
- Still, shortest Q_T-join J satisfies

  $$\ell(J) \leqslant \tfrac{2}{3} \cdot \ell(\mathrm{OPT}). \qquad \text{[Hoogeveen, 1991]}$$

  Proof: Together, OPT and T contain three Q_T-joins.
- This algorithm is only $\frac{5}{3}$-approximate on some instances.

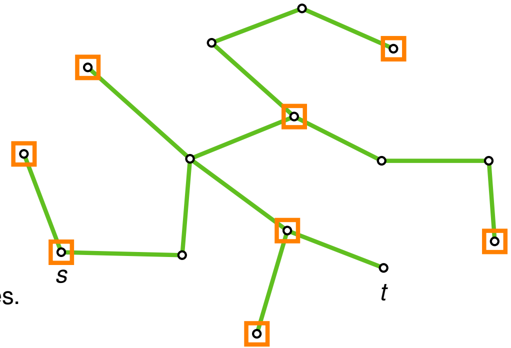


Christofides' approach for path TSP [Hoogeveen, 1991]

- Shortest spanning tree T: $\ell(T) \leqslant \ell(\mathrm{OPT})$.
- But: OPT does not contain two disjoint Q_T-joins.
- Still, shortest Q_T-join J satisfies

$$\ell(J) \leqslant \tfrac{2}{3} \cdot \ell(\mathrm{OPT}). \qquad \text{[Hoogeveen, 1991]}$$

  Proof: Together, OPT and T contain three Q_T-joins.
- This algorithm is only $\frac{5}{3}$-approximate on some instances.



$Q_T := \mathrm{odd}(T) \triangle \{s, t\}$

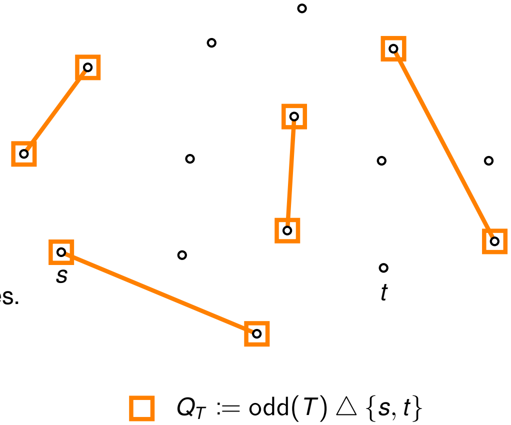
Christofides' approach for path TSP [Hoogeveen, 1991]

- Shortest spanning tree T: $\ell(T) \leqslant \ell(\mathrm{OPT})$.
- But: OPT does not contain two disjoint Q_T-joins.
- Still, shortest Q_T-join J satisfies

$$\ell(J) \leqslant \tfrac{2}{3} \cdot \ell(\mathrm{OPT}).$$ [Hoogeveen, 1991]

Proof: Together, OPT and T contain three Q_T-joins.

- This algorithm is only $\frac{5}{3}$-approximate on some instances.

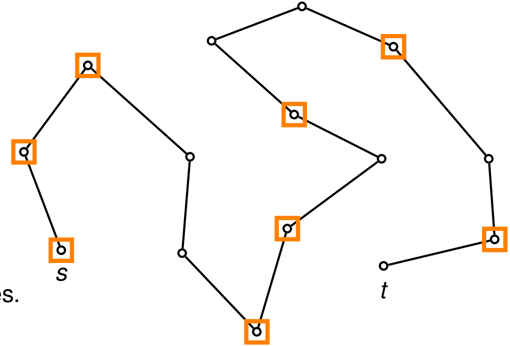


Christofides' approach for path TSP [Hoogeveen, 1991]

- Shortest spanning tree T: $\ell(T) \leqslant \ell(\mathrm{OPT})$.
- But: OPT does not contain two disjoint Q_T-joins.
- Still, shortest Q_T-join J satisfies

  $$\ell(J) \leqslant \tfrac{2}{3} \cdot \ell(\mathrm{OPT}).$$ [Hoogeveen, 1991]

  Proof: Together, OPT and T contain three Q_T-joins.
- This algorithm is only $\frac{5}{3}$-approximate on some instances.



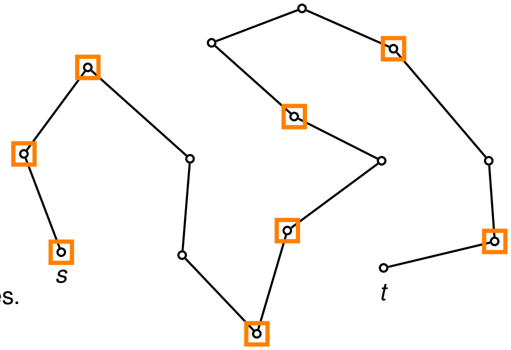
 $Q_T := \mathrm{odd}(T) \triangle \{s, t\}$

Christofides' approach for path TSP [Hoogeveen, 1991]

- Shortest spanning tree T: $\ell(T) \leqslant \ell(\mathrm{OPT})$.
- But: OPT does not contain two disjoint Q_T-joins.
- Still, shortest Q_T-join J satisfies

  $$\ell(J) \leqslant \frac{2}{3} \cdot \ell(\mathrm{OPT}). \qquad \text{[Hoogeveen, 1991]}$$

  Proof: Together, OPT and T contain three Q_T-joins.
- This algorithm is only $\frac{5}{3}$-approximate on some instances.



Goal: Find tree T with $\ell(T) \leqslant \ell(\mathrm{OPT})$ and s.t. shortest Q_T-join J satisfies $\ell(J) \leqslant \frac{1}{2} \cdot \ell(OPT)$.

- Held-Karp polytope for path TSP:

$$P_{\mathrm{HK}} := \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(\delta(v)) = 1 & v \in \{s,t\} \\ x(\delta(v)) = 2 & v \in V \setminus \{s,t\} \\ x(\delta(C)) \geqslant 1 & \forall C \subseteq V,\ |C \cap \{s,t\}| = 1 \\ x(\delta(C)) \geqslant 2 & \forall C \subsetneq V,\ C \neq \emptyset,\ |C \cap \{s,t\}| = 0 \end{array} \right\}$$

Where Wolsey's analysis fails

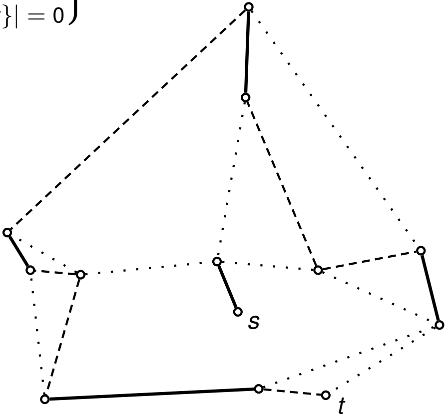
- Held-Karp polytope for path TSP:

$$P_{\mathrm{HK}} := \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(\delta(v)) = 1 & v \in \{s,t\} \\ x(\delta(v)) = 2 & v \in V \setminus \{s,t\} \\ x(\delta(C)) \geqslant 1 & \forall C \subseteq V,\ |C \cap \{s,t\}| = 1 \\ x(\delta(C)) \geqslant 2 & \forall C \subsetneq V,\ C \neq \emptyset,\ |C \cap \{s,t\}| = 0 \end{array} \right\}$$

$\cdots\cdots$ $x^*(e) = 1/3$

$----$ $x^*(e) = 2/3$

——— $x^*(e) = 1$



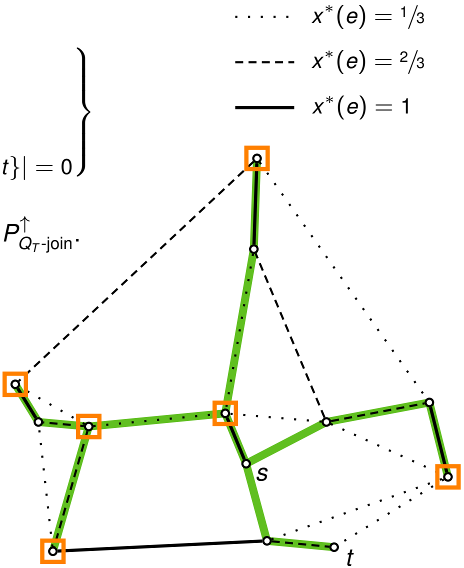
Where Wolsey's analysis fails

- Held-Karp polytope for path TSP:

$$P_{\text{HK}} := \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(\delta(v)) = 1 & v \in \{s,t\} \\ x(\delta(v)) = 2 & v \in V \setminus \{s,t\} \\ x(\delta(C)) \geqslant 1 & \forall C \subseteq V,\ |C \cap \{s,t\}| = 1 \\ x(\delta(C)) \geqslant 2 & \forall C \subsetneq V,\ C \neq \emptyset,\ |C \cap \{s,t\}| = 0 \end{array} \right\}$$

- Problem: $\frac{x^*}{2}$ for $x^* \in \operatorname{argmin}\{\ell^\top x \mid x \in P_{\text{HK}}\}$ infeasible for $P^{\uparrow}_{Q_T\text{-join}}$.

$$P^{\uparrow}_{Q_T\text{-join}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; x(\delta(C)) \geqslant 1 \quad \begin{array}{l} \forall C \subseteq V, \\ |C \cap Q_T| \text{ odd} \end{array} \right\}$$



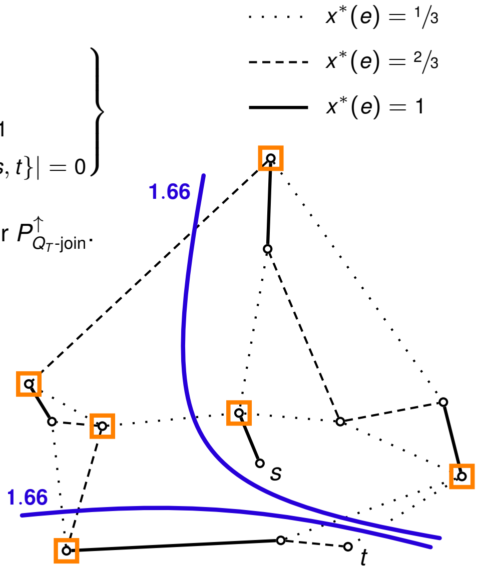
- Held-Karp polytope for path TSP:

$$P_{\text{HK}} := \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(\delta(v)) = 1 & v \in \{s,t\} \\ x(\delta(v)) = 2 & v \in V \setminus \{s,t\} \\ x(\delta(C)) \geqslant 1 & \forall C \subseteq V,\ |C \cap \{s,t\}| = 1 \\ x(\delta(C)) \geqslant 2 & \forall C \subsetneq V,\ C \neq \emptyset,\ |C \cap \{s,t\}| = 0 \end{array} \right\}$$

- Problem: $\frac{x^*}{2}$ for $x^* \in \operatorname{argmin}\{\ell^\top x \mid x \in P_{\text{HK}}\}$ infeasible for $P^\uparrow_{Q_T\text{-join}}$.

$$P^\uparrow_{Q_T\text{-join}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; x(\delta(C)) \geqslant 1 \quad \begin{array}{l} \forall C \subseteq V, \\ |C \cap Q_T| \text{ odd} \end{array} \right\}$$

- Infeasibility caused by *narrow* cuts:
  - $\rightarrow$ cuts C with $x^*(\delta(C)) < 2$.
  - $\rightarrow$ s-t-cuts, form a chain.
  - $\rightarrow$ appear in $P^\uparrow_{Q_T\text{-join}}$ only if $|T \cap \delta(C)|$ even.



