## Remote Estimation of Two-State Markov Sources: a State Estimation Entropy Perspective

BIRS Workshop

Andrea Munari – joint work with Gianluigi Liva, Giuseppe Cocco (UPC, Spain) DLR - Institute of Communications and Navigation Banff. 2024.03.12



How to maintain an accurate knowledge from remote sensors





The challenge of maintaining an accurate knowledge from remote sensors

- monitor state of remotely-deployed sensor nodes
  - smart agriculture
  - environmental monitoring
  - asset tracking, ...
- o possibly massive number of battery-powered, low-complexity devices
  - transmit only terminals
  - constraints on protocols complexity
- sporadic traffic, following non-regular patterns
  - high cost for coordination and resource assignment (grant-based)



The challenge of maintaining an accurate knowledge from remote sensors



- o random access procedures commonly used for medium sharing
  - advanced schemes (modern random access)<sup>1</sup>
  - ALOHA-like access employed in most practical systems, e.g. LoRaWAN, Sigfox

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- o non-trivial challenges hinder performance
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- o goal is to maintain accurate knowledge of the monitored sources at the receiver
  - first step is definition of metrics that can capture this capability
  - pioneering role played by age of information (AoI)<sup>2</sup>

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### Age of Information A measure of freshness of available information





 $\sigma(t):$  time stamp of last rcvd update



- Aol focuses on timeliness of delivered information
  - in some applications, knowledge at the receiver may be critical (e.g., actuation)
- o other metrics proposed to try to capture this aspect
  - age of incorrect information<sup>3</sup>

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    - AoI matters at time instants where receiver needs to use information, e.g. actuation

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    - $_{\circ}~$  Aol matters at time instants where receiver needs to use information, e.g. actuation
  - information theory inspired metrics

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#### mutual information<sup>₅</sup>

- $I(X_t; W^t) = H(X_t) H(X_t | W^t)$
- if close to 0, received samples  $W^t$  carry little information and deemed obsolete

<sup>5.</sup> Y. Sun and B. Cyr, "Information aging through queues: A mutual information perspective," in Proc. IEEE SPAWC Workshop, 2018.

<sup>6.</sup> M. Rezaeian, B. Vo, J. S. Evans, "The optimal observability of partially observable Markov decision processes: Discrete state space," IEEE Trans. Autom. Control, 2010

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- entropy on tracked process conditioned on received samples<sup>6,7</sup>
  - $H(X_t \mid W^t)$
  - measure of uncertainty at the receiver on the current status of the source

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#### Remote monitoring in random access channels Current status and open questions



- performance well-known in terms of traditional metrics (e.g., throughput, delay)
- o recent results characterized behavior in terms of Aol
  - ALOHA-based systems<sup>8,9,10,11</sup>
  - modern random access schemes<sup>12,13</sup>
- $\circ~$  first studies for age of incorrect information  $^{\rm 14,15}$

#### o behavior in terms of other metrics largely unexplored

<sup>8.</sup> S. Kaul, R. Yates, "Status updates over unreliable multiaccess channels," in Proc. IEEE ISIT, 2017

<sup>9.</sup> R. Yates, S. Kaul, "Age of information in uncoordinated unslotted updating," in Proc. IEEE ISIT, 2020

<sup>10.</sup> X. Chen, K. Gatsis, H. Hassani, S. Bidokhti, "Age of information in random access channels," in IEEE Trans. Inf. Theory, 2022

<sup>11.</sup> O. Yavascan and E. Uysal, "Analysis of slotted ALOHA with an age threshold," IEEE J. Sel. Areas Commun., 2021.

<sup>12.</sup> A. Munari, "Modern random access: an age of information perspective on irregular repetition slotted ALOHA," IEEE Trans. Commun., 2021.

<sup>13.</sup> A. Munari, F. Lazaro, G. Durisi, G. Liva, "The dynamic behavior of frameless ALOHA: Drift analysis, throughput and age of information," IEEE Trans. Commun., 2023.

<sup>14.</sup> A. Nayak, A. Kalor, F. Chiariotti, P. Popovski, "A decentralized policy for minimization of age of incorrect information in slotted ALOHA systems," in Proc. IEEE ICC., 2023

<sup>15.</sup> A. Munari, "Monitoring IoT sources over random access channels: age of incorrect information and missed detection probability," in Proc. IEEE ICC, 2024.



