

# NONREPETITIVE COLORING

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(joint work with V. Dujmović, G. Joret, B. Walczak, and D. Wood)

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## SQUARE-FREE WORDS

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## **Observation** (Alon, Grytczuk, Haluszczak, Riordan 2002)

If  $G$  is a tree, then  $\pi(G) \leq 4$ .

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Any path has a 4-coloring such that every repetitive walk is boring.

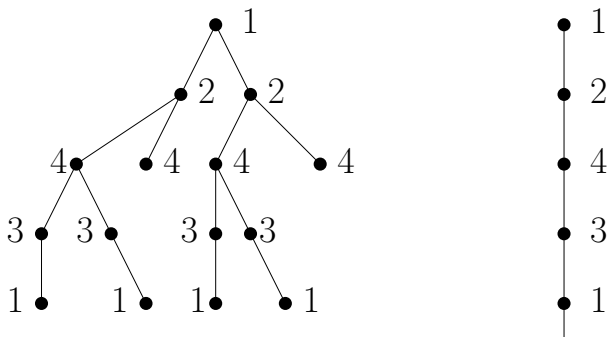
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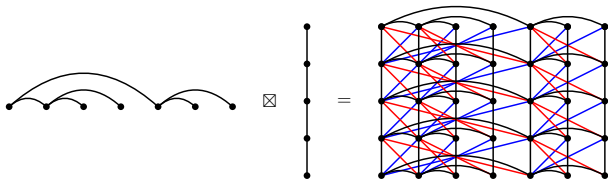
Do planar graphs have bounded nonrepetitive chromatic number?



# A PRODUCT STRUCTURE IN PLANAR GRAPHS

denoted by  $A \boxtimes B$ , is the graph with vertex set  $V(A) \times V(B)$ , where distinct vertices  $(v, x), (w, y) \in V(A) \times V(B)$  are adjacent if

- $v = w$  and  $xy \in E(B)$ , or
- $x = y$  and  $vw \in E(A)$ , or
- $vw \in E(A)$  and  $xy \in E(B)$ .



**Theorem** (Dujmović, Joret, Micek, Morin, Ueckerdt, Wood 2019)

Every planar graph is a subgraph of  $H \boxtimes P$  for some graph  $H$  with treewidth at most 8 and some path  $P$ .

## STRONGLY NONREPETITIVE COLORING

A coloring is **strongly nonrepetitive** if for any repetitive walk  $v_1, \dots, v_{2t}$ , there exist  $1 \leq i \leq t$  such that  $v_i = v_{i+t}$ .

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**Theorem** (Kündgen, Pelsmajer 2008)

If  $G$  has treewidth  $k$ , then  $\pi^*(G) \leq 4^k$ .

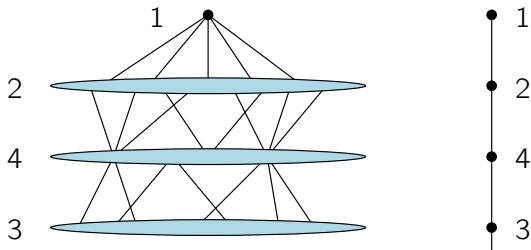
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# APPLICATION TO NONREPETITIVE COLORING

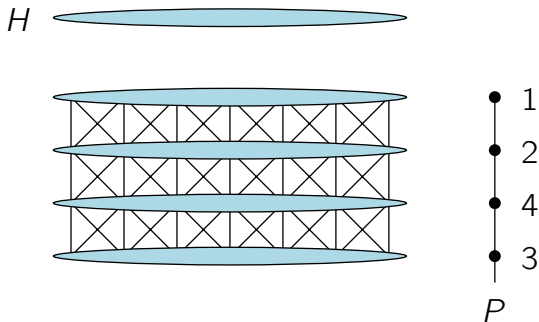
**Lemma** (Dujmović, Esperet, Joret, Walczak, Wood 2019)

For every graph  $H$  and every path  $P$ , we have  $\pi^*(H \boxtimes P) \leq 4 \pi^*(H)$ .