Where Wolsey's analysis fails

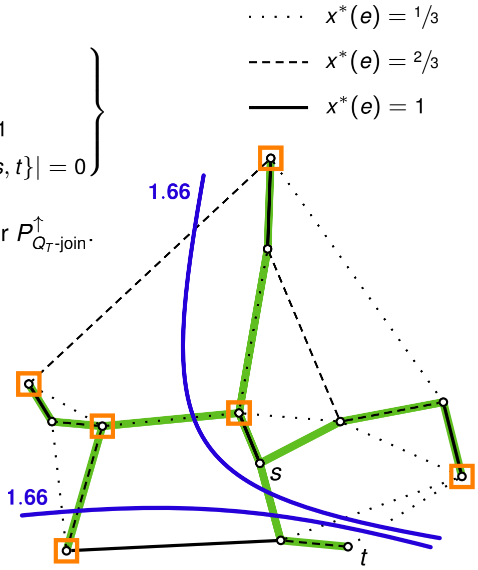
- Held-Karp polytope for path TSP:

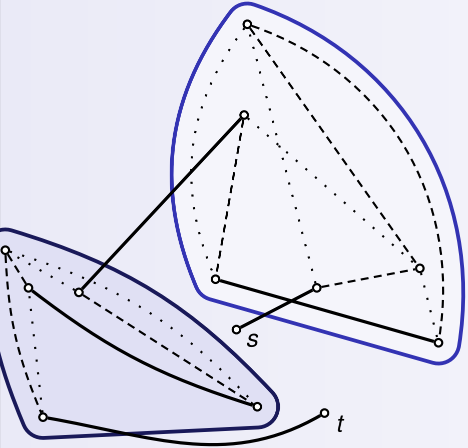
$$P_{\text{HK}} := \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; \begin{array}{ll} x(\delta(v)) = 1 & v \in \{s,t\} \\ x(\delta(v)) = 2 & v \in V \setminus \{s,t\} \\ x(\delta(C)) \geqslant 1 & \forall C \subseteq V,\ |C \cap \{s,t\}| = 1 \\ x(\delta(C)) \geqslant 2 & \forall C \subsetneq V,\ C \neq \emptyset,\ |C \cap \{s,t\}| = 0 \end{array} \right\}$$

- Problem: $\frac{x^*}{2}$ for $x^* \in \text{argmin}\{\ell^\top x \mid x \in P_{\text{HK}}\}$ infeasible for $P^{\uparrow}_{Q_T\text{-join}}$.

$$P^{\uparrow}_{Q_T\text{-join}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; x(\delta(C)) \geqslant 1 \quad \begin{array}{l} \forall C \subseteq V, \\ |C \cap Q_T| \text{ odd} \end{array} \right\}$$

- Infeasibility caused by *narrow* cuts:
  - $\rightarrow$ cuts C with $x^*(\delta(C)) < 2$.
  - $\rightarrow$ s-t-cuts, form a chain.
  - $\rightarrow$ appear in $P^{\uparrow}_{Q_T\text{-join}}$ only if $|T \cap \delta(C)|$ even.





1.5-approximation: The high-level plan

A new ingredient: Finding $\mathcal{B}(x^*)$-good points

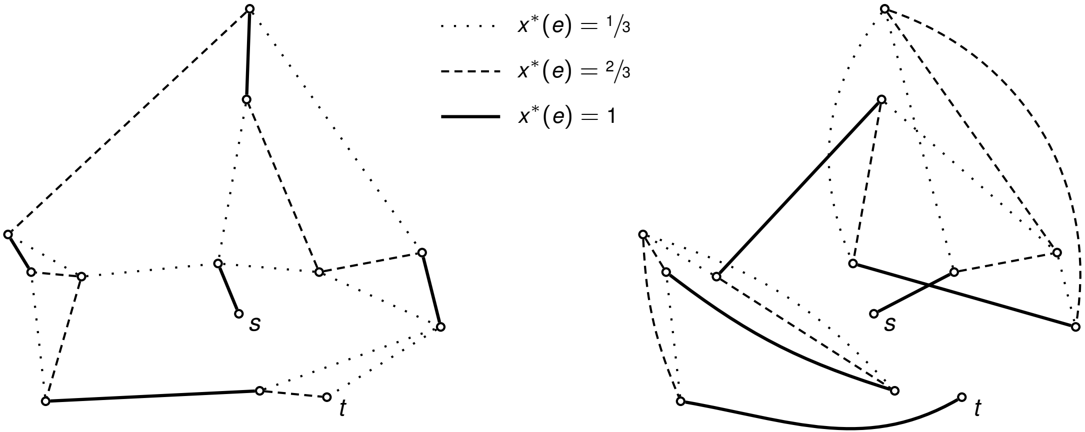
- Let $x^* \in \operatorname{argmin}\{\ell(x) \mid x \in P_{\text{HK}}\}$.
- Let
$$\mathcal{B}(x^*) := \{C \subseteq V \mid s \in C,\ t \notin C,\ x^*(\delta(C)) < 3\} \quad .$$
By Karger's result, $|\mathcal{B}(x^*)|$ is polynomially bounded. [Karger 1993]
- We will find a shortest point $y \in P_{\text{HK}}$ that is $\mathcal{B}(x^*)$-good:

> For each $B \in \mathcal{B}(x^*)$, either
> - $y(\delta(B)) \geqslant 3$, or
> - $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.

A new ingredient: Finding $\mathcal{B}(x^*)$-good points

$\mathcal{B}(x^*)$-good

$y \in P_{HK}$ is $\mathcal{B}(x^*)$-good: For all $B \in \mathcal{B}(x^*)$, ► $y(\delta(B)) \geqslant 3$, or ► $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.



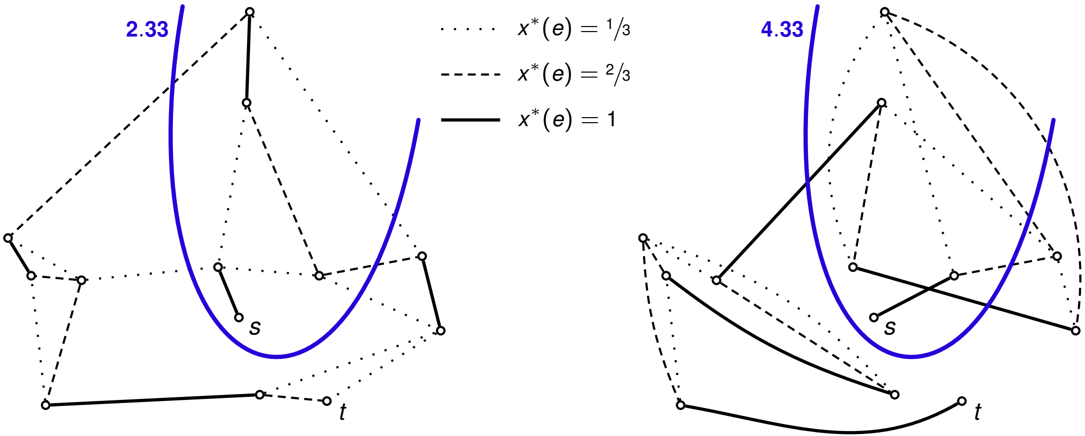
$x^* \in P_{HK}$

$\mathcal{B}(x^*)$-good point y

A new ingredient: Finding $\mathcal{B}(x^*)$-good points

$\mathcal{B}(x^*)$-good

$y \in P_{HK}$ is $\mathcal{B}(x^*)$-good: For all $B \in \mathcal{B}(x^*)$, ► $y(\delta(B)) \geqslant 3$, or ► $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.

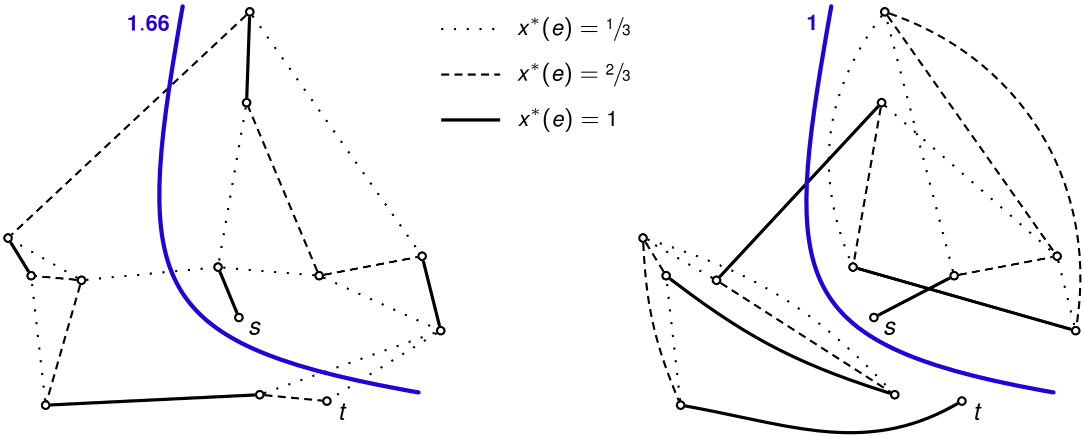


$x^* \in P_{HK}$ $\mathcal{B}(x^*)$-good point y

A new ingredient: Finding $\mathcal{B}(x^*)$-good points

$\mathcal{B}(x^*)$-good

$y \in P_{HK}$ is $\mathcal{B}(x^*)$-good: For all $B \in \mathcal{B}(x^*)$,
- $y(\delta(B)) \geqslant 3$, or
- $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.



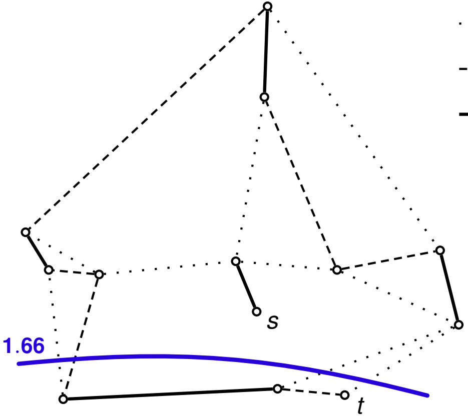
$x^* \in P_{HK}$

$\mathcal{B}(x^*)$-good point y

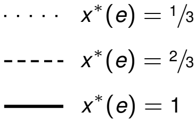
A new ingredient: Finding $\mathcal{B}(x^*)$-good points

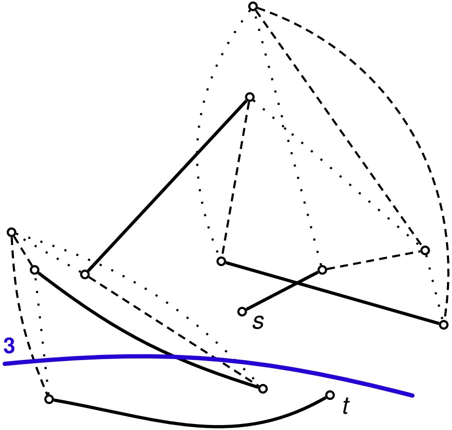
$\mathcal{B}(x^*)$-good

$y \in P_{HK}$ is $\mathcal{B}(x^*)$-good: For all $B \in \mathcal{B}(x^*)$, ▶ $y(\delta(B)) \geqslant 3$, or ▶ $y(\delta(B)) = 1$ and y is 0/1 on $\delta(B)$.