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o study behavior in terms of age of information and state estimation entropy



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• analytical results via hidden Markov models and density evolution

• insights on impact of role played by these metrics in terms of protocol operations

## Outline



## (1) System Model and Preliminaries

2) State Estimation Entropy Analysis

Hidden Markov Models and APP Logarithmic ratio

Recursive Calculation of  $\lambda_n$ 

**Density Evolution Analysis** 

Numerical Results and Discussion



 $\circ~M$  independent sources, modeled as two-state Markov processes of alphabet  $\{0,1\}$ 





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- common channel to a receiver (slotted time)
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- o common channel to a receiver (slotted time)
  - at each slot, nodes decide whether to send packet containing current source value
- slotted ALOHA access, collision channel model
  - destructive collisions, singleton slots always decoded
  - no feedback nor retransmissions



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- $\ominus$ ,  $\oplus$ : observation of state (0 or 1) from another source



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• denote sequence of observations up to time n as  $Y^n = [Y_0, Y_1, \dots, Y_n]$ 

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  - observation  $Y_n \in \{\mathtt{I},\mathtt{C},\ominus,\oplus\}$  carries information on the source of interest
    - $X_{n-1} = 0$ , state known at receiver
    - $\circ~Y_n=\ominus \rightarrow X_n=0:$  a state change would have induced a collision



 $\circ\,$  for a given sequence of observations  $Y^n=y^n,$  uncertainty at the receiver on current state of the source measured by the entropy

$$\mathsf{h}(y^n) = \mathsf{H}\left(X_n \,|\, Y^n = y^n\right)$$



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single node, random transmission policy;  $q_{10} = 0.2$ ,  $q_{01} = 0.01$ , i.e.,  $\pi_0 \simeq 0.95$ ,  $\pi_1 \simeq 0.05$  $H(X) = -\pi_0 \log_2 \pi_0 - \pi_1 \log_2 \pi_1$ 



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• we consider the expected value of the r.v.  $H_n = h(Y^n)$ 

 $\mathsf{E}[H_n] = \mathsf{H}\left(X_n \,|\, Y^n\right)$ 



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 $\circ\,$  in particular, we are interested in the limiting behavior as n grows large, giving the average state estimation entropy

$$H_{\infty} = \lim_{N \to \infty} \frac{1}{N} \sum_{n=0}^{N-1} \mathsf{H}(X_n \,|\, Y^n)$$

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$$P(x_n, y^n) = \sum_{x_{n-1}} P(x_n | x_{n-1}) P(y_n | x_n) P(x_{n-1}, y^{n-1})$$

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o observation probabilities can easily be computed

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$$\mathsf{P}[Y_n = I \mid X_n = x_n] = (1 - \alpha)^M$$

• 
$$\mathsf{P}[Y_n = \ominus | X_n = x_n] = (M-1)\pi_0 \alpha (1-\alpha)^{M-1}$$

•  $\mathsf{P}[Y_n = 0 | X_n = 0] = \alpha (1 - \alpha)^{M-1}$ ,  $\mathsf{P}[Y_n = 0 | X_n = 1] = 0$ 

Random transmission strategy



- o channel access independent of source evolution
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  - observation of I,  $\ominus$ ,  $\oplus$ , C does not provide information on the state of the source



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- $S_n$ : number of sources (other than reference one) in state 0 at time n
- channel output depends on state through conditional probability function  $P(Y_n | \sigma_{n-1}, \sigma_n)$



• consider a posteriori probability (APP) logarithmic ratio

$$\lambda_n := \ln \frac{\mathsf{P}[X_n = 0 \mid Y^n = y^n]}{\mathsf{P}[X_n = 1 \mid Y^n = y^n]}$$



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$$X_n \to \Lambda_n \to Y^n$$
, i.e.  $P(x_n \mid \lambda_n, y^n) = P(x_n \mid \lambda_n)$ 



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- proof sketch: Fisher-Neyman factorization th., writing  $P(y^n | x_n) = a(x_n, \lambda_n)b(y^n)$ , with  $a(\cdot)$ ,  $b(\cdot)$  non-negative functions



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 $\circ~$  recalling that  $\mathsf{P}[X_n=0\,|\,\Lambda_n=\lambda_n]+\mathsf{P}[X_n=1\,|\,\Lambda_n=\lambda_n]=1,$  we get

$$\mathsf{P}[X_n = x_n \,|\, \Lambda_n = \lambda_n] = \frac{e^{-x_n \lambda_n}}{1 + e^{-\lambda_n}}$$