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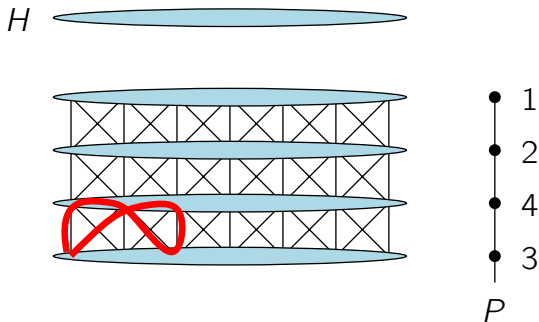
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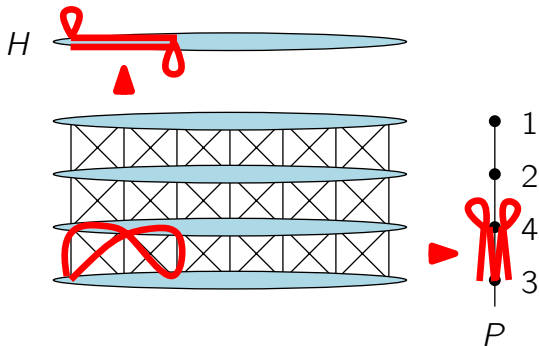




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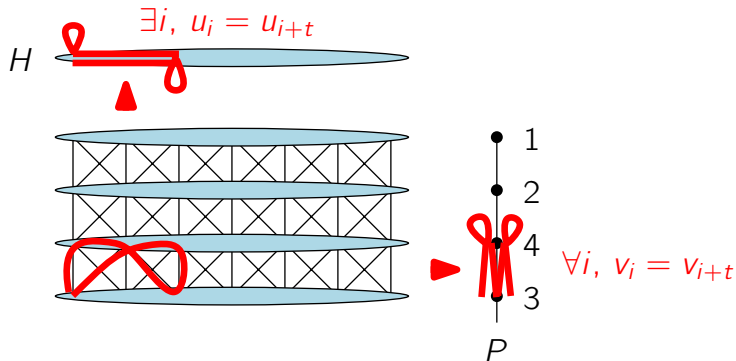
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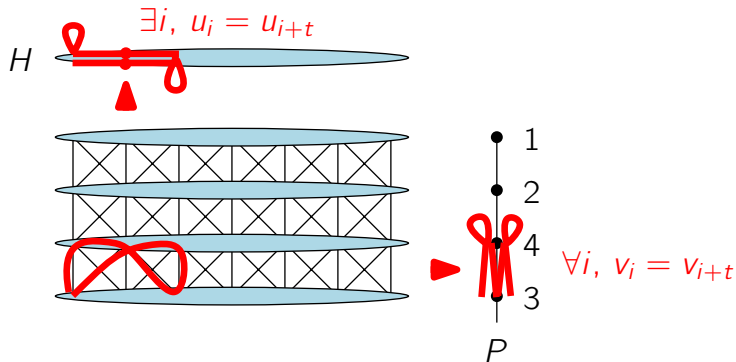
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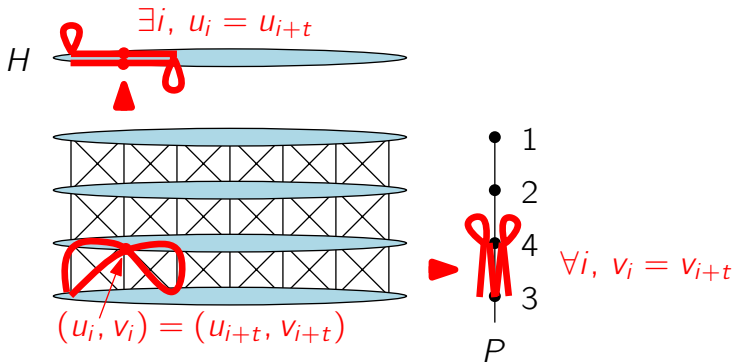
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Graphs of Euler genus  $g$  have nonrepetitive chromatic number at most  $256 \cdot \max(3, 2g)$ .



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**Theorem** (Dujmović, Esperet, Joret, Walczak, Wood 2019)

Graphs avoiding a fixed minor or topological minor have bounded non-repetitive chromatic number.

## OPEN PROBLEMS

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- + Several questions of Gwen and David about the connections between bounded expansion and nonrepetitive chromatic number.