$x^* \in P_{HK}$



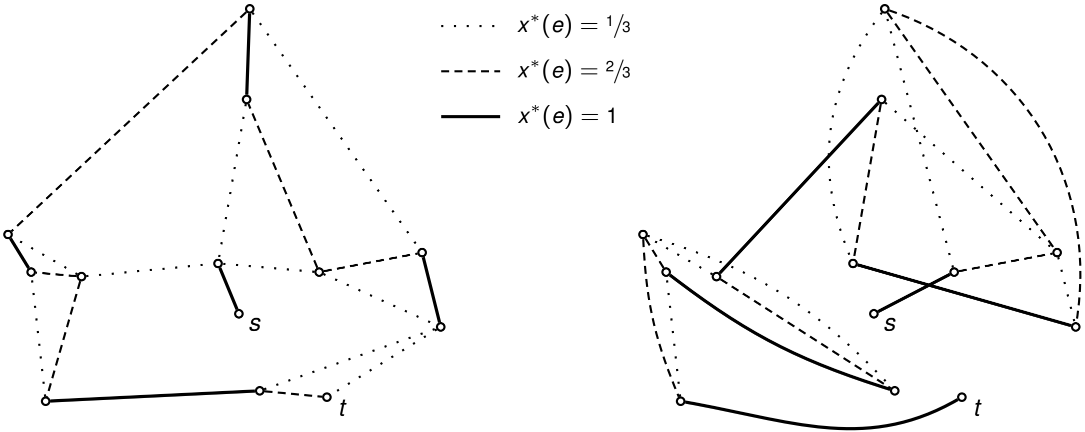


$\mathcal{B}(x^*)$-good point y

A new ingredient: Finding $\mathcal{B}(x^*)$-good points

$\mathcal{B}(x^*)$-good

$y \in P_{HK}$ is $\mathcal{B}(x^*)$-good: For all $B \in \mathcal{B}(x^*)$, ▶ $y(\delta(B)) \geqslant 3$, or ▶ $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.



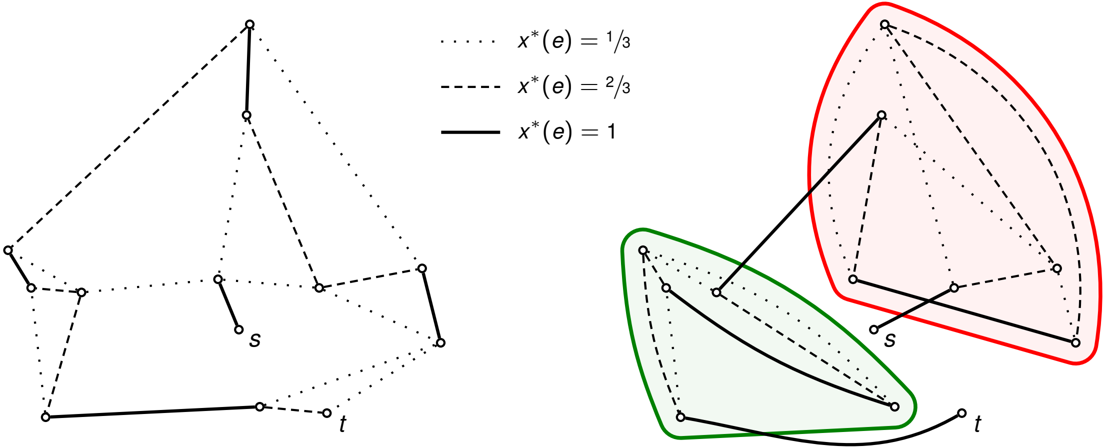
$x^* \in P_{HK}$

$\mathcal{B}(x^*)$-good point y

A new ingredient: Finding $\mathcal{B}(x^*)$-good points

$\mathcal{B}(x^*)$-good

$y \in P_{HK}$ is $\mathcal{B}(x^*)$-good: For all $B \in \mathcal{B}(x^*)$, ► $y(\delta(B)) \geqslant 3$, or ► $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.



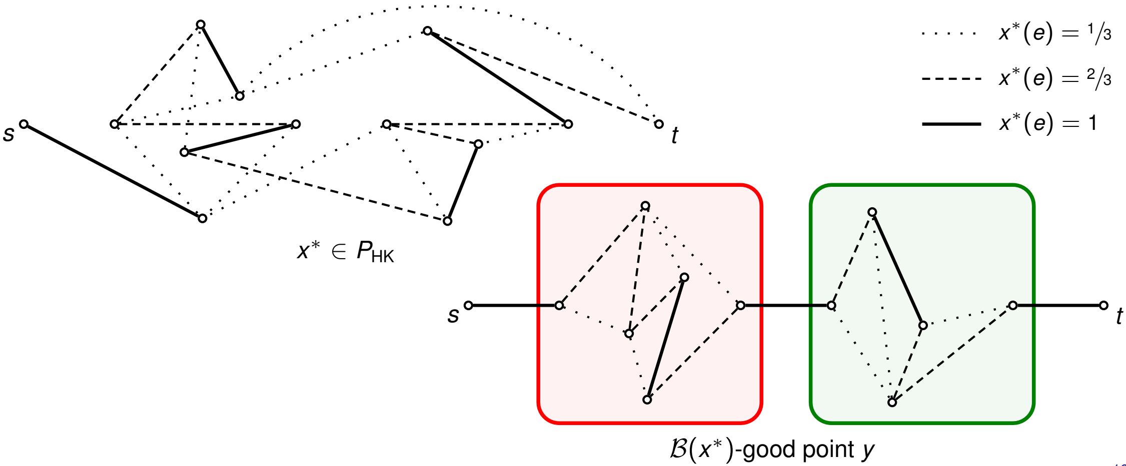
$x^* \in P_{HK}$

$\mathcal{B}(x^*)$-good point y

A new ingredient: Finding $\mathcal{B}(x^*)$-good points

$\mathcal{B}(x^*)$-good

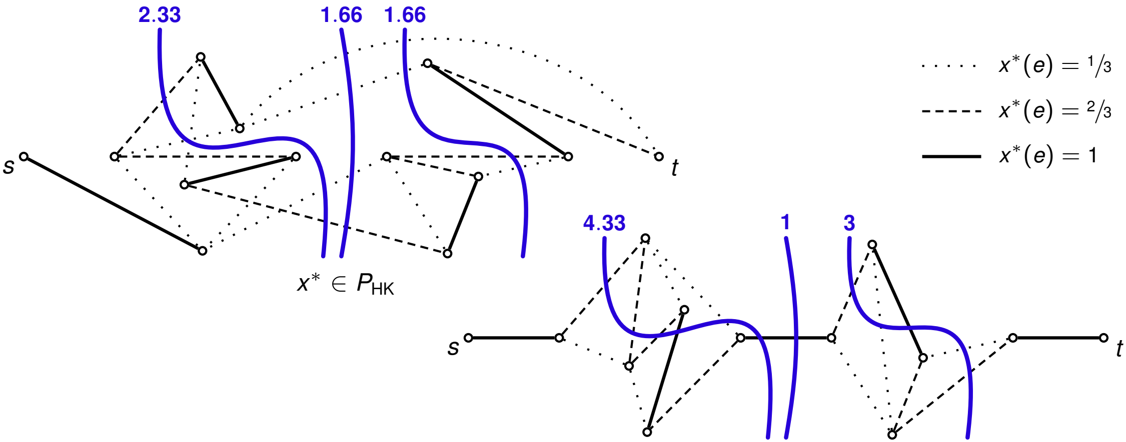
$y \in P_{\text{HK}}$ is $\mathcal{B}(x^*)$-good: For all $B \in \mathcal{B}(x^*)$, ▶ $y(\delta(B)) \geqslant 3$, or ▶ $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.



A new ingredient: Finding $\mathcal{B}(x^*)$-good points

$\mathcal{B}(x^*)$-good

$y \in P_{\text{HK}}$ is $\mathcal{B}(x^*)$-good: For all $B \in \mathcal{B}(x^*)$, ▶ $y(\delta(B)) \geqslant 3$, or ▶ $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.



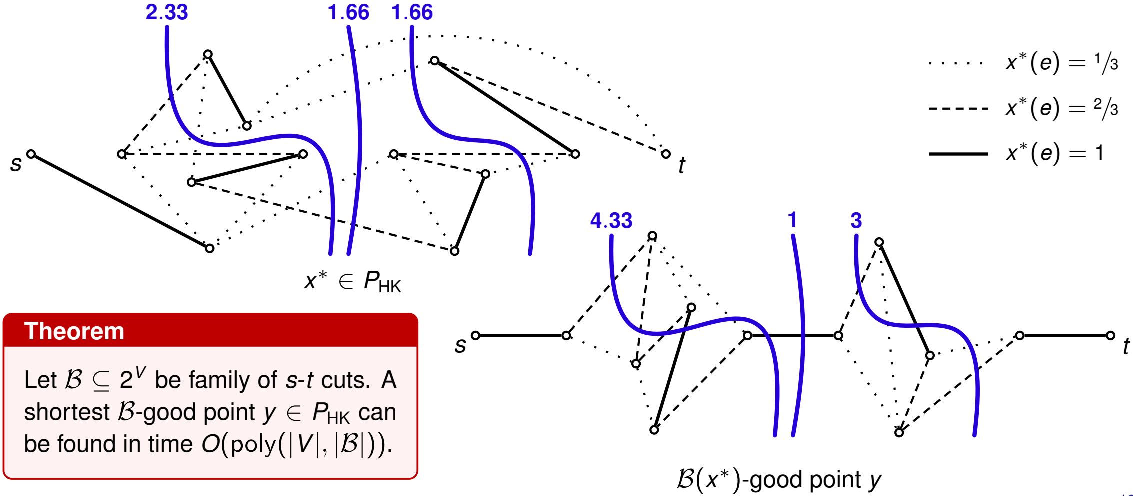
$x^* \in P_{\text{HK}}$

$\mathcal{B}(x^*)$-good point y

A new ingredient: Finding $\mathcal{B}(x^*)$-good points

$\mathcal{B}(x^*)$-good

$y \in P_{\text{HK}}$ is $\mathcal{B}(x^*)$-good: For all $B \in \mathcal{B}(x^*)$, ▶ $y(\delta(B)) \geqslant 3$, or ▶ $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.