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• by the data-processing inequality

$$\mathsf{h}(y^n) = \mathsf{H}(X_n \,|\, Y^n = y^n) = \mathsf{H}(X_n \,|\, \Lambda_n = \lambda_n)$$



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- $\circ~\lambda_n$  suffices to compute uncertainty at receiver at time n
  - recursive calculation of  $\lambda_n$  as function of  $\lambda_{n-1}$  and observation  $y_n$



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Random transmission policy



 $\circ\,$  from forward recursion on HMM we have

$$P(x_n, y^n) = \sum_{x_{n-1} \in \{0,1\}} P(y_n \mid x_n) P(x_n \mid x_{n-1}) P(x_{n-1}, y^{n-1})$$

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• leaning on this, we obtain

$$\lambda_n = \ln \frac{\mathsf{P}[X_n = 0, Y^n = y^n]}{\mathsf{P}[X_n = 1, Y^n = y^n]} = \ln \frac{P(y_n \mid 0)}{P(y_n \mid 1)} + \ln \frac{q_{00} + q_{10}e^{-\lambda_{n-1}}}{q_{01} + q_{11}e^{-\lambda_{n-1}}} := f(y_n, \lambda_{n-1})$$

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- all quantities are known
  - $q_{00}$ ,  $q_{01}$ ,  $q_{10}$ ,  $q_{11}$  are source transition probabilities
  - observation probabilities  $P(y_n | x_n)$  derived earlier

Reactive transmission policy



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Reactive transmission policy



- consider now reactive transmission strategy
- o following the same approach, we obtain

$$P(\sigma_n, y^n) = \sum_{\sigma_{n-1}} P(y_n \mid \sigma_{n-1}, \sigma_n) P(\sigma_n \mid \sigma_{n-1}) P(\sigma_{n-1}, y^{n-1})$$
$$\lambda_n = \ln \frac{\sum_{\sigma_n \in \{0\} \times S} P(\sigma_n, y^n)}{\sum_{\sigma_n \in \{1\} \times S} P(\sigma_n, y^n)}$$

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$$P(\sigma_n, y^n) = \sum_{\sigma_{n-1}} P(y_n \mid \sigma_{n-1}, \sigma_n) P(\sigma_n \mid \sigma_{n-1}) P(\sigma_{n-1}, y^{n-1})$$
$$\lambda_n = \ln \frac{\sum_{\sigma_n \in \{0\} \times S} P(\sigma_n, y^n)}{\sum_{\sigma_n \in \{1\} \times S} P(\sigma_n, y^n)}$$

•  $P(y_n | \sigma_{n-1}, \sigma_n)$ ,  $P(\sigma_{n-1}, \sigma_n)$  obtained from statistics of  $X_n$  and access strategy

Reactive transmission policy



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• however, complexity of calculation grows with  $M^2$ : impractical for larger networks



 $\circ$  to simplify calculation of  $\lambda_n$ , introduce myopic surrogate model



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  - reference source operates following reactive strategy



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 $\circ\,$  under this hypothesis, observation no longer depends on  $S_n$ 


## Recursive calculation of $\lambda_n,$ reactive case $_{\rm Myopic\ surrogate\ model}$



 $\circ\,$  for the myopic HMM, the recursion simplifies to

$$P(x_n, y^n) = \sum_{x_{n-1} \in \{0,1\}} P(y_n \,|\, x_{n-1}, x_n) P(x_n \,|\, x_{n-1}) P(x_{n-1}, y^{n-1})$$

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- complexity independent of the number of sources
- HMM gives in general an approximation, yet exact for symmetric sources  $(q_{10} = q_{01})$

#### Outline



#### (1) System Model and Preliminaries

#### 2 State Estimation Entropy Analysis

Hidden Markov Models and APP Logarithmic ratio

Recursive Calculation of  $\lambda_n$ 

Density Evolution Analysis

Numerical Results and Discussion

#### State Estimation Entropy Derivation Statistics of the APP logarithmic ratio



o to compute state estimation entropy, we are interested in

$$\mathsf{E}[H_n] = \mathsf{H}(X_n \,|\, Y^n) = \mathsf{H}(X_n \,|\, \Lambda_n) = \mathsf{E}[\mathsf{h}(\Lambda_n)]$$

