# OpAMP: Linear Operator Approximate Message Passing

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BIRS Workshop Algorithmic Structures for Uncoordinated Communications and Statistical Inference in Exceedingly Large Spaces

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**Power Method:**  
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  Partition rows of the data matrix into J equally-sized submatrices: M = M<sub>1</sub> M<sub>2</sub> : .
- Give each submatrix to a server.







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- Stragglers: What if one or more servers do not respond by the deadline?
- Coded Computing: coding for matrix multiplication with erasures. **Dutta et al. 2016, Lee et al. 2017, Yu et al. 2017** and many more.
- Can we just ignore the missing computations? (We are just refining an estimate.)

# **Running Example: Spiked Matrix Estimation**

• Throughout the talk, we will evaluate our theorems and numerical experiments for the following spiked matrix model:

$$M = \frac{\lambda}{n} \theta \theta^\top + Z$$

where  $\theta \in \mathbb{R}^n$  is the spike and the noise Z is GOE(n).

- Goal: Estimate  $\theta$  with the highest possible correlation  $\frac{1}{n} \langle \theta, \hat{\theta} \rangle$ .
- Recall that  $Z \sim \text{GOE}(n)$  means
  - $Z \in \mathbb{R}^{n \times n}$  is symmetric,
  - independent N(0, 1/n) entries above the diagonal,
  - independent N(0, 2/n) entries on the diagonal.
- This is primarily for direct comparison with prior AMP literature.
- Our theory holds more generally.

# **Distributed Power Method: Ignoring Erasures**



- Row erasures are i.i.d. Bernoulli(0.9).
- Setting the missing entries to zero does not work.

• Concisely summarize erasures via  $\delta_t \in \{0,1\}^n$ 

$$\delta_{t,i} = \begin{cases} 0 & i^{\text{th}} \text{ row of } M \text{ is erased at iteration t} \\ 1 & \text{otherwise} \end{cases}$$

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# $\left(\begin{array}{c} \text{Ignoring Erasures:} \\ x_t = \delta_t \circ M \hat{\theta}_{t-1} \qquad \hat{\theta}_t = \frac{\sqrt{n}}{\|x_t\|} x_t \end{array}\right)$

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• Why not retain the values of the previous iterate in erased coordinates?

One-Step Memory:  

$$x_t = \delta_t \circ M\hat{\theta}_{t-1} + (1 - \delta_t) \circ x_{t-1} \qquad \hat{\theta}_t = \frac{\sqrt{n}}{\|x_t\|} x_t$$

# **Distributed Power Method: One-Step Memory**



• Keeping the past iterate in erased coordinates is much better.

# Related Work: Power Method and Subspace Tracking

- Many variations on this problem have been considered in the literature.
- An incomplete sampling:
  - Noisy Power Method [Hardt and Price 2014, Balcan et al. 2016, Xu and Li 2022]
  - Coordinate-wise Power Method [Lei et al. 2016]
  - Power Method with Momentum [Xu et al. 2018]
  - Adaptive Power Method [Shin et al. 2023]
  - Distributed Streaming PCA [Raja and Bajwa 2020]
  - Communication-Efficient Distributed SVD [Li et al. 2021]
  - Oja's Method [Oja 1982, Oja and Karhunen 1985]
  - Subspace Tracking with Missing Data [Balzano et al. 2018, Wang et al. 2018]
- This Talk: Approximate Message Passing (AMP) perspective on erasures.
  - Per-iteration performance guarantees via coupling to a Gaussian process.
  - (Ultimately) simple correction terms.
  - More efficient computation?

# Approximate Message Passing (AMP)

**Basic AMP:**  $x_t = M f_t(x_{t-1}) - b_t f_{t-1}(x_{t-2})$ 

- Data Matrix:  $M \in \mathbb{R}^{n \times n}$
- Denoising Functions:  $f_t : \mathbb{R}^n \to \mathbb{R}^n$
- Debiasing Coefficient:  $b_t \in \mathbb{R}$
- Early work on AMP was motivated by compressed sensing [Donoho et al. 2009, Bayati and Montanari 2011, Javanmard and Montanari 2013].
- Many other applications to regression, matrix estimation, channel coding, massive random access, etc. See recent survey [Feng et al. 2022].
- Most work has focused on separable denoisers, we follow the framework of [Berthier et al. 2020] that allows non-separable denoisers.

# Approximate Message Passing (AMP)

**Basic AMP:**  $x_t = Z f_t(x_{t-1}) - b_t f_{t-1}(x_{t-2})$ 

- (Centered) Data Matrix:  $Z \in \mathbb{R}^{n \times n}$
- Denoising Functions:  $f_t : \mathbb{R}^n \to \mathbb{R}^n$
- Debiasing Coefficient:  $b_t \in \mathbb{R}$
- Deterministic Initialization:  $f_0 \in \mathbb{R}^n$
- State Evolution:  $\{y_t\}$ , a zero-mean Gaussian process with covariance

$$\begin{aligned} \mathsf{Cov}(y_0) &= \frac{1}{n} \|f_0\|^2 \,\mathrm{I}_n \\ \mathsf{Cov}(y_s, y_t) &= \frac{1}{n} \mathbb{E}[\langle f_s(y_{s-1}), f_t(y_{t-1}) \rangle] \,\mathrm{I}_n \,\,, \quad 0 \le s \le t. \end{aligned}$$

• Assumption 1: Each  $f_t \colon \mathbb{R}^n \to \mathbb{R}^n$  is *L*-Lipschitz continuous and satisfies  $\frac{1}{\sqrt{n}} \|f_t(0)\| \leq C$  where C, L are positive numbers that do not depend on n.