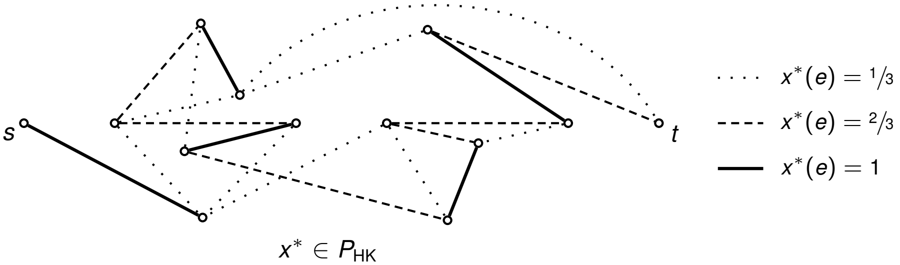


Theorem

Let $\mathcal{B} \subseteq 2^V$ be family of s-t cuts. A shortest $\mathcal{B}$-good point $y \in P_{\text{HK}}$ can be found in time $O(\text{poly}(|V|, |\mathcal{B}|))$.

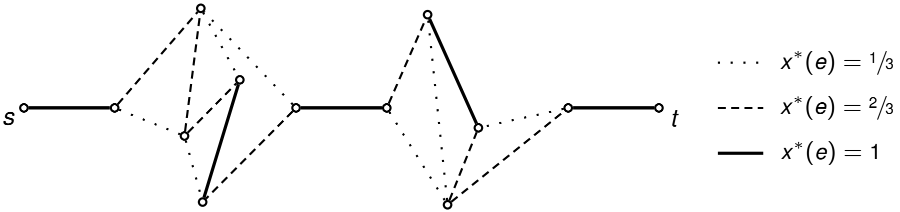
From short $\mathcal{B}(x^*)$-good points to 1.5-approx.

1. Let $x^* \in \mathrm{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be a shortest spanning tree in $(V, \mathrm{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.



From short $\mathcal{B}(x^*)$-good points to 1.5-approx.

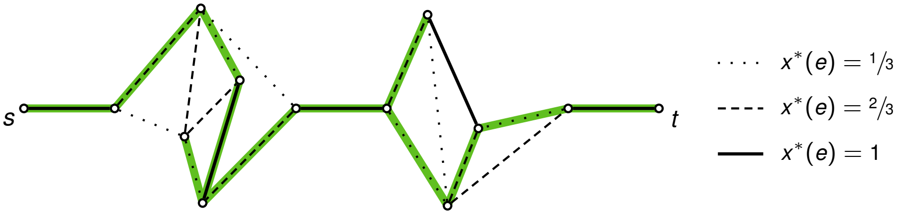
1. Let $x^* \in \text{argmin}\{\ell(x) \mid x \in P_{\text{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be a shortest spanning tree in $(V, \text{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.



$\mathcal{B}(x^*)$-good $y \in P_{\text{HK}}$

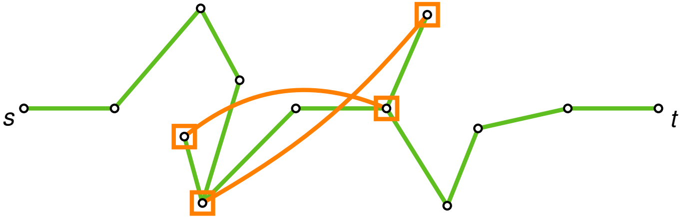
From short $\mathcal{B}(x^*)$-good points to 1.5-approx.

1. Let $x^* \in \text{argmin}\{\ell(x) \mid x \in P_{\text{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be a shortest spanning tree in $(V, \text{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.



From short $\mathcal{B}(x^*)$-good points to 1.5-approx.

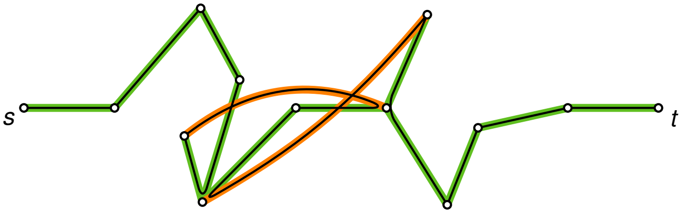
1. Let $x^* \in \text{argmin}\{\ell(x) \mid x \in P_{\text{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be a shortest spanning tree in $(V, \text{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.



$Q_T := \text{odd}(T) \triangle \{s, t\}$

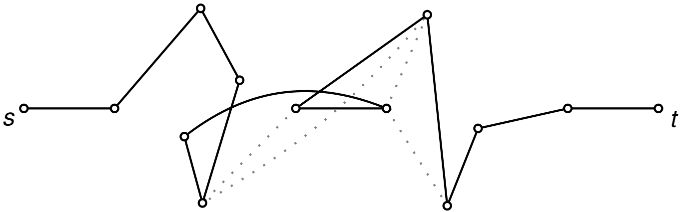
From short $\mathcal{B}(x^*)$-good points to 1.5-approx.

1. Let $x^* \in \operatorname{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be a shortest spanning tree in $(V, \operatorname{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.



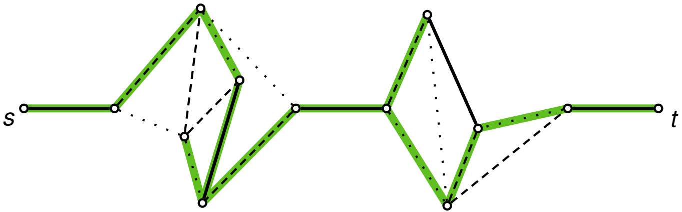
From short $\mathcal{B}(x^*)$-good points to 1.5-approx.

1. Let $x^* \in \text{argmin}\{\ell(x) \mid x \in P_{\text{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be a shortest spanning tree in $(V, \text{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.



The spanning tree T is cheap: $\ell(T) \leqslant \ell(\mathrm{OPT})$

1. Let $x^* \in \operatorname{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let ***T*** be an MST in $(V, \mathrm{supp}(y))$.
4. Let ***J*** be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of ***T*** and ***J***.

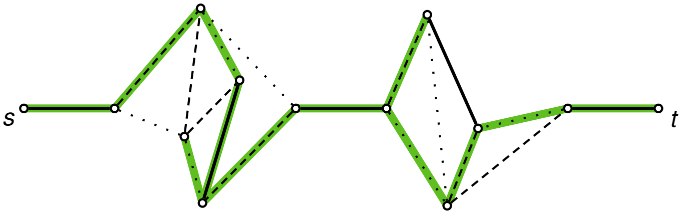


The spanning tree T is cheap: $\ell(T) \leqslant \ell(\mathrm{OPT})$

- We have $y \in P_{\mathrm{HK}} \subseteq P_{\mathrm{ST}}$.

$$\implies \ell(T) \leqslant \ell^\top y \ .$$

1. Let $x^* \in \operatorname{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be an MST in $(V, \operatorname{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.



The spanning tree T is cheap: $\ell(T) \leqslant \ell(\mathrm{OPT})$

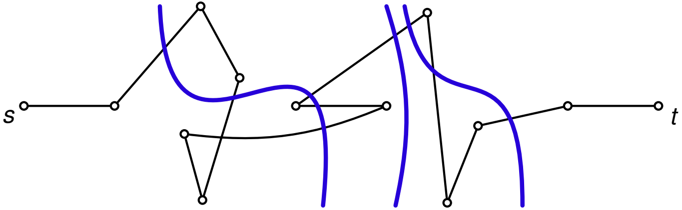
- We have $y \in P_{\mathrm{HK}} \subseteq P_{\mathrm{ST}}$.

$$\implies \ell(\boldsymbol{T}) \leqslant \ell^\top y \ .$$

- OPT is $\mathcal{B}$-good for any family $\mathcal{B}$ of s-t cuts.

$$\implies \ell^\top y \leqslant \ell(\mathrm{OPT}) \ .$$

1. Let $x^* \in \operatorname{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let $\boldsymbol{T}$ be an MST in $(V, \operatorname{supp}(y))$.
4. Let $\boldsymbol{J}$ be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of $\boldsymbol{T}$ and $\boldsymbol{J}$.



The spanning tree T is cheap: $\ell(T) \leqslant \ell(\mathrm{OPT})$

- We have $y \in P_{\mathrm{HK}} \subseteq P_{\mathrm{ST}}$.

$$\implies \ell(T) \leqslant \ell^\top y \ .$$

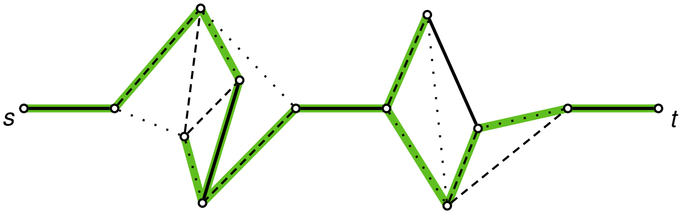
- OPT is $\mathcal{B}$-good for any family $\mathcal{B}$ of s-t cuts.

$$\implies \ell^\top y \leqslant \ell(\mathrm{OPT}) \ .$$

- Together, we conclude

$$\ell(T) \leqslant \ell^\top y \leqslant \ell(\mathrm{OPT}) \ .$$

1. Let $x^* \in \mathrm{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be an MST in $(V, \mathrm{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.



The Q_T-join J is cheap: $\ell(J) \leqslant \frac{1}{2}\ell(\mathrm{OPT})$

- We show $\frac{1}{4}x^* + \frac{1}{4}y \in P^{\uparrow}_{Q_T\text{-join}}$.

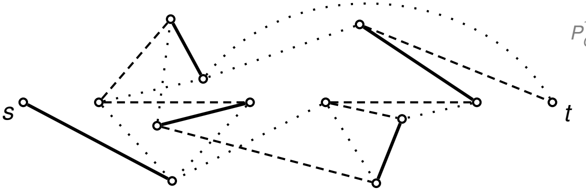
$$\Longrightarrow \ \ell(J) \leqslant \frac{1}{4}\left(\ell^\top x^* + \ell^\top y\right) \leqslant \frac{1}{2}\ell(\mathrm{OPT}) \ .$$

- Distinguish cases:

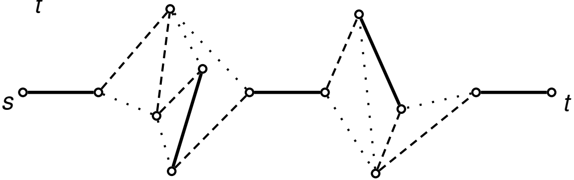
1. **2.** **3.** **4.**

1. Let $x^* \in \operatorname{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be an MST in $(V, \operatorname{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.

$$P^{\uparrow}_{Q_T\text{-join}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; x(\delta(C)) \geqslant 1 \quad \forall C \subseteq V,\ |C \cap Q_T| \text{ odd} \right\}$$



$x^* \in P_{\mathrm{HK}}$



$\mathcal{B}(x^*)$-good point y

The Q_T-join J is cheap: $\ell(J) \leqslant \frac{1}{2}\ell(\mathrm{OPT})$

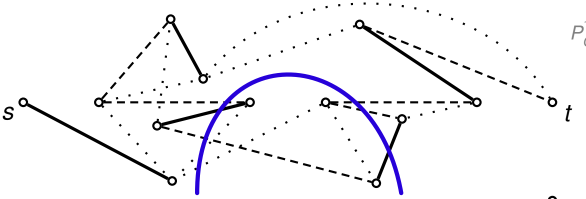
- We show $\frac{1}{4}x^* + \frac{1}{4}y \in P^{\uparrow}_{Q_T\text{-join}}$.

$$\implies \ell(J) \leqslant \frac{1}{4}\left(\ell^\top x^* + \ell^\top y\right) \leqslant \frac{1}{2}\ell(\mathrm{OPT}) \ .$$

- Distinguish cases:

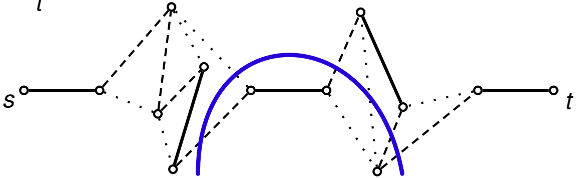
1. Non s-t cuts. **2.** **3.** **4.**

1. Let $x^* \in \operatorname{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be an MST in $(V, \mathrm{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.