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- $\circ\,$  analytical derivation of SEE available via probability distribution of r.v.  $\Lambda_n$
- $\circ\,$  quantized density evolution, i.e. quantize values taken by  $\lambda_n$

$$H(X_n | Y^n) = \sum_{\lambda_n} P(\lambda_n) h(\lambda_n)$$
$$h(\lambda_n) = \sum_{x_n \in \{0,1\}} \frac{e^{-x_n \lambda_n}}{1 + e^{-\lambda_n}} \log_2 \frac{e^{-x_n \lambda_n}}{1 + e^{-\lambda_n}}$$

#### Quantized Density Evolution Recursive calculation of $P(\lambda_n, x_n)$



 $\circ\,$  distribution of  $\Lambda_n$  computed recursively, given distribution of  $\Lambda_{n-1}$  and conditional distribution of  $Y_n\,|\,X_n$ 

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Quantized Density Evolution

Recursive calculation of  $P(\lambda_n, x_n)$ 

 $\circ\,$  for the random transmission strategy

$$P(\lambda_n, x_n) = \sum_{\substack{x_{n-1} \in \{0,1\} \\ f(y_n, \lambda_{n-1}) = \lambda_n}} \sum_{\substack{y_n, \lambda_{n-1}: \\ f(y_n, \lambda_{n-1}) = \lambda_n}} P(\lambda_{n-1}, y_n, x_n, x_{n-1})$$
$$= \sum_{\substack{y_n, \lambda_{n-1}: \\ f(y_n, \lambda_{n-1}) = \lambda_n}} P(y_n | x_n) \sum_{\substack{x_{n-1} \in \{0,1\} \\ x_{n-1} \in \{0,1\}}} P(x_n | x_{n-1}) P(\lambda_{n-1}, x_{n-1})$$

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- distribution of  $\Lambda_n$  computed recursively, given distribution of  $\Lambda_{n-1}$  and conditional distribution of  $Y_n | X_n$
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o similar approach for reactive case, relying on myopic surrogate model





#### Outline



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### Average AoI Performance

Random and reactive strategies



 $\circ\,$  for slotted ALOHA access, assuming i.i.d. behavior across slots, average AoI inversely proportional to throughput T

$$\Delta = \frac{1}{2} + \frac{M}{\mathsf{T}}$$

 $\,\circ\,$  minimal AoI attained for transmission probability 1/M

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- $\,\circ\,$  minimal AoI attained for transmission probability 1/M
- $\circ\,$  random transmission strategy shall be operated accordingly, setting  $\alpha=1/M$
- o reactive transmission strategy: transmission probability driven by source statistics

$$\tilde{\alpha} = \pi_0 q_{01} + \pi_1 q_{10}$$

### Average Aol Performance

Symmetric sources





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### Average State Estimation Entropy

Symmetric sources, random and reactive transmission policies





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Symmetric sources, random and reactive transmission policies





### Average State Estimation Entropy

Symmetric sources, random and reactive transmission policies

- SEE increases in larger networks
  - more congestion, higher loss probability
  - $H_\infty$  tends to source entropy as M grows
- lower SEE when sources transition less often  $(q_{10} = q_{01} = 0.01)$ 
  - fewer updates needed to track evolution
  - lower channel congestion (reactive strategy)
- reactive strategy offers better performance
  - only relevant updates are sent
  - for M=2, perfect knowledge at the receiver is achieved



#### Average State Estimation Entropy Asymmetric sources, $q_{10} = 0.1$ , $q_{01} = 0.01$





number of sources, M

### Concluding Remarks



- remote source monitoring in IoT systems
  - two-state Markov sources
  - slotted ALOHA-based channel access, random and reactive strategies
- analytical characterization of state estimation entropy
  - capture uncertainty at receiver
- different protocol operation insights when considering AoI and SEE
  - when freshness is of relevance, transmission strategy maximizing throughput is optimal
  - when uncertainty on source state is important, reactive strategy is convenient

### Concluding Remarks



#### $\circ\,$ metric choice based on targeted application

- freshness relevant e.g., in tracking, monitoring (Aol valuable proxy for complex systems)
- uncertainty on source state relevant for, e.g., actuation
- results also for error probability, false alarm and missed detection probability<sup>16</sup>
- open research topic
  - identify more advanced, optimal policies
  - tracking of more practical, possibly correlated, source processes

<sup>16.</sup> A. Munari, G. Cocco, G. Liva, "Remote Monitoring of Markov Sources over Random Access Channels: False Alarm and Detection Probability," in Proc. IEEE Asilomar 2023.