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- Debiasing Coefficient:  $b_t \in \mathbb{R}$

#### Theorem

Suppose Assumption 1 holds,  $Z \sim \text{GOE}(n)$ , and  $b_t = \frac{1}{n} \operatorname{tr}(\mathbb{E}[\mathsf{D}f_t(y_{t-1})])$  Then, for any fixed number of iterations T, there exists a sequence (in n) of couplings between  $x_{\leq T}$  and  $y_{\leq T}$  such that  $\frac{\|x_{\leq T} - y_{\leq T}\|}{\sqrt{n}} \xrightarrow[n \to \infty]{} 0.$ 

# **AMP Power Method:**

$$x_t = M\hat{\theta}_{t-1} - \hat{\theta}_{t-1} \qquad \hat{\theta}_t = \frac{\sqrt{n}}{\|x_t\|} x_t$$

• State evolution provides a "single-letter" characterization of the performance at each iteration.

# **AMP** Power Method



- AMP correction term leads to empirical speedup.
- Accurate predictions from state evolution (SE).

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AMP Power Method with Erasures???  

$$x_t = \delta_t \circ M\hat{\theta}_{t-1} + (1 - \delta_t) \circ \hat{\theta}_{t-1} - \text{correction} \qquad \hat{\theta}_t = \frac{\sqrt{n}}{\|x_t\|} x_t$$

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- Can we find a good correction term that establishes a rigorous state evolution?
- Naïve Approach: Just reuse the Basic AMP correction term.

**AMP Power Method with Erasures???**  
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- Standard correction term is not helpful: slow convergence and no state evolution.
- How can we derive the correction term? Generalize.

# Linear Operator Approximate Message Passing (OpAMP)

**OpAMP with Full Memory:** 

$$x_t = \mathcal{L}_t(Z) f_t(x_0, \dots, x_{t-1}) - \sum_{s < t} B_{ts} f_s(x_0, \dots, x_{s-1})$$

- (Centered) Data Matrix:  $Z \in \mathbb{R}^{n imes n}$
- Linear Operators:  $\mathcal{L}_t(Z) = \sum_{k=1}^K L_{tk} Z R_{tk}$  (non-unique decomposition)
- Denoising Functions:  $f_t : \mathbb{R}^{n \times t} \to \mathbb{R}^n$
- Matrix-Valued Debiasing Coefficients:  $B_{ts} \in \mathbb{R}^{n imes n}$
- Assumption 1: Each  $f_t \colon \mathbb{R}^{n \times t} \to \mathbb{R}^n$  is *L*-Lipschitz continuous and satisfies  $\frac{1}{\sqrt{n}} ||f_t(0)|| \le C$  where C, L are positive numbers that do not depend on n.
- Assumption 2:  $||R_{tk}||_{op}$ ,  $||L_{tk}||_{op} \leq C'$  for all  $t, k \in \mathbb{N}_0$  where C', K are positive numbers that do not depend on n.

# Linear Operator Approximate Message Passing (OpAMP)

• State Evolution:  $\{y_t\}$ , a zero-mean Gaussian process with covariance

$$\mathsf{Cov}(y_s, y_t) = \sum_{l,k=1}^{K} \frac{1}{n} \mathbb{E}[\langle R_{sl} f_s(y_{< s}), R_{tk} f_t(y_{< t}) \rangle] L_{sl} L_{tk}^{\top}$$

#### Theorem

Suppose Assumptions 1 and 2 hold,  $Z \sim \text{GOE}(n)$ , and

$$B_{ts} = \sum_{k,l=1}^{K} \frac{1}{n} \operatorname{tr}(R_{tk} \mathbb{E}[\mathsf{D}_s f_t(y_{< t})] L_{sl}) L_{tk} R_{sl}$$

Then, for any fixed number of iterations T, there exists a sequence (in n) of couplings between  $x_{\leq T}$  and  $y_{\leq T}$  such that  $\frac{\|x_{\leq T} - y_{\leq T}\|}{\sqrt{n}} \xrightarrow[n \to \infty]{} 0.$ 

#### **OpAMP Proof Sketch: "Lifted" Recursion**

• Define a doubly-indexed, full-memory AMP recursion:

$$w_{tk} = Zg_{tk}(w_{
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• Can show 
$$x_t = \sum_{k=1}^{K} L_{tk} w_{tk}$$
.

• State evolution  $u_{tk}$  for  $w_{tk}$  using full-memory AMP [Gerbelot and Berthier 2023].