$x^* \in P_{\mathrm{HK}}$

$$P^{\uparrow}_{Q_T\text{-join}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \,\middle|\, x(\delta(C)) \geqslant 1 \quad \forall C \subseteq V,\ |C \cap Q_T| \text{ odd} \right\}$$



$\mathcal{B}(x^*)$-good point y

$x^*(\delta(\boldsymbol{B})) \geqslant 2$ $\qquad$ $y(\delta(\boldsymbol{B})) \geqslant 2$

The Q_T-join J is cheap: $\ell(J) \leqslant \frac{1}{2}\ell(\mathrm{OPT})$

- We show $\frac{1}{4}x^* + \frac{1}{4}y \in P^{\uparrow}_{Q_T\text{-join}}$.

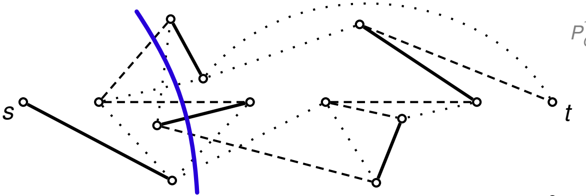
$$\implies \ell(J) \leqslant \frac{1}{4}\left(\ell^\top x^* + \ell^\top y\right) \leqslant \frac{1}{2}\ell(\mathrm{OPT}) \ .$$

- Distinguish cases:

1. **2. s-t cuts not in $\mathcal{B}(x^*)$.** **3.** **4.**

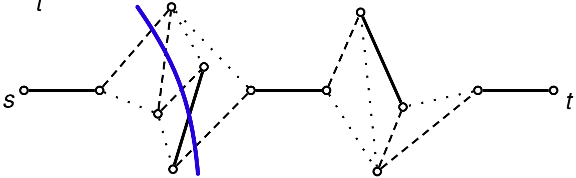
1. Let $x^* \in \operatorname{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be an MST in $(V, \operatorname{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.

$$P^{\uparrow}_{Q_T\text{-join}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \,\middle|\, x(\delta(C)) \geqslant 1 \quad \forall C \subseteq V,\ |C \cap Q_T| \text{ odd} \right\}$$



$x^* \in P_{\mathrm{HK}}$

$x^*(\delta(\boldsymbol{B})) \geqslant 3$ $\qquad$ $y(\delta(\boldsymbol{B})) \geqslant 1$



$\mathcal{B}(x^*)$-good point y

The Q_T-join J is cheap: $\ell(J) \leqslant \frac{1}{2}\ell(\mathrm{OPT})$

- We show $\frac{1}{4}x^* + \frac{1}{4}y \in P^{\uparrow}_{Q_T\text{-join}}$.

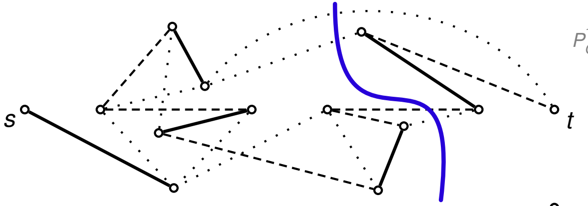
$$\Longrightarrow \ \ell(J) \leqslant \frac{1}{4}\left(\ell^\top x^* + \ell^\top y\right) \leqslant \frac{1}{2}\ell(\mathrm{OPT}) \ .$$

- Distinguish cases:

1. **2.** **3.** s-t cuts $B \in \mathcal{B}(x^*)$ with $y(\delta(B)) \geqslant 3$. **4.**

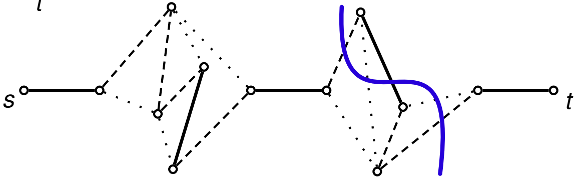
1. Let $x^* \in \operatorname{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be an MST in $(V, \operatorname{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.

$$P^{\uparrow}_{Q_T\text{-join}} = \left\{x \in \mathbb{R}^E_{\geqslant 0} \,\middle|\, x(\delta(C)) \geqslant 1 \quad \forall C \subseteq V,\ |C \cap Q_T| \text{ odd}\right\}$$



$x^* \in P_{\mathrm{HK}}$

$x^*(\delta(\boldsymbol{B})) \geqslant 1$ $\qquad$ $y(\delta(\boldsymbol{B})) \geqslant 3$



$\mathcal{B}(x^*)$-good point y

The Q_T-join J is cheap: $\ell(J) \leqslant \frac{1}{2}\ell(\mathrm{OPT})$

- We show $\frac{1}{4}x^* + \frac{1}{4}y \in P^{\uparrow}_{Q_T\text{-join}}$.

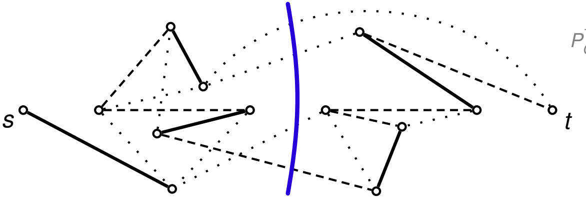
$$\implies \ell(J) \leqslant \frac{1}{4}\left(\ell^\top x^* + \ell^\top y\right) \leqslant \frac{1}{2}\ell(\mathrm{OPT}) \ .$$

- Distinguish cases:

1. **2.** **3.** **4.** s-t cuts $B \in \mathcal{B}(x^*)$ with $y(\delta(B)) = 1$.

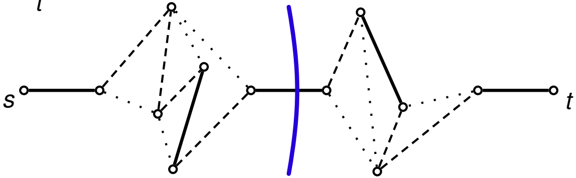
1. Let $x^* \in \operatorname{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be an MST in $(V, \mathrm{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.

$$P^{\uparrow}_{Q_T\text{-join}} = \left\{ x \in \mathbb{R}^E_{\geqslant 0} \,\middle|\, x(\delta(C)) \geqslant 1 \quad \forall C \subseteq V,\ |C \cap Q_T| \text{ odd} \right\}$$



$x^* \in P_{\mathrm{HK}}$

$x^*(\delta(\boldsymbol{B})) \geqslant 1$ $\qquad$ $y(\delta(\boldsymbol{B})) = 1$



$\mathcal{B}(x^*)$-good point y

The Q_T-join J is cheap: $\ell(J) \leqslant \frac{1}{2}\ell(\mathrm{OPT})$

- We show $\frac{1}{4}x^* + \frac{1}{4}y \in P^{\uparrow}_{Q_T\text{-join}}$.

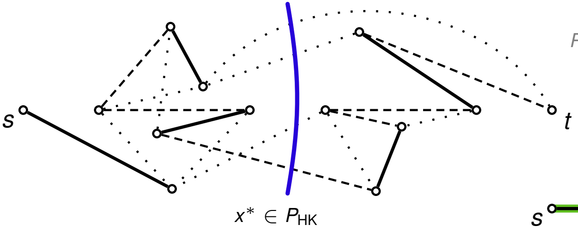
$$\Longrightarrow \ell(J) \leqslant \frac{1}{4}\left(\ell^\top x^* + \ell^\top y\right) \leqslant \frac{1}{2}\ell(\mathrm{OPT}) \ .$$

- Distinguish cases:

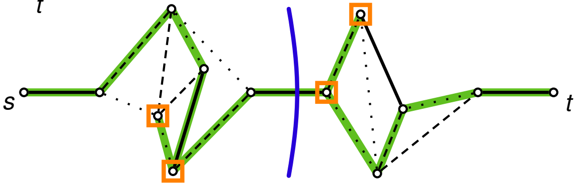
1. **2.** **3.** **4.** s-t cuts $B \in \mathcal{B}(x^*)$ with $y(\delta(B)) = 1$.

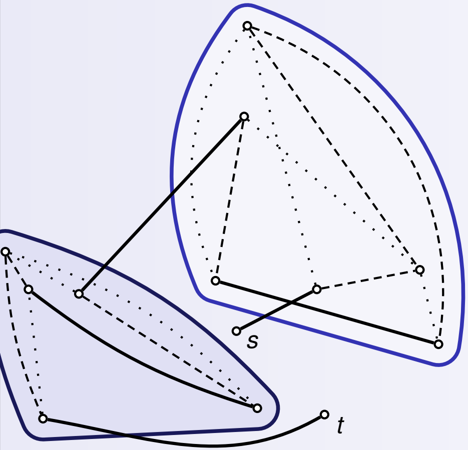
1. Let $x^* \in \operatorname{argmin}\{\ell(x) \mid x \in P_{\mathrm{HK}}\}$.
2. Let y be a shortest $\mathcal{B}(x^*)$-good point.
3. Let T be an MST in $(V, \operatorname{supp}(y))$.
4. Let J be a shortest Q_T-join.
5. Return shortcutted tour in multiunion of T and J.

$$P^{\uparrow}_{Q_T\text{-join}} = \left\{x \in \mathbb{R}^E_{\geqslant 0} \;\middle|\; x(\delta(C)) \geqslant 1 \quad \forall C \subseteq V,\ |C \cap Q_T| \text{ odd}\right\}$$



$|B \cap Q_T|$ even





The dynamic program

The DP: Finding shortest $\mathcal{B}$-good points

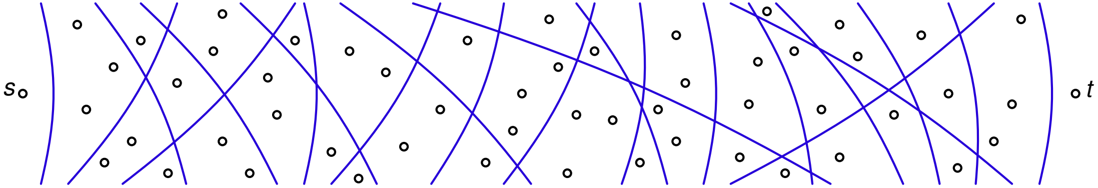
Theorem

Let $\mathcal{B} \subseteq 2^V$ a family of s-t cuts. A shortest $\mathcal{B}$-good point $y \in P_{\text{HK}}$ can be found in time $O(\text{poly}(|V|, |\mathcal{B}|))$.