## Additional Insights

#### State Estimation at the Receiver Estimation Error Probability



- o receiver aims at estimating the state of the reference source
  - denote by  $\widehat{X}_n$  the estimate at time n of the state  $X_n$
- o evaluate performance in terms of state estimation error probability

$$P_e^{(n)} = \mathsf{P}\Big[\widehat{X}_n \neq X_n\Big]$$

• in particular, interested in average error probability

$$P_e = \lim_{N \to \infty} \frac{1}{N} \sum_{n=0}^{N-1} P_e^{(n)}$$

## State Estimation at the Receiver MAP and Decode&Hold



- $\circ$   $P_e$  minimized via maximum a posteriori (MAP) estimator
  - APP logarithmic ratio combined with threshold test
  - $\hat{x}_n = 0$  if  $\lambda_n > 0$ ,  $\hat{x}_n = 1$  if  $\lambda_n < 0$

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- o average error probability of MAP estimator follows from density evolution analysis

$$P_e = \lim_{n \to \infty} \sum_{\lambda_n \le 0} P(\lambda_n, 0) + \lim_{n \to \infty} \sum_{\lambda_n \ge 0} P(\lambda_n, 1)$$

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- MAP estimator entails computational complexity that may be critical, e.g. battery-powered or computationally-limited receivers
  - introduce simpler decode and hold (D&H) estimator

# Decode and Hold Estimator Definition and Trade-Offs



 decode and hold: update estimate only upon receiving update from source, otherwise keep last received value

$$\widehat{X}_n = \begin{cases} Y_n & \text{if } Y_n \in \{0,1\} \\ \widehat{X}_{n-1} & \text{if } Y_n \in \{\mathsf{I},\mathsf{C},\ominus,\oplus\} \,. \end{cases}$$

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- o simpler implementation, yet suboptimal performance
  - consider two-source network, reactive strategy
  - D&H in error as soon as a collision occurs
  - MAP estimator experiences no error: perfect knowledge out of collisions and idle slots

### Decode and Hold Estimator

Markov Model, random transmission strategy



- error probability obtained by tracking Markov process  $(X_n, \widehat{X}_n)$
- $\circ~$  for random strategy, defining  $\omega = \alpha (1-\alpha)^{M-1}$



#### Decode and Hold Estimator Markov Model, random transmission strategy



- error probability obtained by tracking Markov process  $(X_n, \widehat{X}_n)$
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- o aperiodic, irreducible and thus ergodic chain
  - error probability can be derived from stationary distribution  $\pi_{i,j}$ ,  $(i,j) \in \{0,1\} \times \{0,1\}$

$$P_e = \pi_{0,1} + \pi_{1,0}$$

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o computing the stationary distribution, we get

$$P_e = \frac{2q_{01}q_{10}\left(1-\omega\right)}{\left(q_{01}+q_{10}\right)\left[\omega+(1-\omega)(q_{01}+q_{10})\right]}$$

### Decode and Hold Estimator

Markov Model, reactive transmission strategy



- similar approach holds for reactive strategy (myopic approximation)
  - average activation probability  $\tilde{\alpha} = \pi_0 \, q_{01} + \pi_1 \, q_{10}$



#### State Estimation Error Probability MAP and D&H, symmetric sources $(q_{10} = q_{01} = 0.01)$





#### State Estimation Error Probability MAP and D&H, asymmetric sources $(q_{10} = 0.01, q_{01} = 0.1)$





# State Estimation Error Probability MAP and D&H



- $\circ~$  for symmetric sources, D&H matches MAP performance
  - difference only for low  ${\cal M}$
  - · when more sources present, collisions become less informative
  - reactive strategy leads to better estimate

# State Estimation Error Probability MAP and D&H



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  - · when more sources present, collisions become less informative
  - reactive strategy leads to better estimate
- $\circ~$  for asymmetric sources, D&H performs significantly worse
  - D&H may need long time to recover from an error
  - reactive strategy performs worse than random one
    - $_{\circ}\,$  random transmissions can lead to quicker recovery from erroneous estimate