• Obtain state evolution 
$$y_t = \sum_{k=1}^{K} L_{tk} u_{tk}$$
 for  $x_t$ .

# **Projection AMP**

Projection AMP:  $x_t = \Pi_t \left( Z f_t(x_{t-1}) - \sum_{s < t} b_{ts} f_s(x_{s-1}) \right) + \Pi_t^{\perp} x_{t-1}$ 

- Projection Matrices:  $\Pi_t$  (not necessarily diagonal, nor commuting)
- Scalar Debiasing Coefficients:  $b_{ts} \in \mathbb{R}^{n \times n}$
- Define  $C_{ts} = \begin{cases} \mathbf{I}, & s = t \\ \Pi_t^{\perp} \Pi_{t-1}^{\perp} \cdots \Pi_{s+2}^{\perp} \Pi_{s+1}^{\perp} & 0 \le s < t \end{cases}$

#### Theorem

Suppose Assumption 1 holds,  $Z \sim \text{GOE}(n)$ , and  $b_{ts} = \frac{1}{n} \operatorname{tr}(\mathbb{E}[\mathsf{D}f_t(y_{t-1})]C_{t-1,s}\Pi_s)$ Then, for any fixed number of iterations T, there exists a sequence (in n) of couplings between  $x_{\leq T}$  and  $y_{\leq T}$  such that  $\frac{\|x_{\leq T} - y_{\leq T}\|}{\sqrt{n}} \xrightarrow{\mathrm{pr}}{n \to \infty} 0$ .

**OpAMP Power Method:**  
$$x_{t} = \delta_{t} \circ \left( Mx_{t-1} - \sum_{s < t} b_{ts} \hat{\theta}_{s} \right) + (1 - \delta_{t}) \circ x_{t-1} \qquad \hat{\theta}_{t} = \frac{\sqrt{n}}{\|x_{t}\|} x_{t}$$

• If  $\delta_t$  is elementwise i.i.d. Bernoulli $(\gamma)$ , then we can establish a state evolution by setting the debiasing coefficients to  $b_{ts} = \frac{\sqrt{n}}{\|x_{t-1}\|} p_t(s)$  where

$$p_t(s) = \begin{cases} \gamma (1-\gamma)^{t-s-1} & \text{if } s = 1, 2, \dots, t-1\\ (1-\gamma)^{t-1} & \text{if } s = 0 \end{cases}$$

• State evolution has a simple form.

#### **OpAMP** Power Method: i.i.d. Erasures



- With the OpAMP correction term, we can establish a rigorous state evolution.
- Attains the same fixed point as the AMP power method.

# **OpAMP Power Method: Round Robin Updates**

**OpAMP Power Method:**  
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- Consider now the setting where we deliberately only apply a subblock of the data matrix M at each iteration, to reduce the computational load.
- We partition the row indices  $\{1, \ldots, n\}$  into J equally-sized subsets  $\mathcal{A}_0, \ldots, \mathcal{A}_{J-1}$ .

$$\delta_{t,i} = egin{cases} 1 & i \in \mathcal{A}_{(t ext{ mod } J)} \ 0 & ext{otherwise} \end{cases}$$

• Debiasing coefficients and state evolution have a simple form.

# **OpAMP Power Method: Round Robin Updates**



- Round Robin: Update 0.1 rows per iteration according to a schedule.
- Noticeable speedup compared to stochastic erasures.

# **OpAMP** Power Method: Subgaussian Matrices



• 
$$M = \frac{\lambda}{n} \theta \theta^{\top} + Z$$

• Z symmetric from i.i.d. Rademacher

• 
$$\lambda = \sqrt{2}$$

• 
$$n = 5000$$

- $\theta \sim \operatorname{Unif}(\{\pm 1\}^n)$
- mark = empirical
- line = GOE state evolution

- Empirical performance very similar.
- GOE(n) state evolution no longer a theoretical guarantee.

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• So far, we have plotted the performance with respect to iteration t.

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- For instance, we could track the total number of  $n \times n$  matrix multiplications.
  - Full Matrix: Standard AMP that applies the full matrix.
  - Round Robin: Cycle through fixed subsets of rows of M.
  - Random Update: Randomly selected rows of M applied.



- Round Robin: Uses fewer matrix multiplications to converge.
- Random Update: Sometimes uses fewer matrix multiplications to converge.

# Conclusions

- AMP perspective on the distributed power method with erasures.
  - Simple state evolution and scalar debiasing coefficients.
  - Same fixed point as no-erasure setting.
  - Computational speedup for partial updates.
  - Can also consider other denoisers, e.g., Bayes.
- Theoretical results established by first generalizing to linear operator AMP, which may be useful in other settings.
- Some follow-up questions:
  - Orthogonal ensembles?
  - Adaptive updates?
  - 1st-order methods with erasures?
  - Noise instead of erasures?
  - Connection to SGD speedup?
- Acknowledgments: Thanks to Nicholas Sacco and Viveck Cadambe for valuable discussions on the power method with erasures.