$\mathcal{B}$-good point y

For all $B \in \mathcal{B}$, either

- $y(\delta(B)) \geqslant 3$, or
- $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.



The DP: Finding shortest $\mathcal{B}$-good points

Theorem

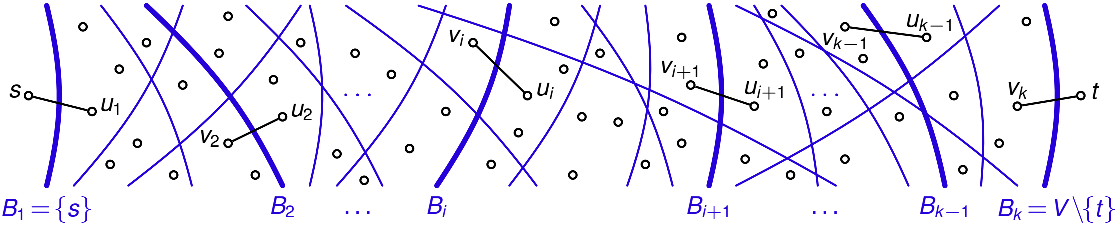
Let $\mathcal{B} \subseteq 2^V$ a family of s-t cuts. A shortest $\mathcal{B}$-good point $y \in P_{\text{HK}}$ can be found in time $O(\text{poly}(|V|, |\mathcal{B}|))$.

$\mathcal{B}$-good point y

For all $B \in \mathcal{B}$, either

- $y(\delta(B)) \geqslant 3$, or
- $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.

- DP can be interpreted as a simplified version of the one used by Traub & Vygen [SODA 2018].
- Key plan:
  - "Guess" cuts $B_1, \ldots, B_k \in \mathcal{B}$ with $y(\delta(B_i)) = 1$, and the single edge in these cuts.
  - Observation: $B_1, \ldots, B_k$ must form a chain $\rightarrow$ can split into subproblems on $B_{i+1} \setminus B_i$.



The DP: Finding shortest $\mathcal{B}$-good points

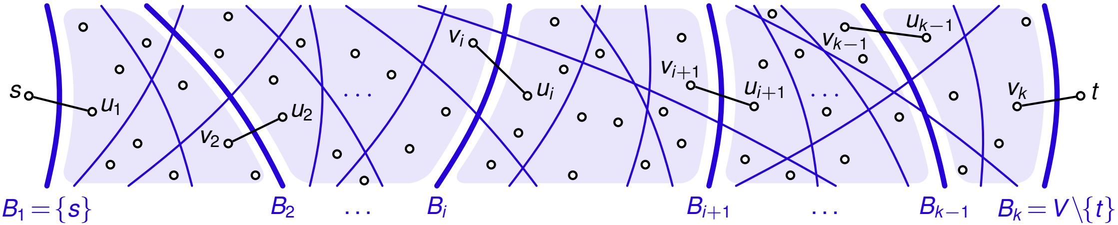
Theorem

Let $\mathcal{B} \subseteq 2^V$ a family of s-t cuts. A shortest $\mathcal{B}$-good point $y \in P_{\text{HK}}$ can be found in time $O(\text{poly}(|V|, |\mathcal{B}|))$.

$\mathcal{B}$-good point y

For all $B \in \mathcal{B}$, either
- $y(\delta(B)) \geqslant 3$, or
- $y(\delta(B)) = 1$ and y is $0/1$ on $\delta(B)$.

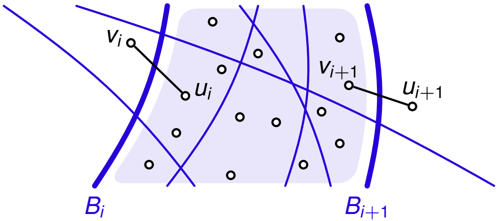
- DP can be interpreted as a simplified version of the one used by Traub & Vygen [SODA 2018].
- Key plan:
  - "Guess" cuts $B_1, \ldots, B_k \in \mathcal{B}$ with $y(\delta(B_i)) = 1$, and the single edge in these cuts.
  - Observation: $B_1, \ldots, B_k$ must form a chain $\rightarrow$ can split into subproblems on $B_{i+1} \setminus B_i$.



Solving a single subproblem

- Restriction to $B_{i+1} \setminus B_i$, start at u_i, end at v_{i+1}.
- Enforce $y(\delta(B)) \geqslant 3$ for $B \in \mathcal{B}$ with $B_i \subsetneq B \subsetneq B_{i+1}$.
- Corresponding LP formulation:

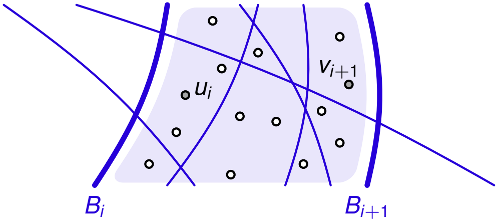
$$\lambda(B_{i+1} \setminus B_i, u_i, v_{i+1}) = \min \ell^\top y$$
$$y \in P_{\mathrm{HK}}(B_{i+1} \setminus B_i, u_i, v_{i+1})$$
$$y(\delta(B)) \geqslant 3 \qquad \forall B \in \mathcal{B} \colon B_i \subsetneq B \subsetneq B_{i+1} \ .$$



Solving a single subproblem

- Restriction to $B_{i+1} \setminus B_i$, start at u_i, end at v_{i+1}.
- Enforce $y(\delta(B)) \geqslant 3$ for $B \in \mathcal{B}$ with $B_i \subsetneq B \subsetneq B_{i+1}$.
- Corresponding LP formulation:

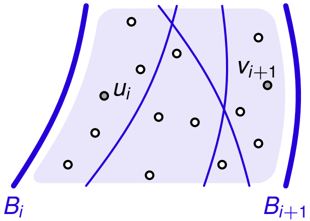
$$\begin{aligned}\lambda(B_{i+1} \setminus B_i, u_i, v_{i+1}) = \min\ & \ell^\top y \\ & y \in P_{\mathrm{HK}}(B_{i+1} \setminus B_i, u_i, v_{i+1}) \\ & y(\delta(B)) \geqslant 3 \qquad \forall B \in \mathcal{B} \colon B_i \subsetneq B \subsetneq B_{i+1} \quad .\end{aligned}$$



Solving a single subproblem

- Restriction to $B_{i+1} \setminus B_i$, start at u_i, end at v_{i+1}.
- Enforce $y(\delta(B)) \geqslant 3$ for $B \in \mathcal{B}$ with $B_i \subsetneq B \subsetneq B_{i+1}$.
- Corresponding LP formulation:

$$\begin{aligned} \lambda(B_{i+1} \setminus B_i, u_i, v_{i+1}) = \min\ & \ell^\top y \\ & y \in P_{\mathrm{HK}}(B_{i+1} \setminus B_i, u_i, v_{i+1}) \\ & y(\delta(B)) \geqslant 3 \qquad \forall B \in \mathcal{B} \colon B_i \subsetneq B \subsetneq B_{i+1} \end{aligned} .$$



- Idea: Advance from one cut B with $y(\delta(B)) = 1$ to another.

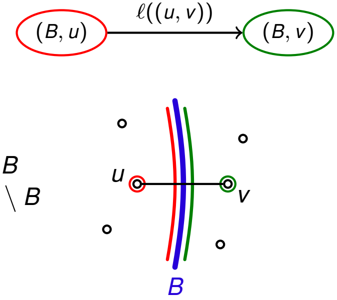
Setting up the DP

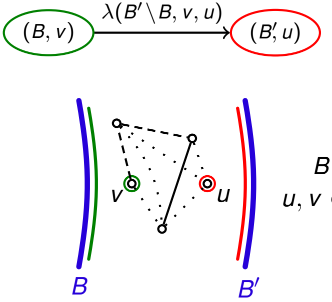
- Idea: Advance from one cut B with $y(\delta(B)) = 1$ to another.
- Formulation as a shortest path problem on auxiliary digraph:

Nodes: Pairs (B, v) for $B \in \mathcal{B}$ and $v \in V$.

Edges: Two types of steps corresponding to extension of a solution.

$u \in B$
$v \in V \setminus B$





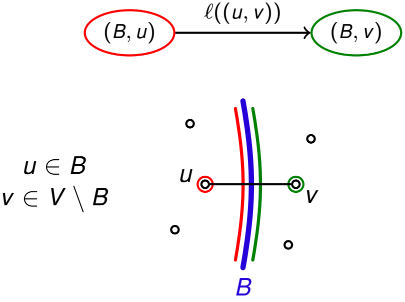
$B \subsetneq B'$
$u, v \in B' \setminus B$

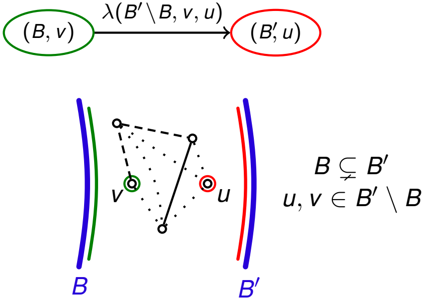
Setting up the DP

- Idea: Advance from one cut B with $y(\delta(B)) = 1$ to another.
- Formulation as a shortest path problem on auxiliary digraph:

  Nodes: Pairs (B, v) for $B \in \mathcal{B}$ and $v \in V$.

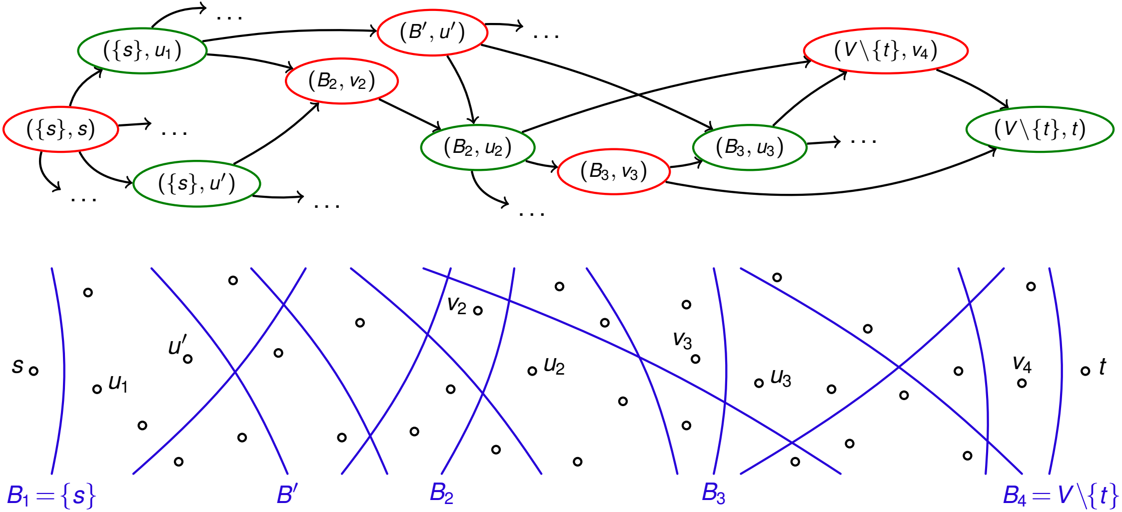
  Edges: Two types of steps corresponding to extension of a solution.



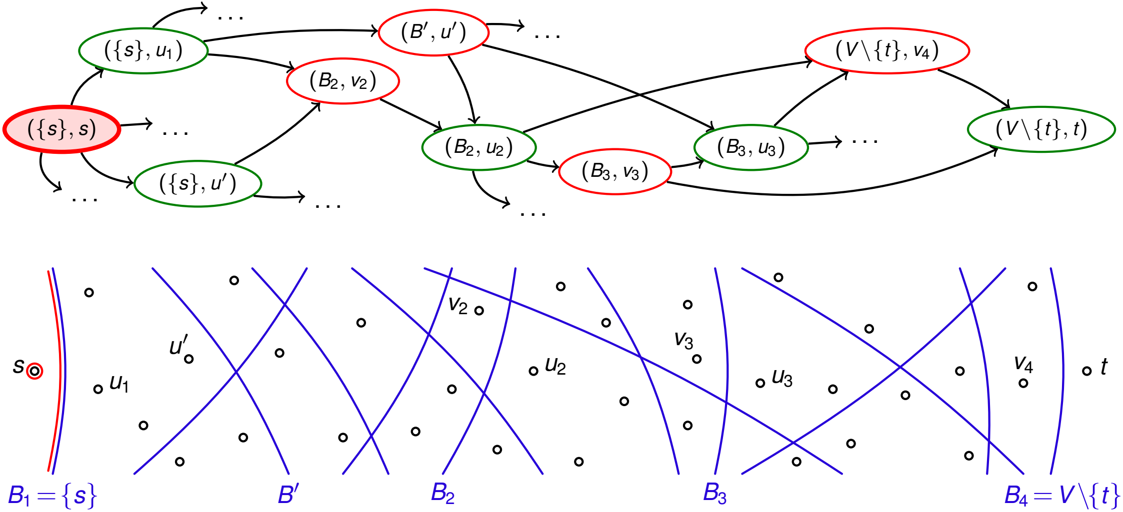


- Optimal solution: Shortest $(\{s\}, s)$-$(V \setminus \{t\}, t)$ path in auxiliary digraph.

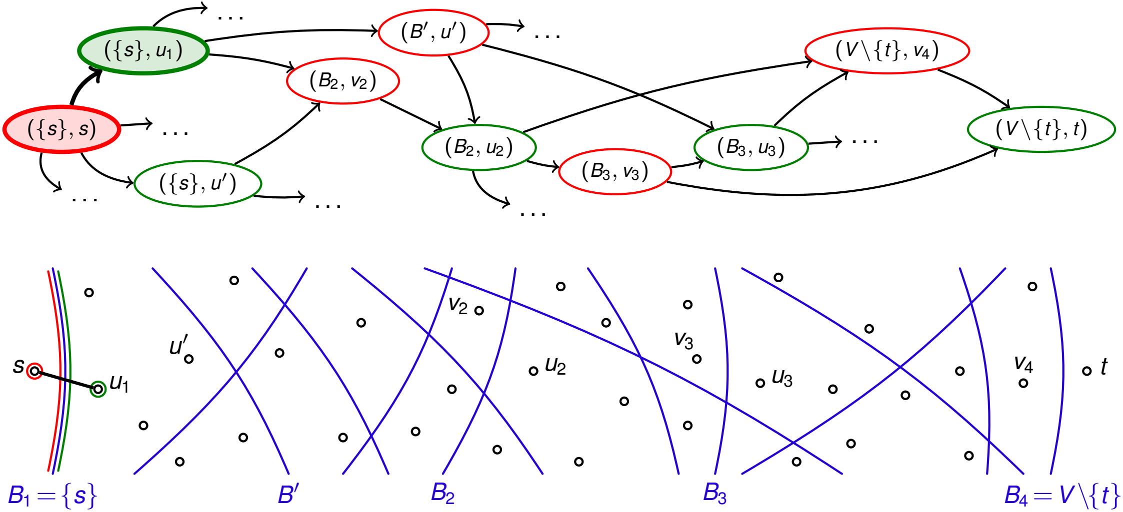
DP auxiliary graph: An example



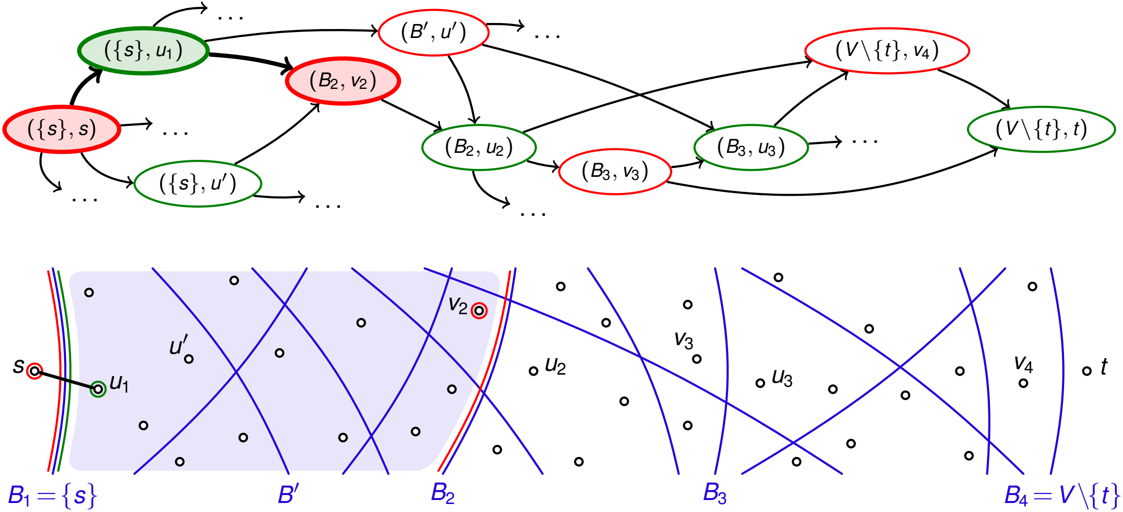
DP auxiliary graph: An example



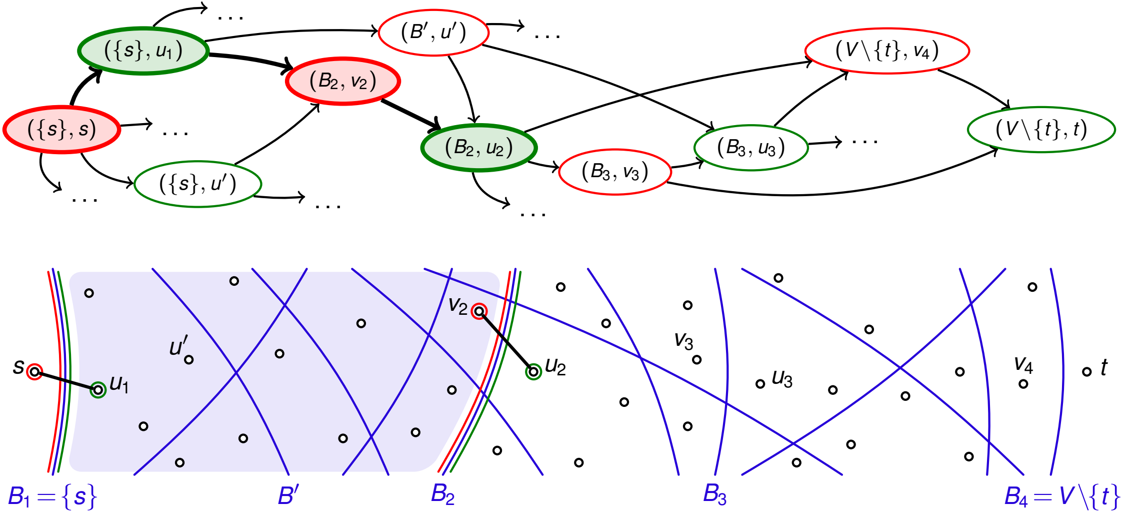
DP auxiliary graph: An example



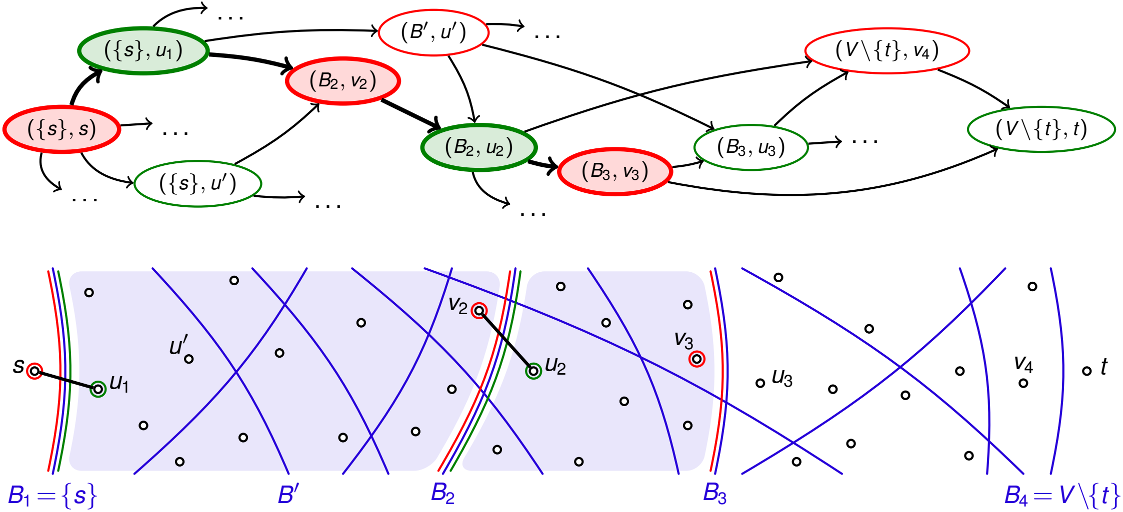
DP auxiliary graph: An example



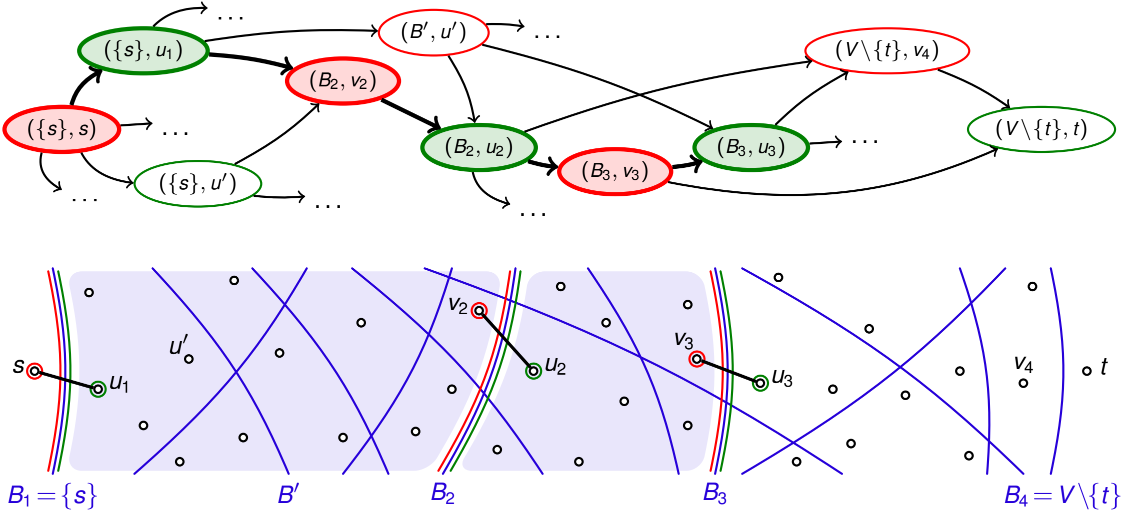
DP auxiliary graph: An example



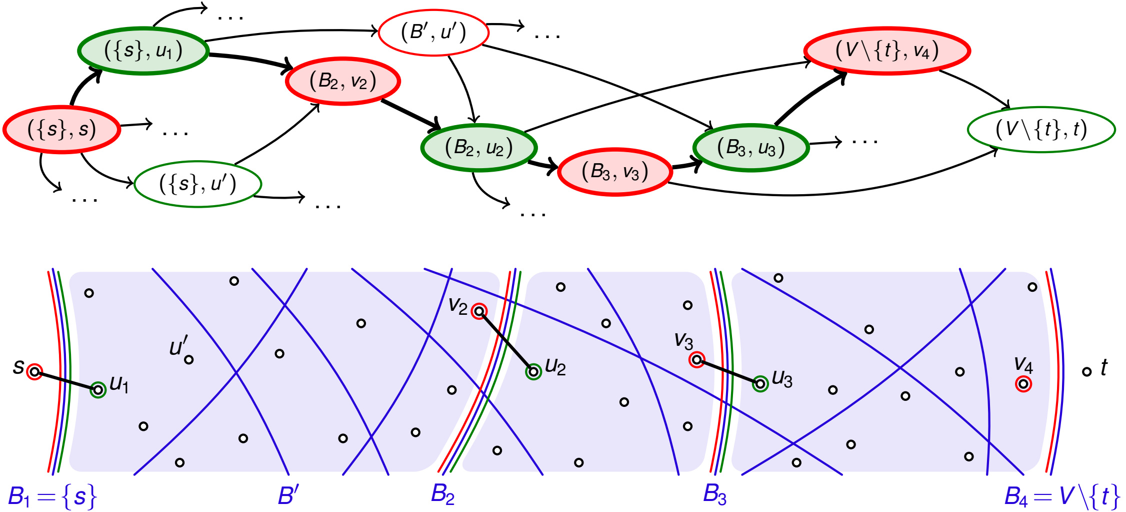
DP auxiliary graph: An example



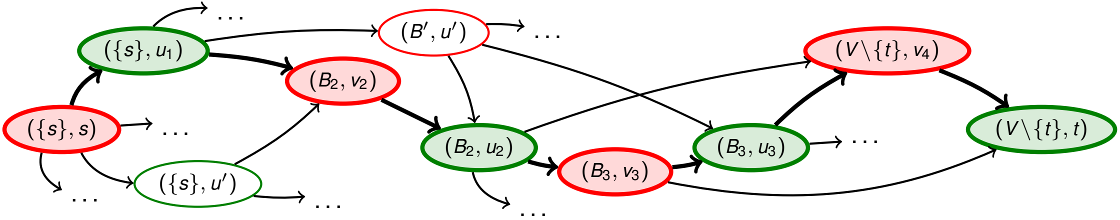
DP auxiliary graph: An example

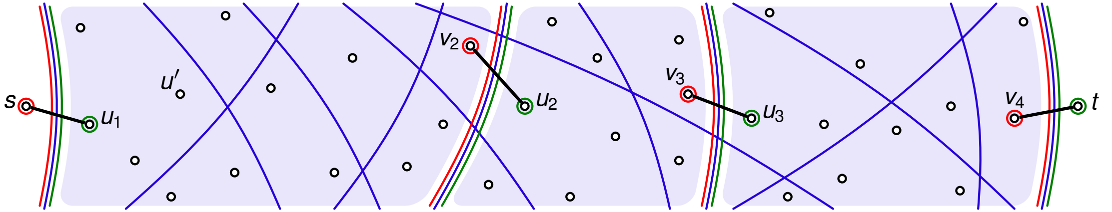


DP auxiliary graph: An example



DP auxiliary graph: An example





Getting solutions through DP

- DP solution: Combination of edges in cuts B_i and partial Held-Karp solutions.

Getting solutions through DP

- DP solution: Combination of edges in cuts B_i and partial Held-Karp solutions.

Theorem (basic properties of DP solutions)

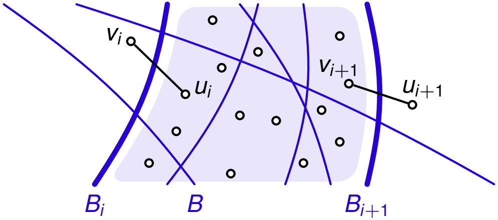
- Any DP solution is in P_{HK}.
- Any DP solution is $\mathcal{B}$-good.
- Shortest DP solution has length at most $\ell(\mathrm{OPT})$.

- DP solution: Combination of edges in cuts B_i and partial Held-Karp solutions.

Theorem (basic properties of DP solutions)

- Any DP solution is in P_{HK}.
- Any DP solution is $\mathcal{B}$-good.
- Shortest DP solution has length at most $\ell(\text{OPT})$.

Proof. Need to show that DP solution y satisfies $y(\delta(B)) \geqslant 3$ for all $B \in \mathcal{B} \setminus \{B_1, \ldots, B_k\}$.



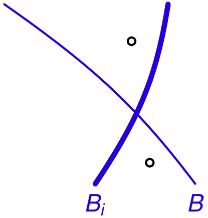
Getting solutions through DP

- DP solution: Combination of edges in cuts B_i and partial Held-Karp solutions.

Theorem (basic properties of DP solutions)

- Any DP solution is in P_{HK}.
- Any DP solution is $\mathcal{B}$-good.
- Shortest DP solution has length at most $\ell(\mathrm{OPT})$.

Proof. Need to show that DP solution y satisfies $y(\delta(B)) \geqslant 3$ for all $B \in \mathcal{B} \setminus \{B_1, \ldots, B_k\}$.



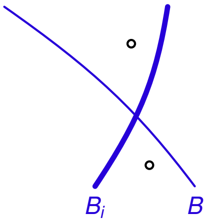
Getting solutions through DP

- DP solution: Combination of edges in cuts B_i and partial Held-Karp solutions.

Theorem (basic properties of DP solutions)

- Any DP solution is in P_{HK}.
- Any DP solution is $\mathcal{B}$-good.
- Shortest DP solution has length at most $\ell(\text{OPT})$.

Proof. Need to show that DP solution y satisfies $y(\delta(B)) \geqslant 3$ for all $B \in \mathcal{B} \setminus \{B_1, \ldots, B_k\}$.



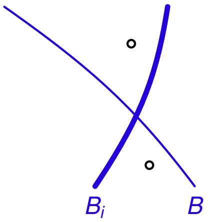
$y(\delta(B))$

- DP solution: Combination of edges in cuts B_i and partial Held-Karp solutions.

Theorem (basic properties of DP solutions)

- Any DP solution is in P_{HK}.
- Any DP solution is $\mathcal{B}$-good.
- Shortest DP solution has length at most $\ell(\mathrm{OPT})$.

Proof. Need to show that DP solution y satisfies $y(\delta(B)) \geqslant 3$ for all $B \in \mathcal{B} \setminus \{B_1, \ldots, B_k\}$.



$$y(\delta(B)) + y(\delta(B_i))$$

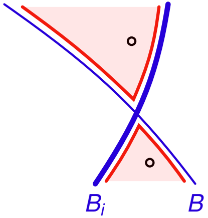
Getting solutions through DP

- DP solution: Combination of edges in cuts B_i and partial Held-Karp solutions.

Theorem (basic properties of DP solutions)

- Any DP solution is in P_{HK}.
- Any DP solution is $\mathcal{B}$-good.
- Shortest DP solution has length at most $\ell(\text{OPT})$.

Proof. Need to show that DP solution y satisfies $y(\delta(B)) \geqslant 3$ for all $B \in \mathcal{B} \setminus \{B_1, \ldots, B_k\}$.



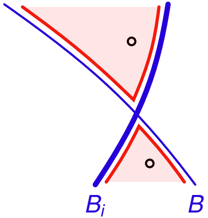
$$y(\delta(B)) + y(\delta(B_i)) \geqslant y(\delta(B_i \setminus B)) + y(\delta(B \setminus B_i))$$

- DP solution: Combination of edges in cuts B_i and partial Held-Karp solutions.

Theorem (basic properties of DP solutions)

- Any DP solution is in P_{HK}.
- Any DP solution is $\mathcal{B}$-good.
- Shortest DP solution has length at most $\ell(\mathrm{OPT})$.

Proof. Need to show that DP solution y satisfies $y(\delta(B)) \geqslant 3$ for all $B \in \mathcal{B} \setminus \{B_1, \ldots, B_k\}$.



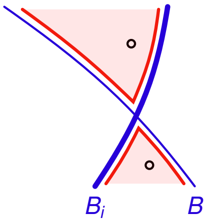
$$y(\delta(B)) + \underbrace{y(\delta(B_i))}_{=1} \geqslant \underbrace{y(\delta(B_i \setminus B))}_{\geqslant 2} + \underbrace{y(\delta(B \setminus B_i))}_{\geqslant 2}$$

- DP solution: Combination of edges in cuts B_i and partial Held-Karp solutions.

Theorem (basic properties of DP solutions)

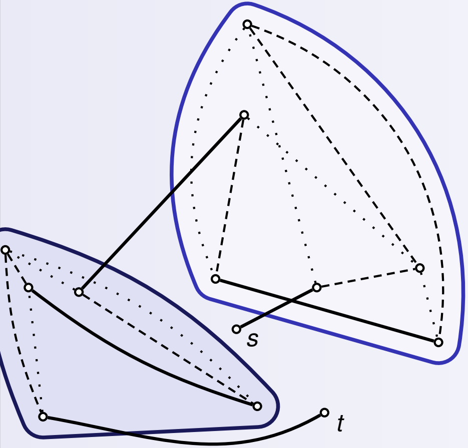
- Any DP solution is in P_{HK}.
- Any DP solution is $\mathcal{B}$-good.
- Shortest DP solution has length at most $\ell(\text{OPT})$.

Proof. Need to show that DP solution y satisfies $y(\delta(B)) \geqslant 3$ for all $B \in \mathcal{B} \setminus \{B_1, \ldots, B_k\}$.



$$y(\delta(B)) + \underbrace{y(\delta(B_i))}_{=1} \geqslant \underbrace{y(\delta(B_i \setminus B))}_{\geqslant 2} + \underbrace{y(\delta(B \setminus B_i))}_{\geqslant 2}$$

$$\Longrightarrow \quad y(\delta(B)) \geqslant 3 \ .$$



Conclusions

Theorem **[Zenklusen, 2018]**

There is a 1.5-approximation for path TSP.

- Approximation factors below 1.5 for TSP (or even path TSP)?
- Show that the integrality gap of Held-Karp relaxation for path TSP is 1.5.